

### Simplifying Modeling Complexity In Dynamic Transportation Systems: A State-space-time Network-based Framework

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Prepared for Workshop IV: Decision Support for Traffic Long Program <u>New Directions in Mathematical Approaches for Traffic Flow Management</u> Institute for Pure & Applied Mathematics, UCLA

## Outline

#### **1**. Introduction

large-scale dynamic traffic assignment and simulation

### 2. From simulation to optimization:

- Modeling next-generation of transportation systems
- Key Questions: modeling challenges
- Problem Definition: VRPPDTW
- Methodology: combination of dynamic programming and Lagrangian relaxation

### 3. Extensions

Traffic flow state estimation, and traffic signal control and train timetabling...

## Background

Xuesong Zhou

Pronounced as "Su-song Joe"

Acronym for extending traffic user equilibrium and system optimum models to next generation

Research areas
Simulation-based mesoscopic dynamic traffic assignment
DYNASMART
DTALite (based on simplified kinematic wave model)
Traffic state estimation and prediction
Train routing and scheduling
→ Vehicle routing and scheduling (new)

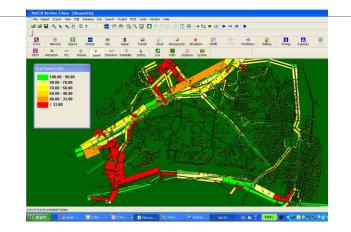
Topic 1: Introduction: large-scale dynamic traffic assignment and simulation

# Open-source Free Software Package

NEXTA: front-end Graphical User Interface GUI (C++)https://github.com/xzhou99/dtalite\_beta\_test

DTALite: Open-source computational engine (C++)

- Light-weight and agent-based DTA
- Simplified kinematic wave model (Newell)
- Built-in OD demand matrix estimation (ODME) program
- Emission prediction (light-weight MOVES interface)
  - Simplified car follow modeling (Newell)



Converting demand flow to vehicles — Neturk Loading for Iteration 16 Pres global path set pres global path set simulation clockiB min, B of vehicles Generated: 0, In network: 0 simulation clockiB min, B of vehicles Generated: 721, In network: 1209 simulation clockiB min, B of vehicles Generated: 1462, In network: 1209 simulation clockiB min, B of vehicles Generated: 1462, In network: 1209 simulation clockiB min, B of vehicles Generated: 285, In network: 1279 simulation clockiB min, B of vehicles Generated: 2865, In network: 1279 simulation clockiB min, B of vehicles Generated: 3685, In network: 1279 simulation clockiB min, B of vehicles Generated: 3685, In network: 1279 simulation clockiB min, B of vehicles Generated: 3685, In network: 1279 simulation clockiB min, B of vehicles Generated: 3685, In network: 1289 simulation clockiB min, B of vehicles Generated: 3687, In network: 1289 simulation clockiB min, B of vehicles Generated: 8577, In network: 1289 simulation clockiB min, B of vehicles Generated: 8577, In network: 1289 simulation clockiB min, B of vehicles Generated: 8577, In network: 1289 simulation clockiB min, B of vehicles Generated: 8577, In network: 1289 simulation clockiB min, B of vehicles Generated: 8577, In network: 1289 simulation clockiB min, B of vehicles Generated: 8577, In network: 1289 simulation clockiB min, B of vehicles Generated: 8577, In network: 1289 simulation clockiB min, B of vehicles Generated: 8577, In network: 1289 simulation clockiB min, B of vehicles Generated: 8577, In network: 1289 simulation clockiB min, B of vehicles Generated: 8577, In network: 1289 simulation clockiB min, B of vehicles Generated: 8577, In network: 1289 simulation clockiB min, B of vehicles Generated: 8577, In network: 1289 simulation clockiB min, B of vehicles Generated: 8577, In network: 1289 simulation clockiB min, B of vehicles Generated: 8577, In network: 1289 simulation clockiB min, B
Processor o is working on shortest path calculation

# Learning Traffic Network Modeling using Open-source tools

#### www.learning-transportation.org

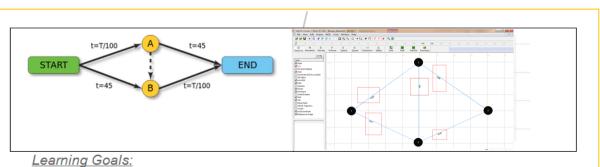
12 lessons:

Lesson 1: Let Us Create a Transportation Network

Lesson 2: From Population to Driving Trips:

Lesson 3: Remove Roads to Speed Travel?

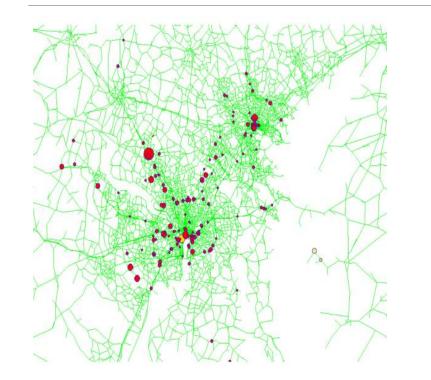
Lesson 4: Optimize Traffic Signal Lights



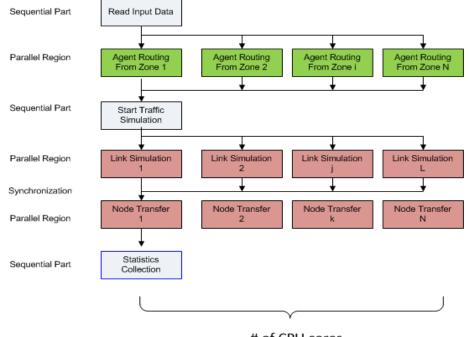
1. Understand modeling principles of user equilibrium, and gap functions

- 2. Know how to setup BPR parameters for special link types
- 3. <u>Understand the impact of random incidents and analyze traffic at link, path and</u> <u>network levels</u>
- 4. <u>Understand different network equilibrium method: method of successive average</u> <u>vs. day to day learning</u>
- 5. <u>Understand the impact of road pricing to resolve Braess's Paradox</u>

# Computational Challenges



Shared memory-based parallel computing for agent-based path finding and mesoscopic traffic simulation (based on OpenMP)

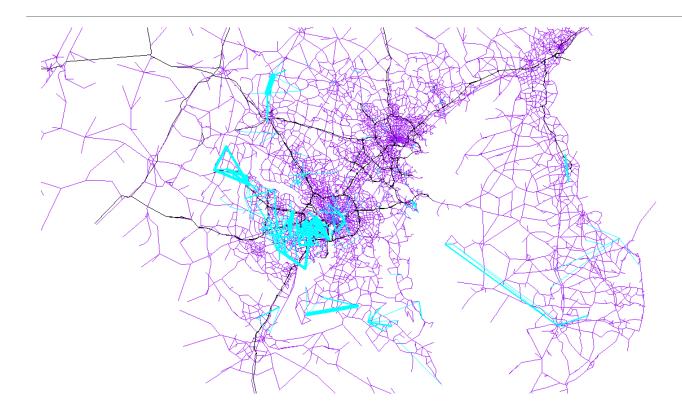


Maryland State-wide model: 20 K nodes, 47K links, 3,000 zones, 18 M agents

# of CPU cores

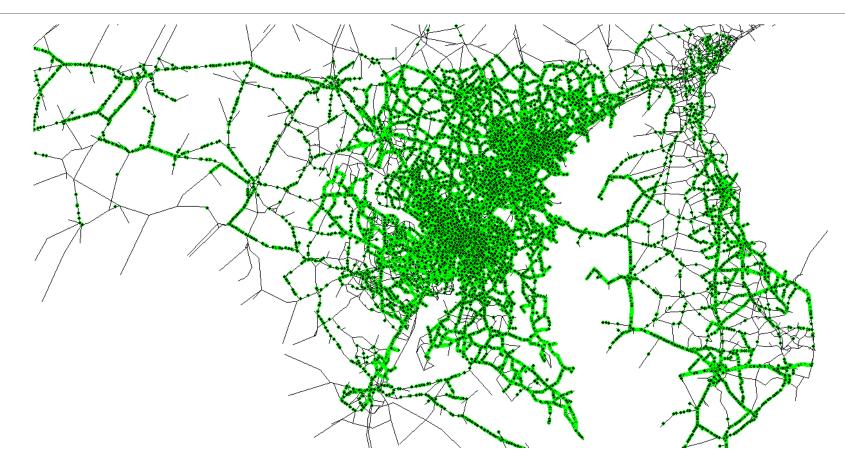
CPU time: 30 min per UE iteration on a 20-core workstation with 194 GB RAM

## Origin-Destination Demand Spatial Distribution Pattern

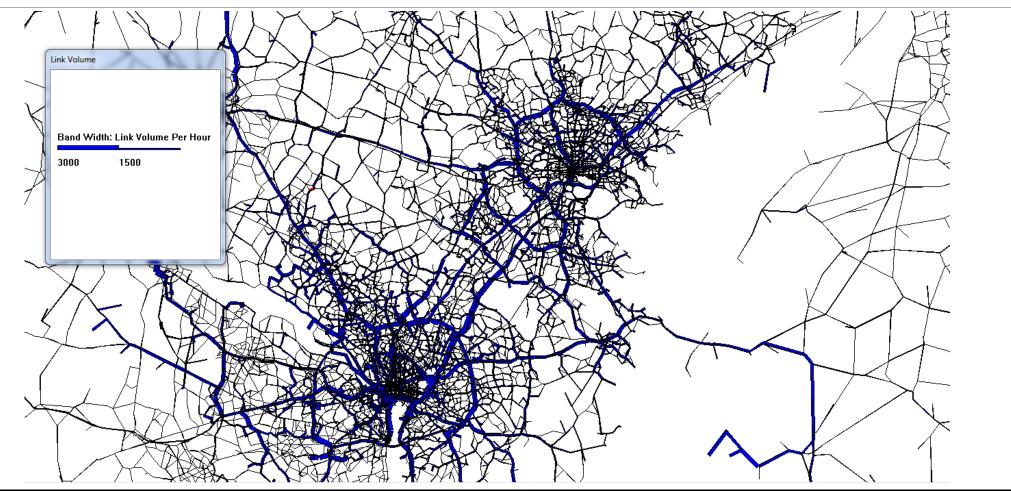


Collaboration with University of Maryland and Maryland <u>State Highway Administration</u> Supported by TRB SHRP II Program

