

# Quantum divergences via convex optimization

Omar Fawzi

The logo for Inria, featuring the word "Inria" in a red, cursive script font.The logo for ENS DE LYON, consisting of the letters "ENS" in a bold, black, sans-serif font with horizontal lines through them, and "ENS DE LYON" in a smaller, black, sans-serif font below it.

IPAM Entropy Inequalities, Quantum Information and Quantum Physics workshop

Based on joint works with Peter Brown and Hamza Fawzi  
arXiv:2007.12575 and arXiv:2007.12576

# Rényi divergence

- Let  $p, q \in \mathbb{R}_+^n$  and  $\alpha \geq 0$ , Rényi divergence

$$D_\alpha(p||q) = \frac{1}{\alpha - 1} \log \sum_{i=1}^n p_i^\alpha q_i^{1-\alpha}.$$

- Data-processing inequality:**  $D_\alpha(Mp||Mq) \leq D_\alpha(p||q)$  for a channel  $M$

Consequence of convexity/concavity of  $(a, b) \mapsto a^\alpha b^{1-\alpha}$   
concave for  $\alpha \in [0, 1]$  and convex for  $\alpha > 1$

- How to extend  $D_\alpha$  to positive semidefinite operators  $\rho, \sigma \geq 0$ ?**  
**Requirement:** Data-processing inequality (DPI)  
**Optional:** Retain properties of classical  $D_\alpha$

# Quantum Rényi divergences

Recall:  $D_\alpha(p\|q) = \frac{1}{\alpha-1} \log Q_\alpha(p\|q)$  with  $Q_\alpha(p\|q) = \sum_{i=1}^n p_i^\alpha q_i^{1-\alpha}$

- [Petz, 1986] Petz:  $\bar{Q}_\alpha(\rho\|\sigma) = \text{tr}(\rho^\alpha \sigma^{1-\alpha})$   
DPI for  $\alpha \in [0, 2]$   
Uses: Hypothesis testing [MH, 2011]
- [MDSFT, WWY 2013] sandwiched:  $\tilde{Q}_\alpha(\rho\|\sigma) = \text{tr}\left(\left(\sigma^{\frac{1-\alpha}{2\alpha}} \rho \sigma^{\frac{1-\alpha}{2\alpha}}\right)^\alpha\right)$   
DPI for  $\alpha \in [\frac{1}{2}, \infty]$   
Uses: Hypothesis testing [MO, 2014], Entropy accumulation [DFR, 2016]
- [Matsumoto, 2013] geometric:  $\hat{Q}_\alpha(\rho\|\sigma) = \text{tr}(\sigma \#_\alpha \rho)$   
DPI for  $\alpha \in [0, 2]$   
Uses: Channel capacity bounds [FF, 2019]
- [Donald, 1986] measured:  $D_\alpha^{\text{M}}(\rho\|\sigma) = \max_{\text{measurement } \mathcal{M}} D_\alpha(\mathcal{M}(\rho)\|\mathcal{M}(\sigma))$   
DPI for  $\alpha \in [0, \infty]$   
Uses: Recoverability bounds [BHOS, 2014]

**This talk:** there are more!

# The Rényi divergence as a convex program

For the rest of the talk  $\alpha > 1$

**Observation:**  $Q_\alpha(p\|q) = \inf_{a \geq 0} \left\{ \sum_{i=1}^n a_i : q_i^{1-\frac{1}{\alpha}} a_i^{\frac{1}{\alpha}} \geq p_i \forall i \right\}$

$\{(a, x, y) \in \mathbb{R}_+^3 : x^{1-\frac{1}{\alpha}} a^{\frac{1}{\alpha}} \geq y\}$  is a convex set

$\implies Q_\alpha(p\|q)$  is jointly convex

$\implies D_\alpha(p\|q)$  satisfies DPI

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**Quantum generalization:** Use matrix geometric mean in place of  $x^{1-\frac{1}{\alpha}} a^{\frac{1}{\alpha}}$

For  $A, B > 0$ , the  $\beta$ -geometric mean is [Kubo-Ando]

$$A \#_\beta B = A^{1/2} \left( A^{-1/2} B A^{-1/2} \right)^\beta A^{1/2}$$

For  $\beta \in [0, 1]$ ,  $\{(A, X, Y) \in \text{Pos}(\mathbb{C}^n)^3 : X \#_\beta A \geq Y\}$  is a convex set

$$Q_\alpha^\#(\rho||\sigma) = \inf_{A \geq 0} \{\text{tr}(A) : \sigma \#_{1/\alpha} A \geq \rho\}$$

# The #-divergence

$$D_{\alpha}^{\#}(\rho\|\sigma) := \frac{1}{\alpha - 1} \log Q_{\alpha}^{\#}(\rho\|\sigma) \quad Q_{\alpha}^{\#}(\rho\|\sigma) := \inf_{A \geq 0} \{ \text{tr}(A) : \sigma \#_{1/\alpha} A \geq \rho \}$$