# Vehicle Animation at Network Level

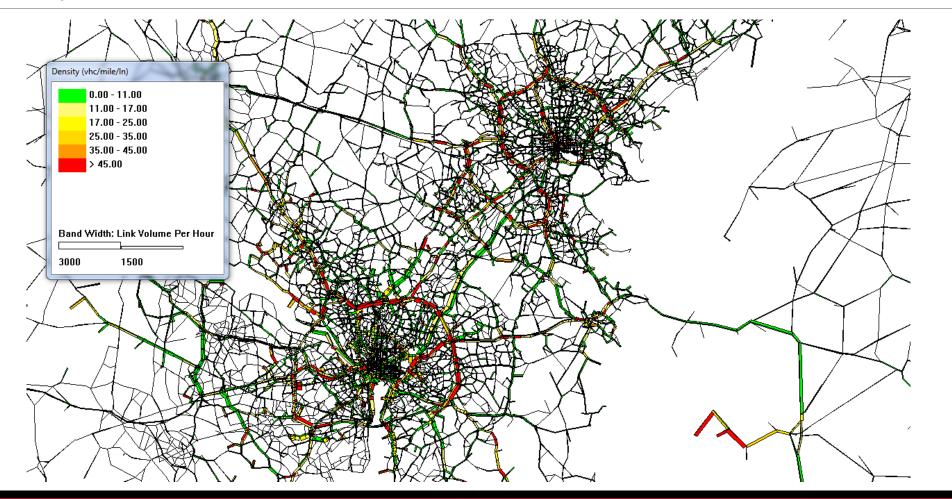


## Volume at Network Level

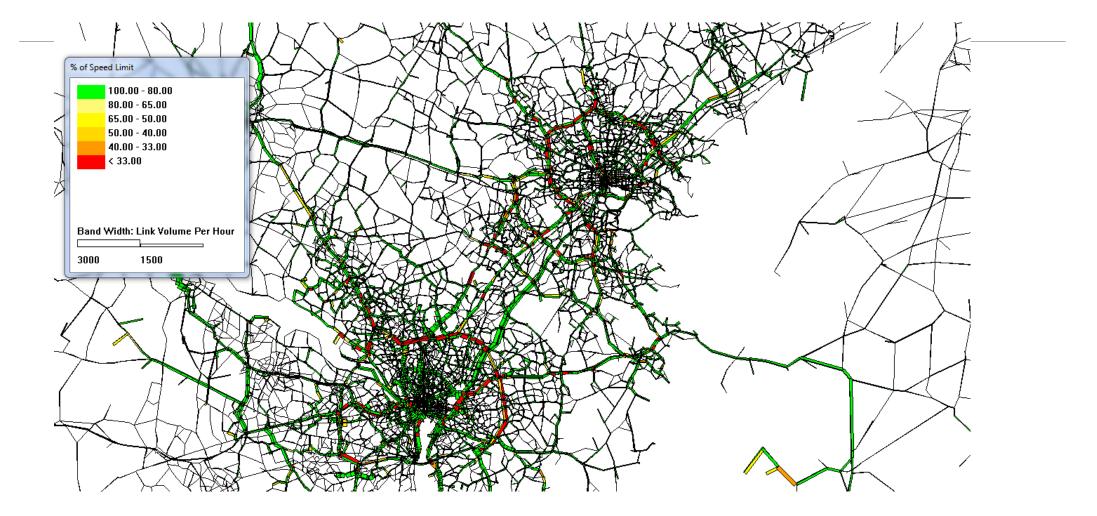


#### Band width of a link is proportional to link volume

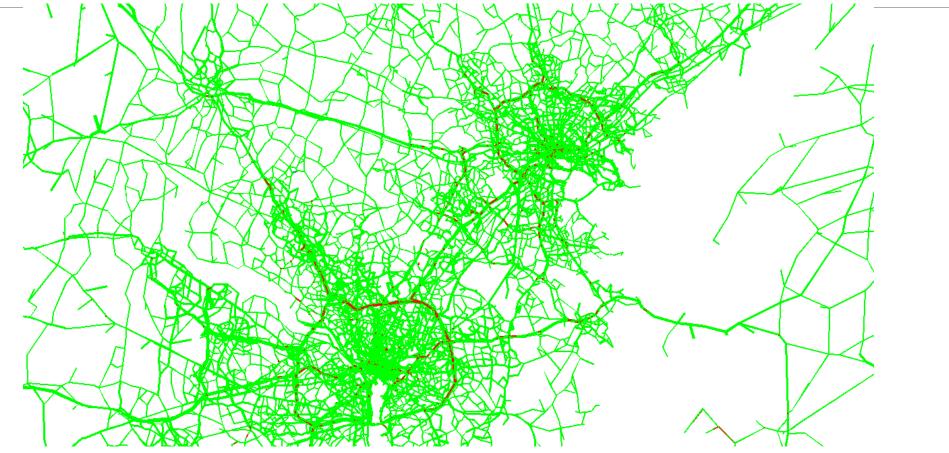
## Density at Network Level



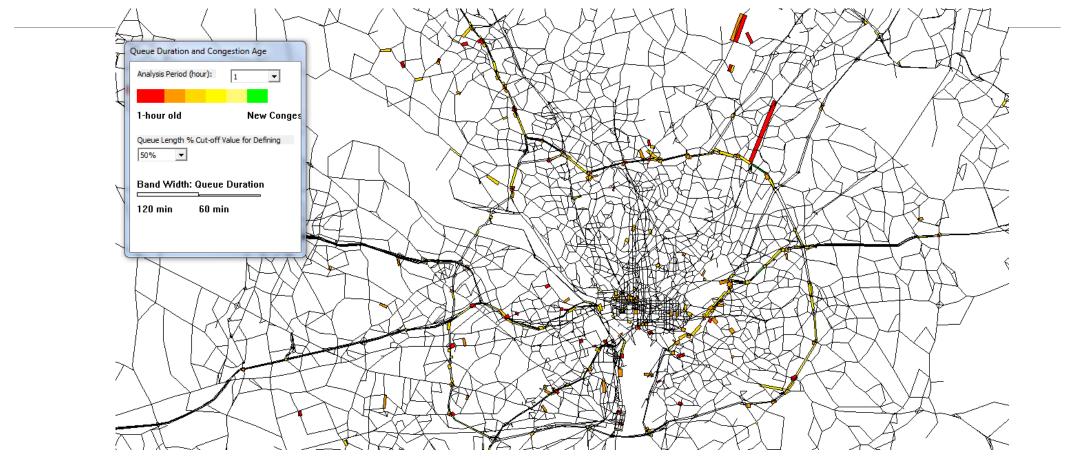
## Speed at Network Level



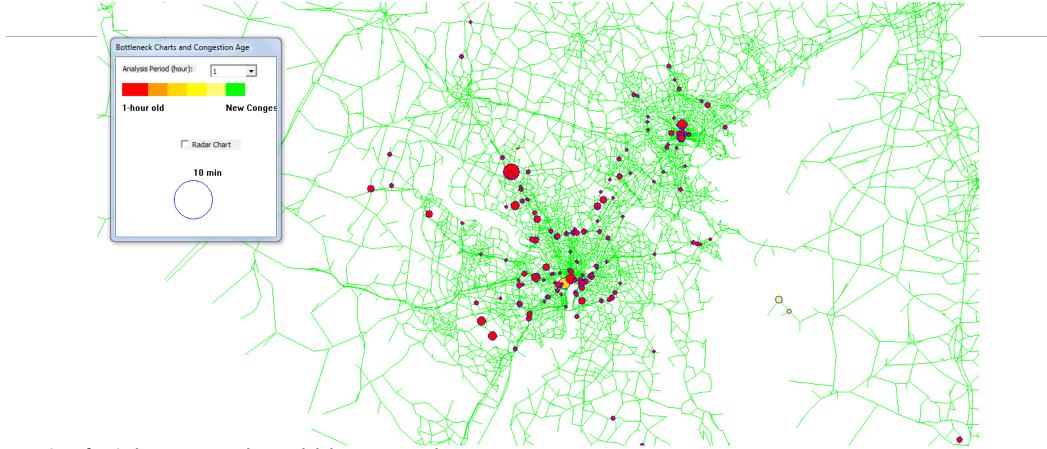
# Queue at Network Level



### Queue Duration at Network Level



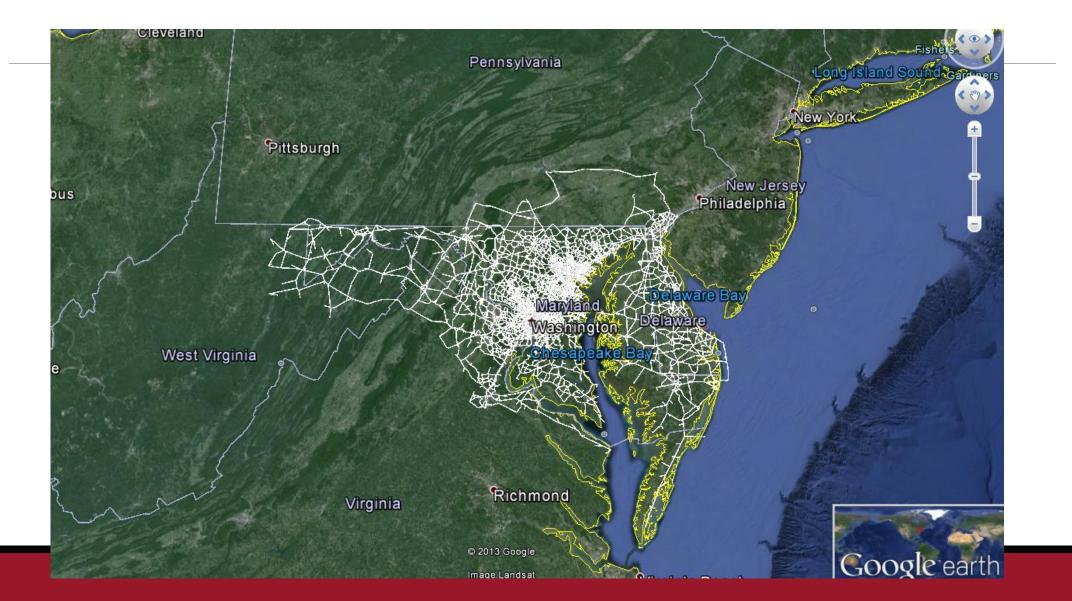
## Time-dependent Bottleneck Locations



Size of a circle represents the total delay at one node

Color of a circle represents the age of congestion (to identify the congestion propagation sequence)

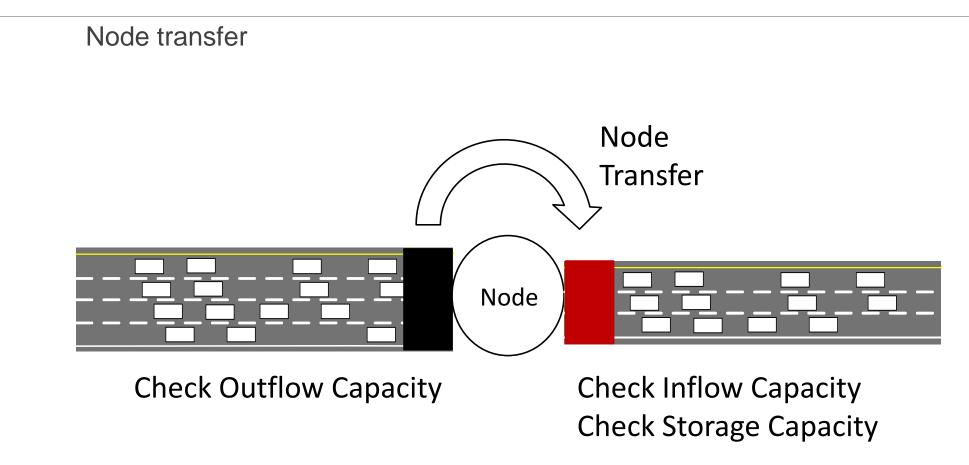
## Statewide Network Coverage in Google Earth



# Volume Display in Google Earth



# Inside: Simplified Event-based Traffic Simulator



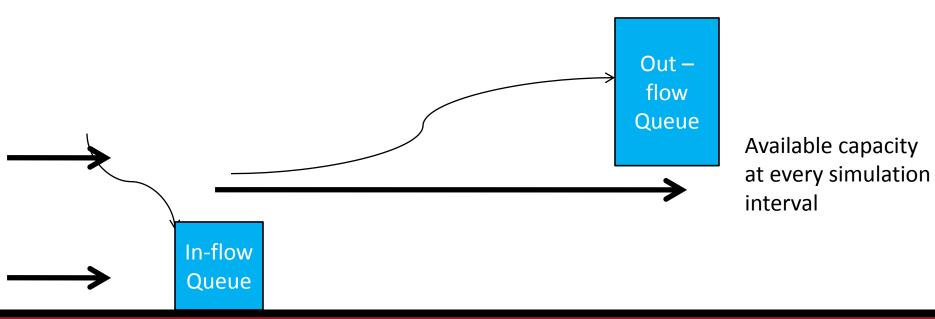
# Multiple Traffic Flow Models

Point queue (relaxed storage constraints)

Spatial queue

Newell's model (i.e. Link Transmission Model)

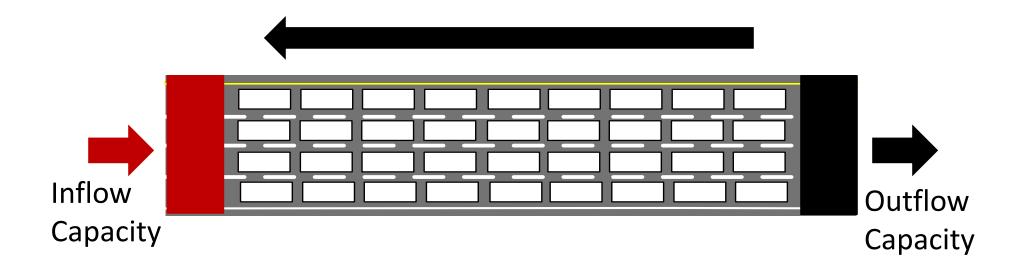
Shockwave propagation



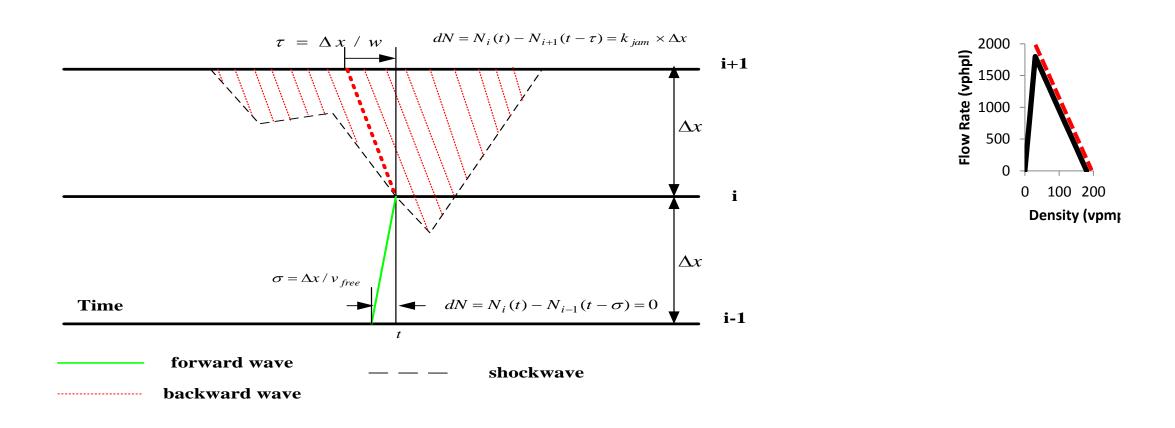
# Traffic Flow Model (on the Link)

Queue propagation

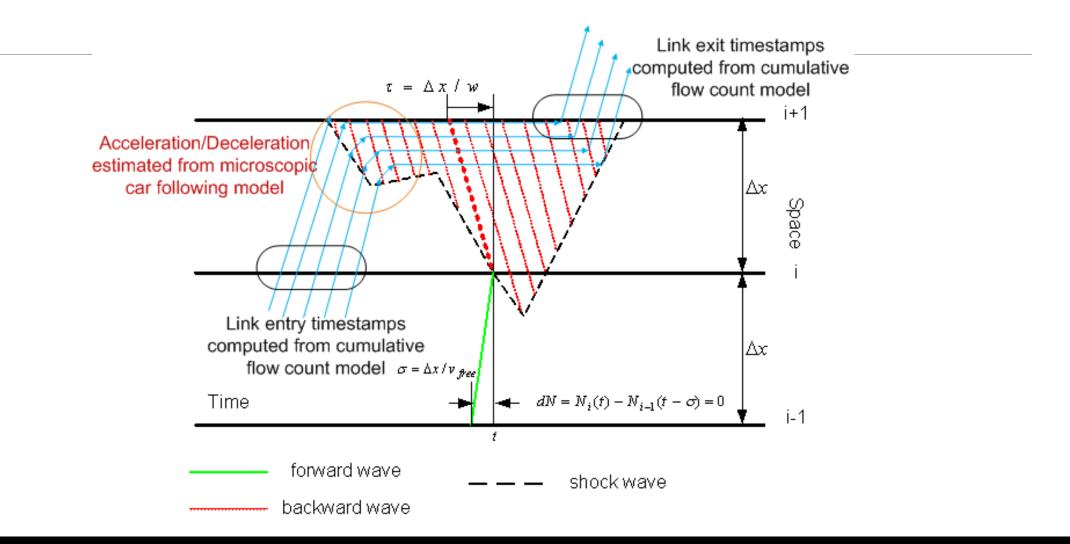
Inflow capacity = outflow capacity



# Illustration of N-Curve Computation For Tracking Queue Spillback

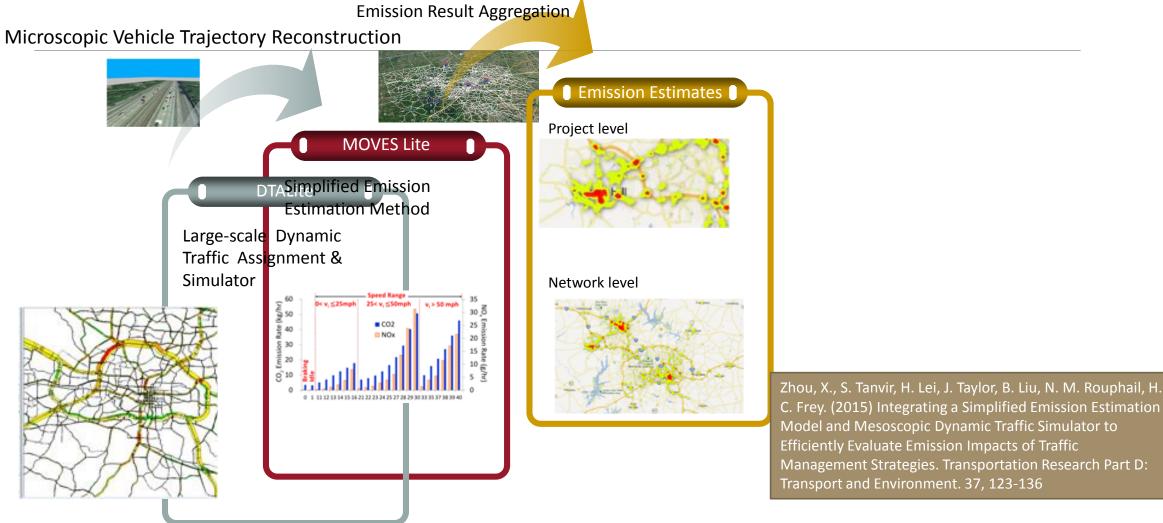


Construct Microscopic Vehicle Trajectory from Mesoscopic Simulation Results using Consistent Simplified Kinematic Wave Model and Simplified Car following Model