- Convex optimization problem. Semidefinite program when  $\alpha$  rational
- $Q_{\alpha}^{\#}(\rho\|\sigma)$  jointly convex in  $(\rho, \sigma) \Rightarrow$  DPI
- If  $\rho, \sigma$  commute  $Q_{\alpha}^{\#}(\rho\|\sigma) = Q_{\alpha}(\rho\|\sigma)$
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## Interesting properties:

- Dual form useful for computing conditional entropies for quantum correlations (NPA hierarchy)

$$Q_{\alpha}^{\#}(\rho\|\sigma) = \max_{C, D} \left\{ \text{tr}(C\rho) + \text{tr}(D\sigma) \right. \\ \left. \text{s.t. non-commutative poly. constraints on } C \text{ and } D \right\}$$

- Chain rule and additivity properties for quantum channels
- Regularizes to the sandwiched Rényi divergence

## Comparison with geometric $\widehat{D}_\alpha$

We always have  $D_\alpha^\#(\rho\|\sigma) \leq \widehat{D}_\alpha(\rho\|\sigma) = \frac{1}{\alpha-1} \log \text{tr}(\sigma\#_\alpha\rho)$

$$\begin{aligned} \text{tr}(\sigma\#_\alpha\rho) &= \min\{\text{tr}(A) : A \geq \sigma\#_\alpha\rho\} \\ &\geq \min\{\text{tr}(A) : \sigma\#_{1/\alpha}A \geq \rho\} \end{aligned}$$



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Inequality can be strict

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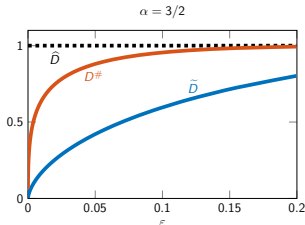
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**Difference can be large:**

- For pure  $\rho = |\phi\rangle\langle\phi|$ ,  $\widehat{D}_\alpha(\rho\|\sigma) = \log(\langle\phi|\sigma^{-1}\phi\rangle) = D_{\max}(\rho\|\sigma)$ .
- Let  $\rho = |\phi\rangle\langle\phi|$  with  $|\phi\rangle = \sqrt{\varepsilon}|00\rangle_{XY} + \sqrt{1-\varepsilon}|11\rangle_{XY}$ , and  $\sigma = I_X \otimes \rho_Y$

$$\widehat{D}_\alpha(\rho\|\sigma) = \begin{cases} 1 & \varepsilon > 0 \\ 0 & \varepsilon = 0 \end{cases}$$



## Comparison with sandwiched $\tilde{D}_\alpha$

**Remark:**  $D_\alpha^\#$  is *not* additive under tensor product, but it is **subadditive**

$$D_\alpha^\#(\rho_1 \otimes \rho_2 \| \sigma_1 \otimes \sigma_2) \leq D_\alpha^\#(\rho_1 \| \sigma_1) + D_\alpha^\#(\rho_2 \| \sigma_2)$$

**Regularization:**

$$\lim_{n \rightarrow \infty} \frac{1}{n} D_\alpha^\#(\rho^{\otimes n} \| \sigma^{\otimes n}) = \tilde{D}_\alpha(\rho \| \sigma)$$

**Proof:** Use measured divergence  $D_\alpha^M(\rho \| \sigma) = \max_{\mathcal{M}} D_\alpha(\mathcal{M}(\rho) \| \mathcal{M}(\sigma))$  and  $\lim_{n \rightarrow \infty} \frac{1}{n} D_\alpha^M(\rho \| \sigma) = \tilde{D}_\alpha(\rho \| \sigma)$

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- Key fact: if  $\rho \leq \rho'$  then  $D_\alpha^\#(\rho \| \sigma) \leq D_\alpha^\#(\rho' \| \sigma)$ .

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Assume  $\sigma = \sum_\lambda \lambda P_\lambda$  spectral decomposition. Consider pinching map

$$\mathcal{P}_\sigma(X) = \sum_{\lambda \in \text{spec}(\sigma)} P_\lambda X P_\lambda.$$

Using pinching inequality  $\rho \leq |\text{spec}(\sigma)| \mathcal{P}_\sigma(\rho)$ :

$$\begin{aligned} Q_\alpha^\#(\rho \| \sigma) &\leq Q_\alpha^\#(|\text{spec}(\sigma)| \mathcal{P}_\sigma(\rho) \| \sigma) \\ &= |\text{spec}(\sigma)|^\alpha Q_\alpha^\#(\mathcal{P}_\sigma(\rho) \| \sigma) \\ &= |\text{spec}(\sigma)|^\alpha Q_\alpha^\#(\mathcal{P}_\sigma(\rho) \| \mathcal{P}_\sigma(\sigma)) \leq |\text{spec}(\sigma)|^\alpha Q_\alpha^M(\rho \| \sigma). \end{aligned}$$