White Paper: DTALite: A queue-based mesoscopic traffic simulator for fast model evaluation and calibration: Cogent Engineering (2014): http://www.tandfonline.com/doi/abs/10.1080/23311916.2014.961345

# Mesoscopic Dynamic Traffic Assignment for Emission Evaluation



C. Frey. (2015) Integrating a Simplified Emission Estimation Model and Mesoscopic Dynamic Traffic Simulator to Efficiently Evaluate Emission Impacts of Traffic Management Strategies. Transportation Research Part D: Transport and Environment. 37, 123-136

# Topic 2: Modeling next-generation of transportation systems: from simulation to optimization

Based on Paper titled "Finding Optimal Solutions for Vehicle Routing Problem with Pickup and Delivery Services with Time Windows: A Dynamic Programming Approach Based on State-space-time Network Representations" Monirehalsadat Mahmoudi, Xuesong Zhou

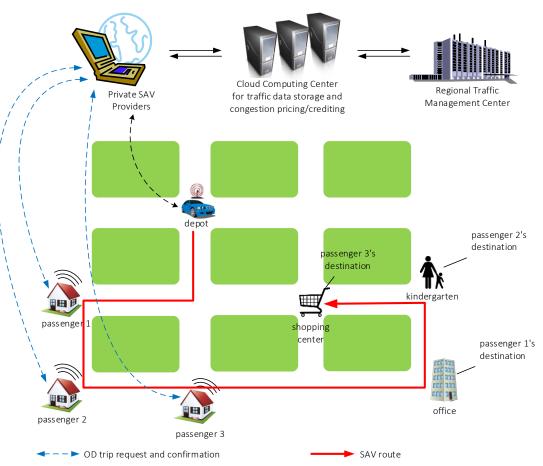
Submitted Transportation Research Part B; http://arxiv.org/abs/1507.02731

### Motivation

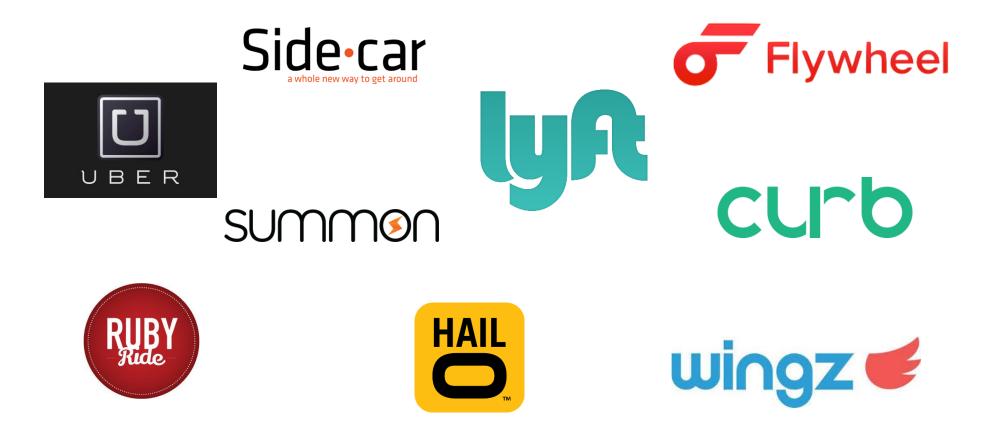
Concept of Ride Sharing:

### Advantages of Ride Sharing:

- Reducing driver stress and driving cost
- Increasing safety
- Increasing road capacity and reducing costs
- Increasing fuel efficiency and reducing pollution

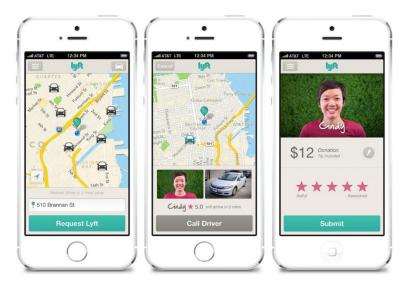


# Ride Sharing Companies



# Ridesharing Apps





One-to-One Matching Innovative Pricing Mechanism Accessibility for Low-income Families

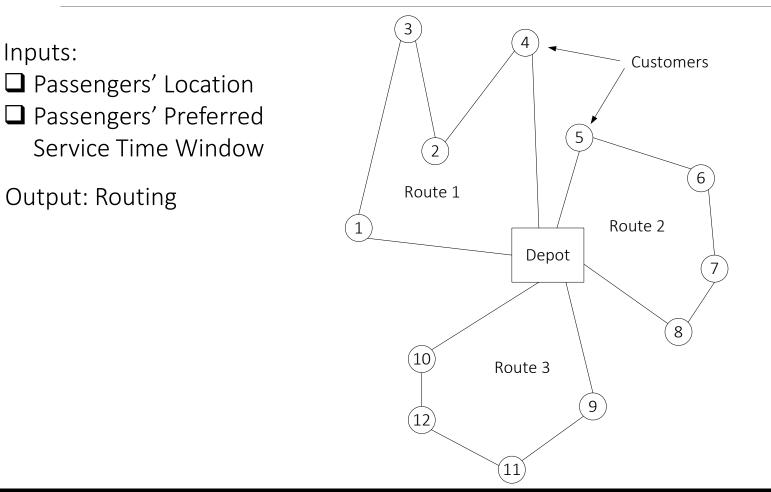
# Key Questions

- How many cars a city should use to support the overall transportation activity demand, at different levels of <u>coordination and pre-trip scheduling</u>?
- How much energy is used at the optimal state (optimal state is a condition in which 100% of travelling is supported by ride sharing)?
- *How much emissions* can be minimized?

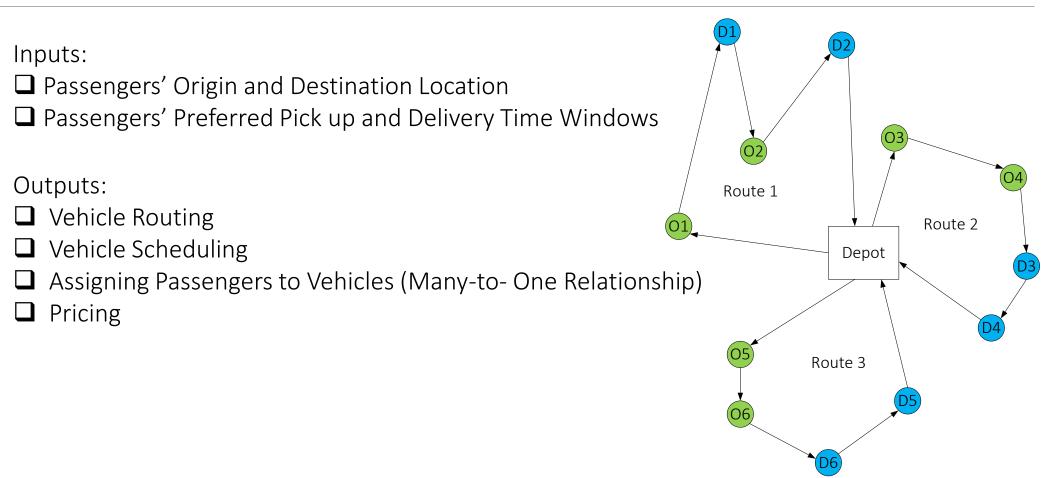
To address the first question:

we propose a new mathematical model for *pickup and delivery problem with time windows (PDPTW)* to present a holistic optimization approach for synchronizing travel activity schedules, transportation services, and infrastructure on urban networks.

# Vehicle Routing Problem (VRP) & Vehicle Routing Problem with Time Windows (VRPTW)



# Vehicle Routing Problem with Pick up and Delivery with Time Windows (VRPPDTW)

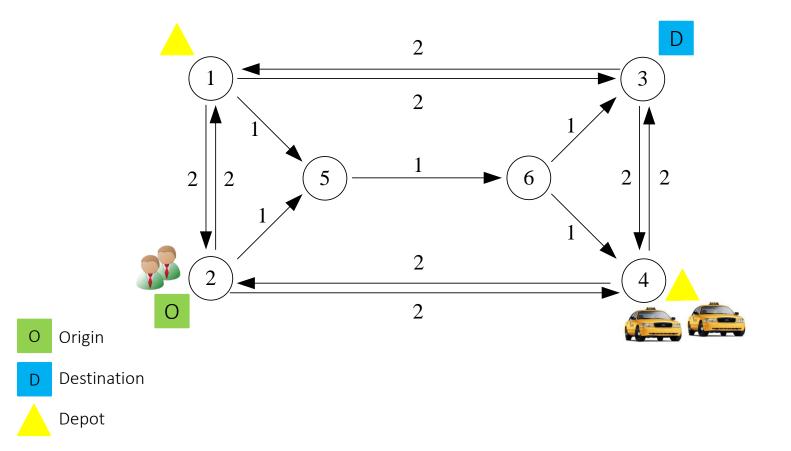


# Optimization-based Approaches for Solving VRPPDTW

Method (Algorithm)	Type of problem	Objective Function
Exact backward dynamic programming	Single vehicle VRPPD	Weighted combination of the total service time and the total customer inconvenience
Forward dynamic programming	Single vehicle VRPPDTW	Sum of waiting and riding times
Benders' decomposition	Single vehicle VRPPD with one sided windows	Total customer inconvenience
Exact forward dynamic programming	single vehicle VRPPDTW	Total distance traveled
Column Generation	Multiple vehicle VRPPDTW	Total travel cost
	Exact backward dynamic programming Forward dynamic programming Benders' decomposition Exact forward dynamic programming	Exact backward dynamic programmingSingle vehicle VRPPDForward dynamic programmingSingle vehicle VRPPDTWBenders' decompositionSingle vehicle VRPPD with one sided windowsExact forward dynamic programmingsingle vehicle VRPPD with one sided windowsExact forward dynamic programmingMultiple vehicle

# Optimization-based Approaches for Solving VRPPDTW (contd.)

Ruland (1995 <i>,</i> 1997)	Polyhedral approach	Single Vehicle VRPPD without capacity constraints	Total travel cost
Savelsbergh and Sol (1998)	Branch-and-price	Multiple vehicle VRPPDTW	Primary objective function: Total number of vehicles, secondary: Total distance traveled
Lu and Dessouky (2004)	Branch-and-cut	Multiple vehicle VRPPDTW	Total travel cost and the fixed vehicle cost
Ropke, Cordeau, Laporte (2007)	Branch-and-cut-and- price	Multiple vehicle VRPPDTW	Total traveled distance
Ropke and Cordeau (2009)	Branch-and-cut-and- price	Multiple vehicle VRPPDTW	Total traveled distance
Baldacci, Bartolini, Mingozzi (2011)	Set-partitioning formulation improved by additional cuts	Multiple vehicle VRPPDTW	Primary objective function: Route costs, secondary: Sum of vehicle fixed costs and then sum of route costs

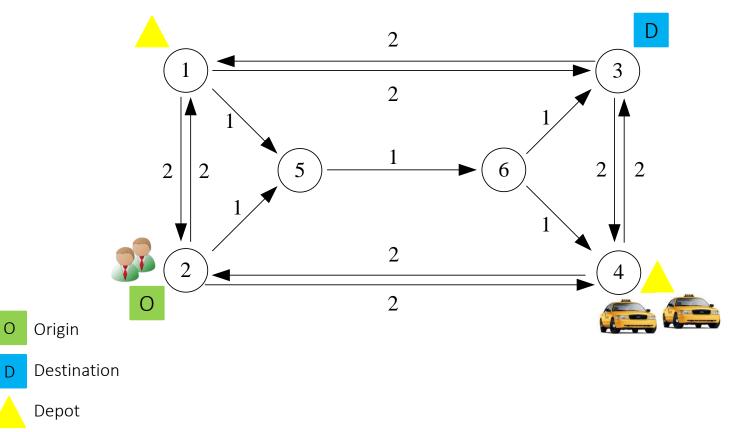


#### Two passengers

#### Two vehicles

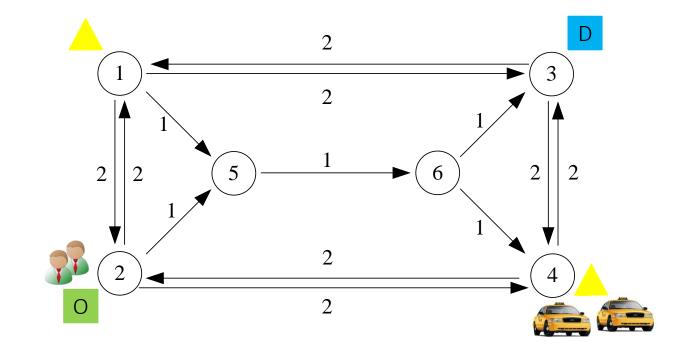
Passenger 1 and 2's origin: node 2 Passenger 1 and 2's destination: node 3

Vehicle 1 and 2's origin depot: node 4 Vehicle 1 and 2's destination depot: node 1

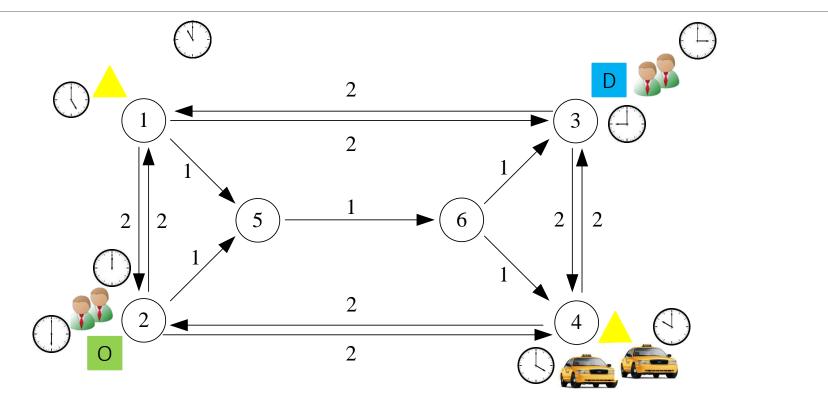


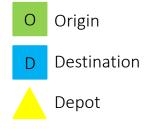
Passenger 1's preferred time window for departure from origin: [4,7] Passenger 1's preferred time for arrival at destination: [9,12]

Passenger 2's preferred time window for departure from origin: [10,12] Passenger 2's preferred time for arrival at destination: [15,17]

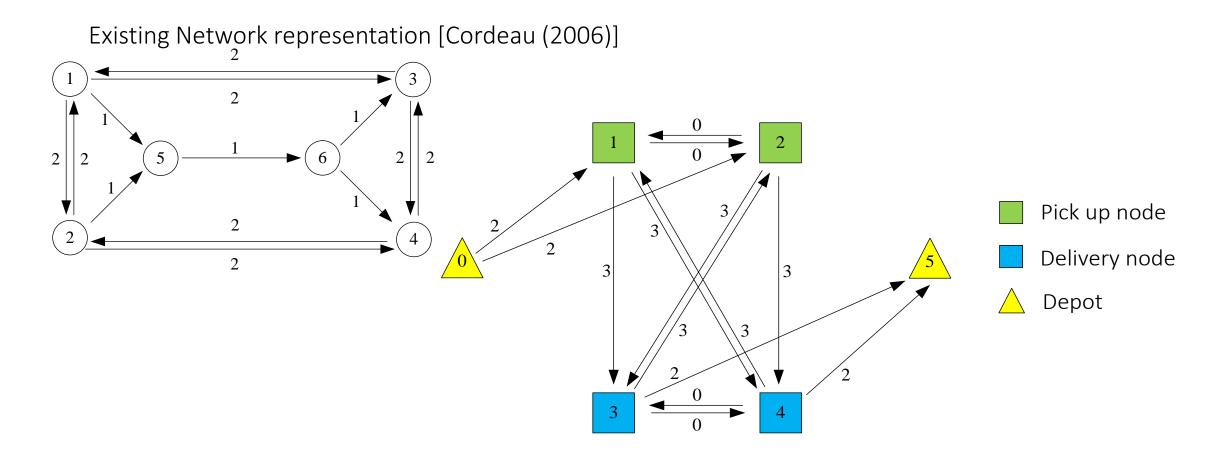








### Opening Statement about Our Method



### Current Mathematical Model for VRPPDTW[Cordeau (2006)]

$Min \sum_{\nu \in V} \sum_{i \in N} \sum_{j \in N} c_{ij}^{\nu} x_{ij}^{\nu}$		objective function: minimizing the total routing cost
$\sum_{v \in V} \sum_{j \in N} x_{ij}^v = 1$	$\forall i \in P$	guarantees that each passenger is definitely picked up
$\sum_{j\in N} x_{ij}^{\nu} - \sum_{j\in N} x_{n+i,j}^{\nu} = 0$	$\forall i \in P, v \in V$	ensure that each passenger's origin and destination are visited exactly once by the same vehicle
$\sum_{j \in N} x_{0j}^{\nu} = 1$	$\forall v \in V$	each vehicle $ u$ starts its route from the origin depot
$\sum_{j\in N} x_{ji}^{\nu} - \sum_{j\in N} x_{ij}^{\nu} = 0$	$\forall i \in P \cup D, v \in V$	flow balance on each node
$\sum_{i\in N} x_{i,2n+1}^{v} = 1$	$\forall v \in V$	each vehicle $ u$ ends its route to the destination depot

### Current Mathematical Model for VRPPDTW[Cordeau (2006)] (contd.)

$x_{ij}^{\nu} \left( B_i^{\nu} + d_i + t_{ij} \right) \le B_j^{\nu}$	Non-linear Constraint	$\forall i \in N, j \in N, v \in V$	validity of the time variables
$x_{ij}^{\nu} (Q_i^{\nu} + q_j) \le Q_j^{\nu}$	Non-linear Constraint	$\forall i \in N, j \in N, v \in V$	validity of the load variables
$L_{i}^{\nu} = B_{n+i}^{\nu} - (B_{i}^{\nu} + d_{i})$		$\forall i \in P, v \in V$	defines each passenger's ride time
$B_{2n+1}^{\nu} - B_0^{\nu} \le T_{\nu}$		$\forall v \in V$	impose maximal duration of each route
$e_i \leq B_i^{\nu} \leq l_i$		$\forall i \in N, v \in V$	impose time windows constraints
$t_{i,n+i} \le L_i^{\nu} \le L$		$\forall i \in P, v \in V$	impose the ride time of each passenge constraints
$max\{0,q_i\} \le Q_i^{\nu} \le min\{Q_{\nu},Q_{\nu}+q_i\}$		$\forall i \in N, v \in V$	impose capacity constraints
$x_{ij}^{\nu} \in \{0,1\}$		$\forall i \in N, j \in N, v \in V$	