## Extension to channels

Definition of divergence for channels:  $\mathcal{N}_{X' \rightarrow Y}, \mathcal{M}_{X' \rightarrow Y}$  completely positive trace-preserving maps

$$\begin{aligned} D_{\alpha}^{\#}(\mathcal{N} \parallel \mathcal{M}) &:= \sup_{\rho \in \mathcal{D}(X \otimes X')} D_{\alpha}^{\#} \left( (\mathcal{N}_{X' \rightarrow Y} \otimes \mathcal{I}_{X \rightarrow X})(\rho) \parallel (\mathcal{M}_{X' \rightarrow Y} \otimes \mathcal{I}_{X \rightarrow X})(\rho) \right) \\ &= \sup_{\omega_X \in \mathcal{D}(X)} D_{\alpha}^{\#} \left( \omega_X^{\frac{1}{2}} J_{XY}^{\mathcal{N}} \omega_X^{\frac{1}{2}} \parallel \omega_X^{\frac{1}{2}} J_{XY}^{\mathcal{M}} \omega_X^{\frac{1}{2}} \right), \end{aligned}$$

$$J_{XY}^{\mathcal{N}} = (\mathcal{I}_{X \rightarrow X} \otimes \mathcal{N}_{X' \rightarrow Y})(|\Phi\rangle\langle\Phi|_{XX'}) \text{ where } |\Phi\rangle_{XX'} = \sum_i |i\rangle_X \otimes |i\rangle_{X'}$$

Can be written as a convex optimization problem:

$$D_{\alpha}^{\#}(\mathcal{N} \parallel \mathcal{M}) = \frac{1}{\alpha - 1} \log \left\{ \min_{A_{XY} \geq 0} \|\text{tr}_Y(A_{XY})\|_{\infty} \text{ s.t. } J_{XY}^{\mathcal{M}} \#_{1/\alpha} A_{XY} \geq J_{XY}^{\mathcal{N}} \right\}$$

# Main properties

- 1 Sub-additive

$$D_{\alpha}^{\#}(\mathcal{N}_1 \otimes \mathcal{N}_2 \| \mathcal{M}_1 \otimes \mathcal{M}_2) \leq D_{\alpha}^{\#}(\mathcal{N}_1 \| \mathcal{M}_1) + D_{\alpha}^{\#}(\mathcal{N}_2 \| \mathcal{M}_2)$$

- 2 Regularization coincides with that of  $\tilde{D}_{\alpha}$ :

$$\inf_n \frac{1}{n} D_{\alpha}^{\#}(\mathcal{N}^{\otimes n} \| \mathcal{M}^{\otimes n}) = \sup_n \frac{1}{n} \tilde{D}_{\alpha}(\mathcal{N}^{\otimes n} \| \mathcal{M}^{\otimes n}) = \tilde{D}_{\alpha}^{\text{reg}}(\mathcal{N} \| \mathcal{M})$$

**Remark:**  $\tilde{D}_{\alpha}^{\text{reg}}(\mathcal{N} \| \mathcal{M}) > \tilde{D}_{\alpha}(\mathcal{N} \| \mathcal{M})$  in general

- 3 Chain rule:

$$D_{\alpha}^{\#}(\mathcal{N}(\rho) \| \mathcal{M}(\sigma)) \leq D_{\alpha}^{\#}(\mathcal{N} \| \mathcal{M}) + D_{\alpha}^{\#}(\rho \| \sigma).$$

## Application:

Channel discrimination, adaptive strategies do not help in some settings  
See David Sutter's talk on Thursday



## Application: computation of $\tilde{D}^{\text{reg}}(\mathcal{N}||\mathcal{M})$

$\tilde{D}_\alpha^{\text{reg}}(\mathcal{N}||\mathcal{M}) = \lim_{n \rightarrow \infty} \frac{1}{n} \tilde{D}_\alpha(\mathcal{N}^{\otimes n}||\mathcal{M}^{\otimes n})$  can be approximated efficiently

For any  $n \geq 1$

$$\frac{D_\alpha^\#(\mathcal{N}^{\otimes n}||\mathcal{M}^{\otimes n})}{n} - \epsilon(\alpha, d, n) \leq \tilde{D}_\alpha^{\text{reg}}(\mathcal{N}||\mathcal{M}) \leq \frac{D_\alpha^\#(\mathcal{N}^{\otimes n}||\mathcal{M}^{\otimes n})}{n}$$

where

$$\epsilon(\alpha, d, n) = \frac{1}{n} \frac{\alpha}{\alpha - 1} (d^2 + d) \log(n + d)$$

and  $d = \dim X \dim Y$  and  $\alpha > 1$

As  $n \rightarrow \infty$  converging hierarchy of upper bound on  $\tilde{D}_\alpha^{\text{reg}}(\mathcal{N}||\mathcal{M})$

# Conclusion

- Divergences defined in terms of a convex optimization program
- Has complementary properties (chain rule, “nice” variational formulation,...) and useful as a proof tool
- Applications to Shannon theory [arXiv:2007.12576](#) and bounding rates for device-independent cryptographic protocols [arXiv:2007.12575](#)

## Open questions:

- $\lim_{\alpha \rightarrow 1} D_{\alpha}^{\#}$ ?
- Other divergences via optimization?  $\alpha < 1$ ?