### Current Theoretical and Computational Challenges

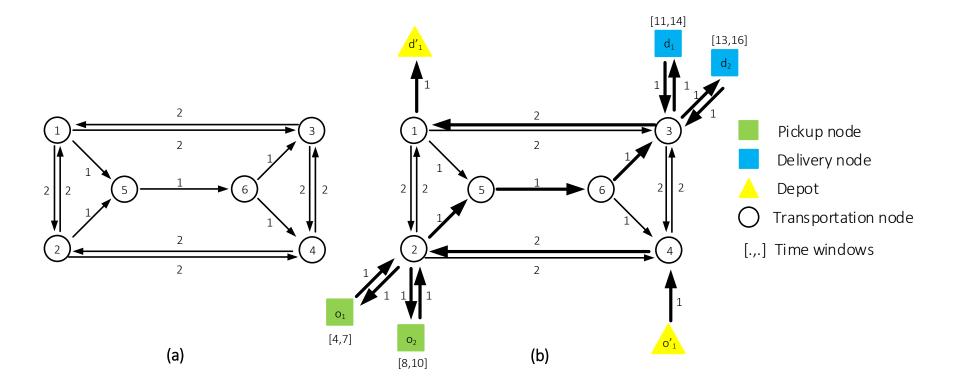
- Single vs multiple vehicles
  - The focus of most research was on solving the PDPTW for a single vehicle (simpler case)
- Single vs multiple depots
- Limited number of transportation requests (passengers)
  - The most Current Algorithm: Baldacci et al. (2011) based on a set-partitioning formulation solved instances of <u>approximately 500</u> requests with tight time windows.
- Time windows
  - Some preprocessing steps to find feasible transportation requests are needed
  - Research only focus on tight time windows to prevent some fluctuation in results

### Current Theoretical and Computational Challenges (contd.)

- Fixed routing cost (travel time) over time
  - Existing network for PDPTW: an offline network in which each link has a fixed routing cost (travel time)
- Existence of sub-tour in the optimal solution
  - Some additional constraints are needed to avoid the existence of any sub-tour in the optimal solution

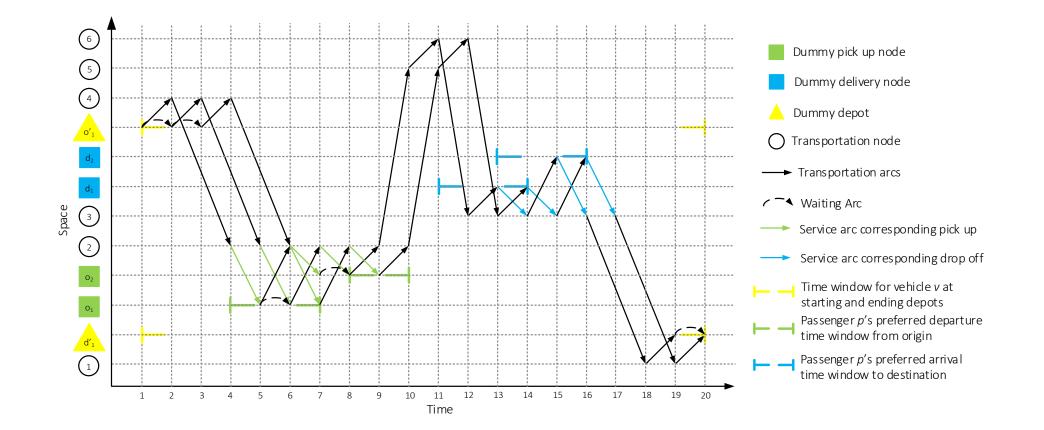
### Opening Statement about Our Method (contd.)

Description of the PDPTW in Space-Time Transportation Network [Mahmoudi, M. & Zhou, X. (2014)]

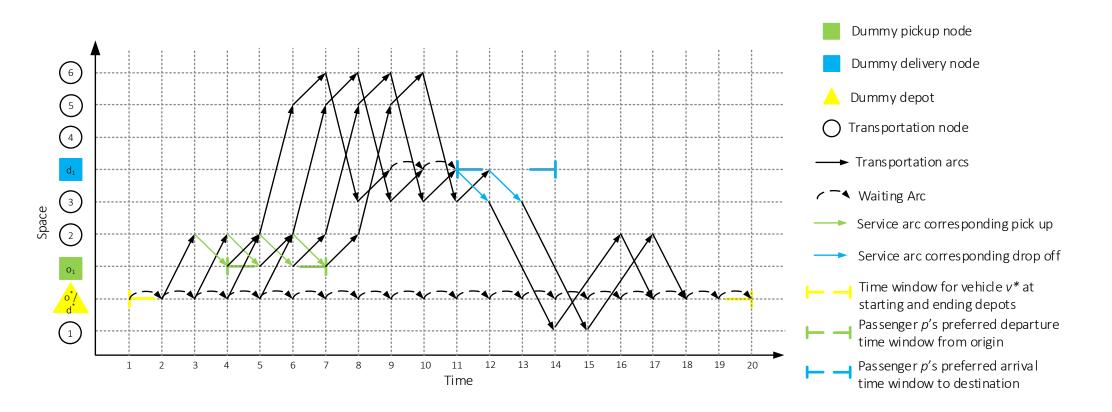


#### http://arxiv.org/abs/1507.02731

### Physical Vehicle's Space-time Transportation Network

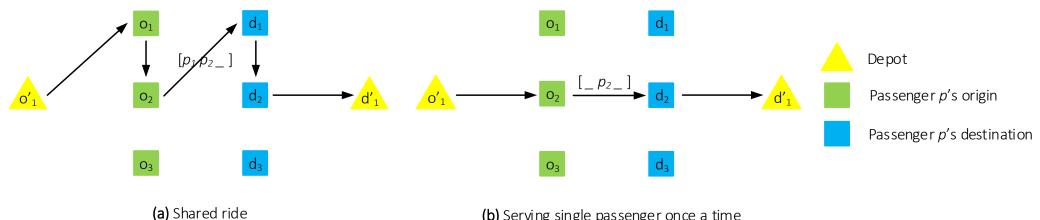


### Virtual Vehicle's Space-time Transportation Network



### Binary representation and equivalent character-based representation for passenger carrying states

Binary	Equivalent character-based		
representation	representation		
[0,0,0]	[]		
[1,0,0]	[p <sub>1</sub> ]		
[0,1,1]	$[_{-}p_{2}p_{3}]$		

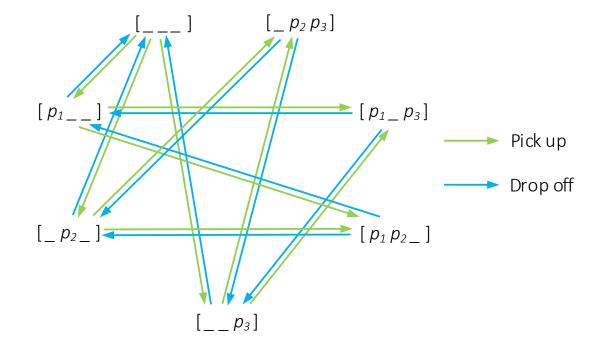


(b) Serving single passenger once a time

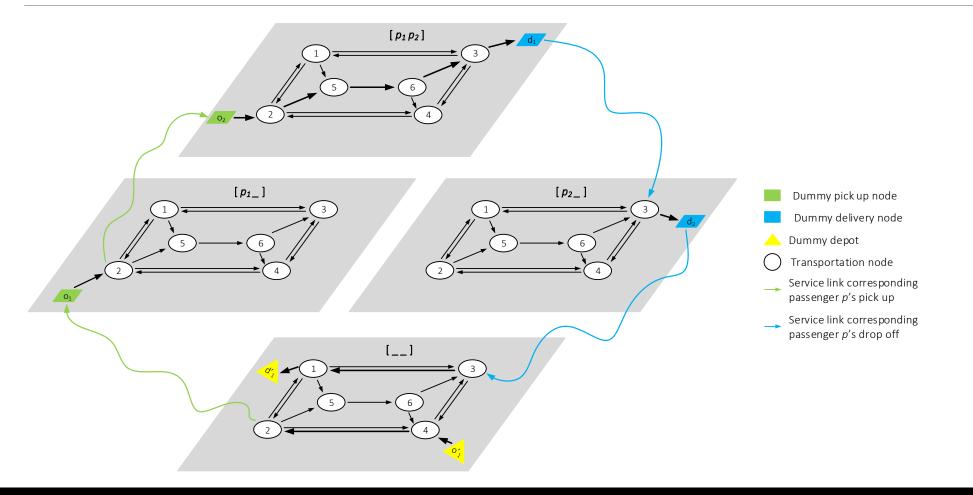
### All possible combinations of passenger carrying states

w w'	[]	[ <i>p</i> <sub>1</sub> ]	[_p <sub>2</sub> _]	[ <i>p</i> <sub>3</sub> ]	[p <sub>1</sub> p <sub>2 _</sub> ]	$[p_1 \_ p_3]$	$[_{-}p_{2}p_{3}]$
[]	no change	pickup	pickup	pickup			
[p <sub>1</sub> ]	drop-off	no change			pickup	pickup	
$[_p_2_]$	drop-off		no change		pickup		pickup
$[-p_3]$	drop-off			no change		pickup	pickup
$[p_1 \ p_2 \ \_]$		drop-off	drop-off		no change		
$[p_{1} _{-} p_{3}]$		drop-off		drop-off		no change	
$[_{-}p_{2}p_{3}]$			drop-off	drop-off			no change

Finite states graph showing all possible passenger carrying state transition (pickup or drop-off)



Projection on state-space network representation for ridesharing path (pick up passenger  $p_1$  and then  $p_2$ ).



### **Problem Definition**

$$Min \ Z = \sum_{v \in (V \cup V^*)} \sum_{(i,j,t,s,w,w') \in B_v} c(v,i,j,t,s,w,w') y(v,i,j,t,s,w,w')$$
  
s.t.

Flow balance constraints at vehicle v's origin vertex

$$\sum_{(i,j,t,s,w,w')\in B_{v}} y(v,i,j,t,s,w,w') = 1 \qquad i = o'_{v}, t = e_{v}, w = w' = w_{0}, \forall v \in (V \cup V^{*})$$

Flow balance constraint at vehicle v's destination vertex

$$\sum_{(i,j,t,s,w,w')\in B_{v}} y(v,i,j,t,s,w,w') = 1 \qquad j = d'_{v}, s = l_{v}, w = w' = w_{0}, \forall v \in (V \cup V^{*})$$

Flow balance constraint at intermediate vertex

$$\sum_{(j,s,w'')} y(v,i,j,t,s,w,w'') - \sum_{(j',s',w')} y(v,j',i,s',t,w',w) = 0 \quad (i,t,w) \notin \{(o'_v,e_v,w_0),(d'_v,l_v,w_0)\}, \forall v \in (V \cup V^*)$$

### Problem Definition (contd.)

Passenger p's pick-up request constraint

$$\sum_{v \in (V \cup V^*)} \sum_{(i,j,t,s,w,w') \in \Psi_{p,v}} y(v,i,j,t,s,w,w') = 1 \qquad \forall p \in P$$

Passenger p's drop-off request constraint

$$\sum_{v \in (V \cup V^*)} \sum_{(i,j,t,s,w,w') \in \Phi_{p,v}} y(v,i,j,t,s,w,w') = 1 \qquad \forall p \in P$$

Binary definitional constraint  $y(v, i, j, t, s, w, w') \in \{0, 1\}$ 

 $\forall (i, j, t, s, w, w') \in B_{v}, \forall v \in (V \cup V^{*})$ 

### Lagrangian Relaxation-based Solution Approach

- Pickup and drop-off constraints : each passenger is picked up and dropped off exactly once by a vehicle (either physical or virtual).
- Flow balance constraints on intermediate nodes force the vehicle to end its route at the destination depot with the empty passenger carrying state.
- o If vehicle v picks up passenger p from his origin, to maintain the flow balance constraints on intermediate nodes, the vehicle must drop-off the passenger at his destination node so that the vehicle comes back to its ending depot with the empty passenger carrying state.
- As a result, constraint (6) is redundant

$$L = \sum_{v \in (V \cup V^*)} \sum_{(i,j,t,s,w,w') \in B_v} c(v,i,j,t,s,w,w')y(v,i,j,t,s,w,w') + \sum_{p \in P} \lambda(p) \left[ \sum_{v \in (V \cup V^*)} \sum_{(i,j,t,s,w,w') \in \Psi_{p,v}} y(v,i,j,t,s,w,w') - 1 \right]$$

### New Relaxed Problem

#### Min L

s.t.

$$\begin{split} \sum_{(i,j,t,s,w,w')\in B_{v}} y(v,i,j,t,s,w,w') &= 1 & i = o'_{v}, t = e_{v}, w = w' = w_{0}, \forall v \in (V \cup V^{*}) \\ \sum_{(i,j,t,s,w,w')\in B_{v}} y(v,i,j,t,s,w,w') &= 1 & j = d'_{v}, s = l_{v}, w = w' = w_{0}, \forall v \in (V \cup V^{*}) \\ \sum_{(j,s,w'')} y(v,i,j,t,s,w,w'') &- & \sum_{(j',s',w')} y(v,j',i,s',t,w',w) &= 0 & (i,t,w) \notin \{(o'_{v},e_{v},w_{0}),(d'_{v},l_{v},w_{0})\}, \forall v \in (V \cup V^{*}) \\ y(v,i,j,t,s,w,w') \in \{0,1\} & \forall (i,j,t,s,w,w') \in B_{v}, \forall v \in (V \cup V^{*}) \end{split}$$

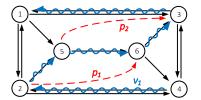
The simplified Lagrangian function L:

$$L = \sum_{v \in (V \cup V^*)} \sum_{(i,j,t,s,w,w') \in B_v} \xi(v,i,j,t,s,w,w') y(v,i,j,t,s,w,w') - \sum_{p \in P} \lambda(p)$$

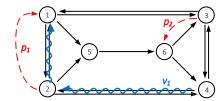
### Time-dependent Forward Dynamic Programming

for each vehicle  $v \in (V \cup V^*)$  do  $L(\ldots) := +\infty;$ node pred of vertex (.,.,.) := -1; time pred of vertex (.,.,.) := -1; state pred of vertex (.,.,.) := -1;  $L(o'_{\nu}, e_{\nu}, w_0) := 0;$ for each time  $t \in [e_v, l_v]$  do for each link (*i*, *j*) do for each state w do derive downstream state w' based on the possible state transition on link (i, j); derive arrival time s = t + TT(i, j, t); if  $(L(i,t,w) + \xi(v,i,j,t,s,w,w') < L(j,s,w'))$  then  $L(j, s, w') := L(i, t, w) + \xi(v, i, j, t, s, w, w')$ ; //label update node pred of vertex (j, s, w') := i; time pred of vertex (j, s, w') := t; state pred of vertex (j, s, w') := w;

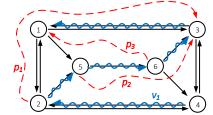
### Computational Results for Testing Different Cases



Scenario I. Two passengers are served by one vehicle through ride-sharing mode.

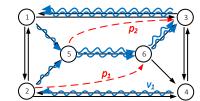


Scenario III. Two passengers and one vehicle; one passenger remains unserved.

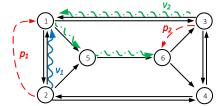


Scenario V. Three passengers are served by one vehicle through ride-sharing mode

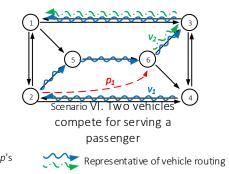
→ Transportation link \_ \_ \_ \_ Representative of passenger p's origin-destination pair



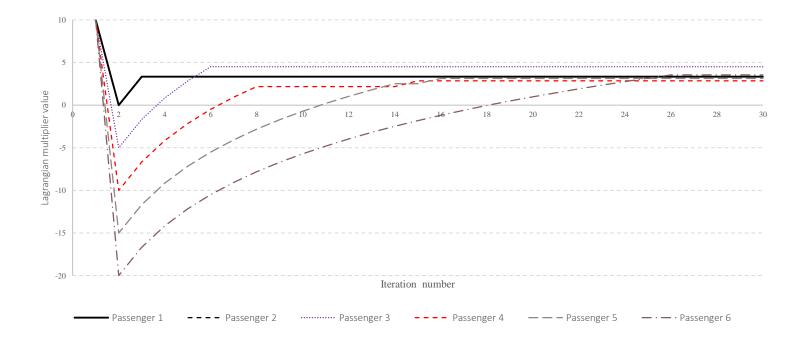
Scenario II. Two passengers are served by one vehicle (not through ride-sharing mode).



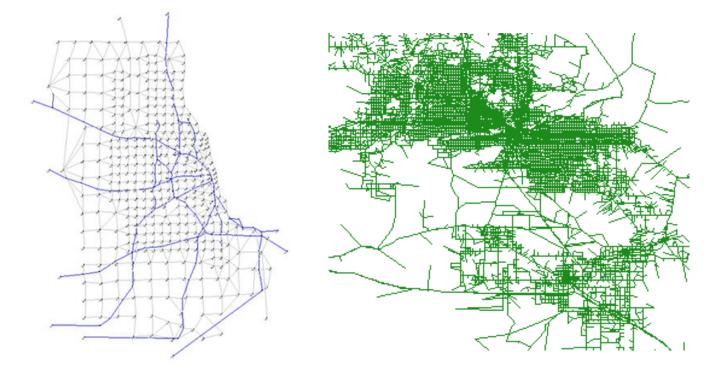
Scenario IV. Two passengers and two vehicles; each vehicle is assigned to a passenger



Lagrangian multipliers along 30 iterations in test case 1 for the six-node transportation network: System Marginal Cost!



# Medium and large-scale transportation networks for computational performance testing



(a) Chicago sketch network

(b) Phoenix metropolitan regional network

# Results for the Chicago network with 933 transportation nodes and 2,967 links

Test case number	Number of iterations	Number of passengers	Number of vehicles	$LB^*$	$UB^*$	Gap (%)	Number of passengers not served	CPU running time (sec)
1	20	2	2	108.43	108.43	0.00%	0	17.43
2	20	11	3	352.97	352.97	0.00%	0	91.87
3	20	20	5	616.66	626.18	1.52%	1	327.51
4	20	46	15	1586.81	1664.07	4.64%	2	4681.52
5	20	60	15	1849.98	1878.55	1.52%	3	7096.50

# Results for the Phoenix network with 13,777 transportation nodes and 33,879 links

Test case number	Number of iterations	Number of passengers	Number of vehicles	$LB^*$	$UB^*$	Gap (%)	Number of passengers not served	CPU running time (sec)
1	6	4	2	70.95	70.95	0.00%	0	110.39
2	6	10	5	191.55	207.05	7.49%	1	398.37
3	6	20	6	310.37	310.37	0.00%	0	1323.18
4	6	40	12	622.23	622.23	0.00%	0	3756.505
5	6	50	15	784.07	784.07	0.00%	0	6983.189

### Short Summary

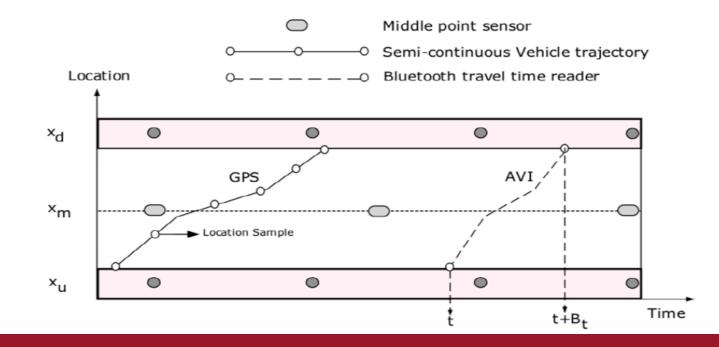
- We propose a new mathematical formulation in state-space-time network for PDPTW
- Based on time-dependent forward dynamic programming approach in the Lagrangian reformulation framework, the main problem is transformed to easy sub-problems (Time-dependent least cost path sub-problems) which is solved independently without much effort
- Unlike former proposed models for PDPTW, this model is now able to solve PDPTW in large scale transportation networks.

Topic 3: Extensions of state-space-time modeling framework: Traffic flow state estimation, and traffic signal control and train timetabling...

## Application 1: Traffic State Estimation

How much information is sufficient?

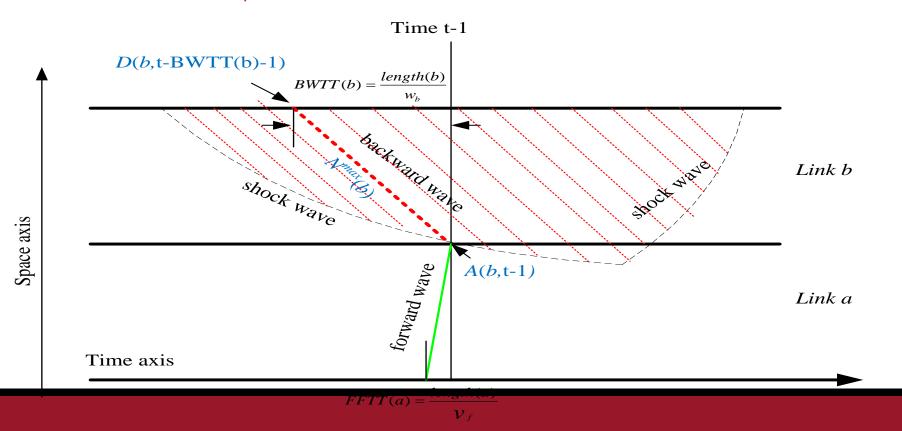
- How to locate point sensors on a traffic segment?
- How to locate Bluetooth reader locations?
- How much AVI/GPS market penetration rate is sufficient?



# Space-state-time network: N(x,t): state N as cumulative flow counts

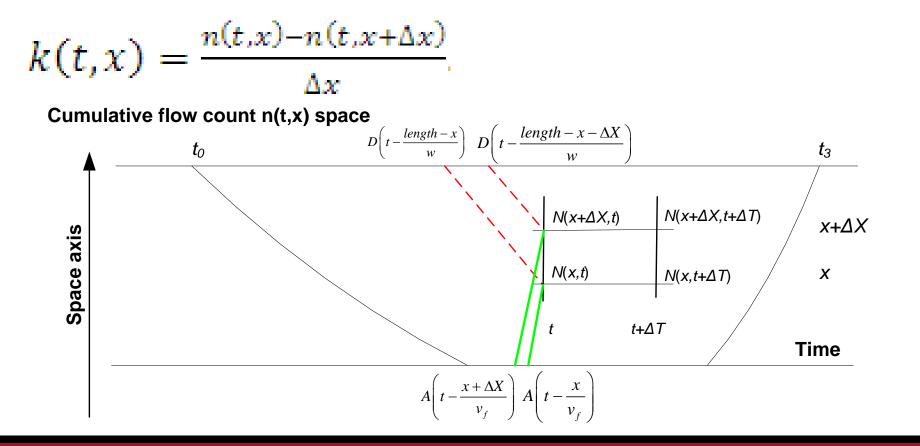
Dr. Newell's three-detector model provides a unified framework

• N(t,x)=Min {N<sub>upstream</sub>(t-BWTT)+Kjam\*distance, N<sub>downstream</sub>(t-FFTT)}



### 1: From Point Sensor Data to Boundary N-curves

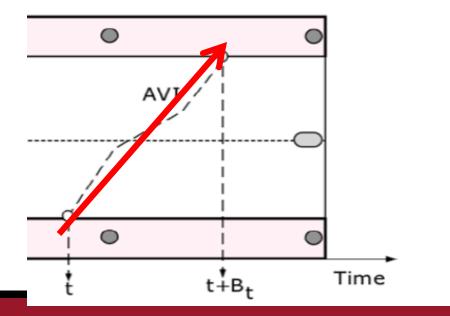
Cell density and flow are all functions of cumulative flow counts



#### 2: From Bluetooth Travel Time to Boundary N-curves

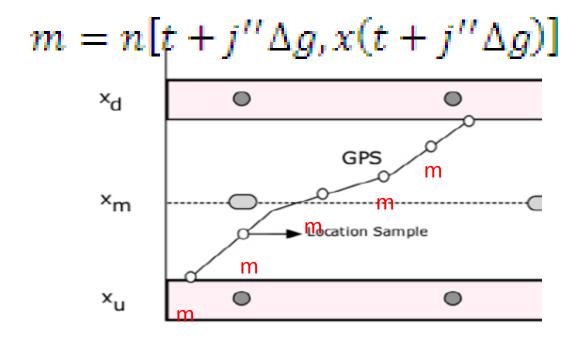
Downstream and upstream N-Curves between two time stamps are connected

$$n_u(t) = n_d(t + B_t)$$



#### 3: From to GPS Trajectory Data to Boundary N-curves

Under FIFO conditions, GPS probe vehicle keeps the same N-Curve number (say m)



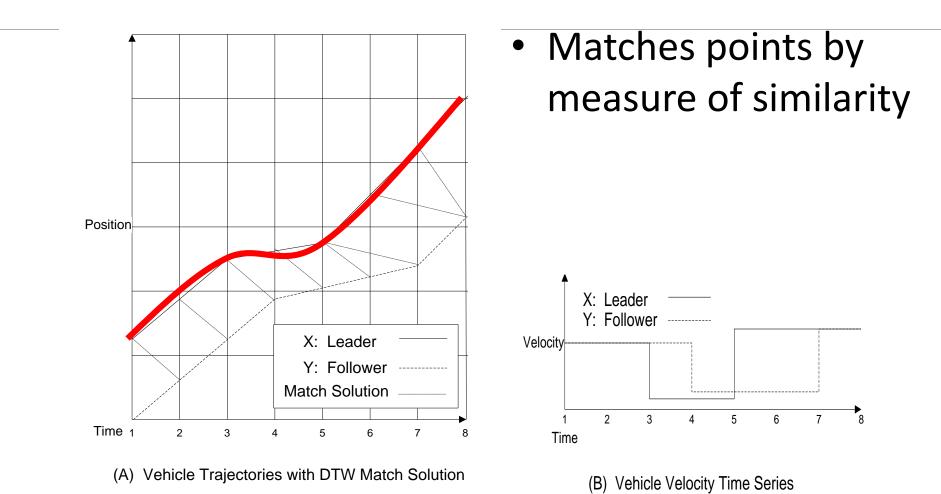
## Stochastic 3-detector Model

All sensors have errors  $\rightarrow$  error propagation

Deng, W. Lei H. ,Zhou, X. (2013) Freeway Traffic State Estimation and Uncertainty Quantification based on Heterogeneous Data Sources: A Three Detector Approach. Transportation Research Part B. 57, 132-157

From single segment to corridor

Lei, H., & Zhou, X. (2014). Linear Programming Model for Estimating High-Resolution Freeway Traffic States from Vehicle Identification and Location Data. Transportation Research Record: Journal of the Transportation Research Board, 2421, 151-160. Application 2: State-space-time path → State as speed for trajectory Use Dynamic Time Warping (DTW) to Estimate dynamic car following model



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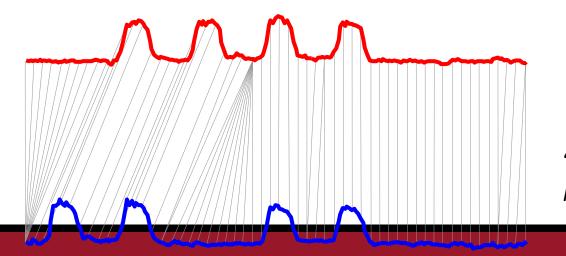
Reference: Eamonn Keogh

Computer Science & Engineering Department University of California - Riverside

### Euclidean Vs Dynamic Time Warping

#### Euclidean Distance

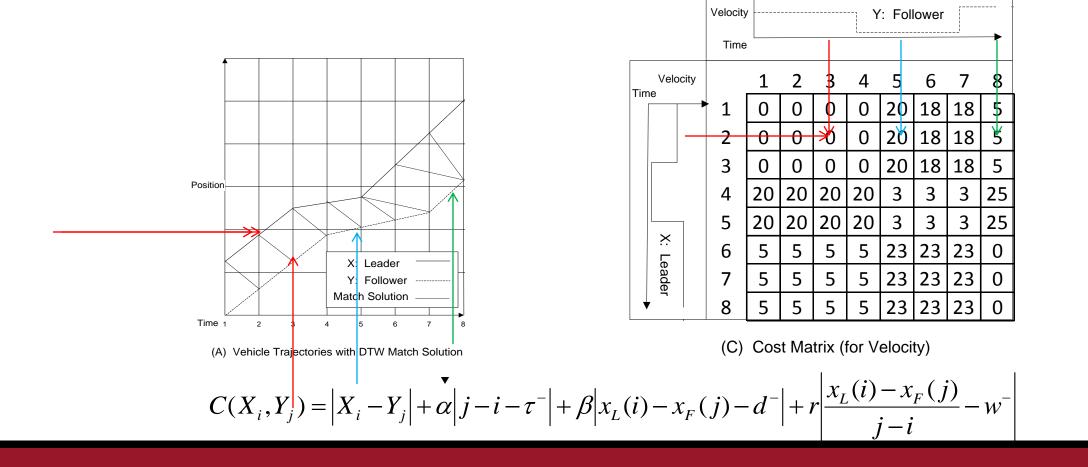
Sequences are aligned "one to one".



### "Warped" Time Axis

Nonlinear alignments are possible.

### Construct Cost Matrix for Traffic Trajectory Matching

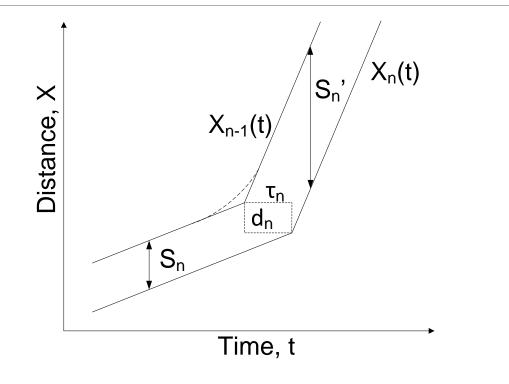


## Application to Newell's Model

Follower separated by leader by reaction time and critical jam spacing

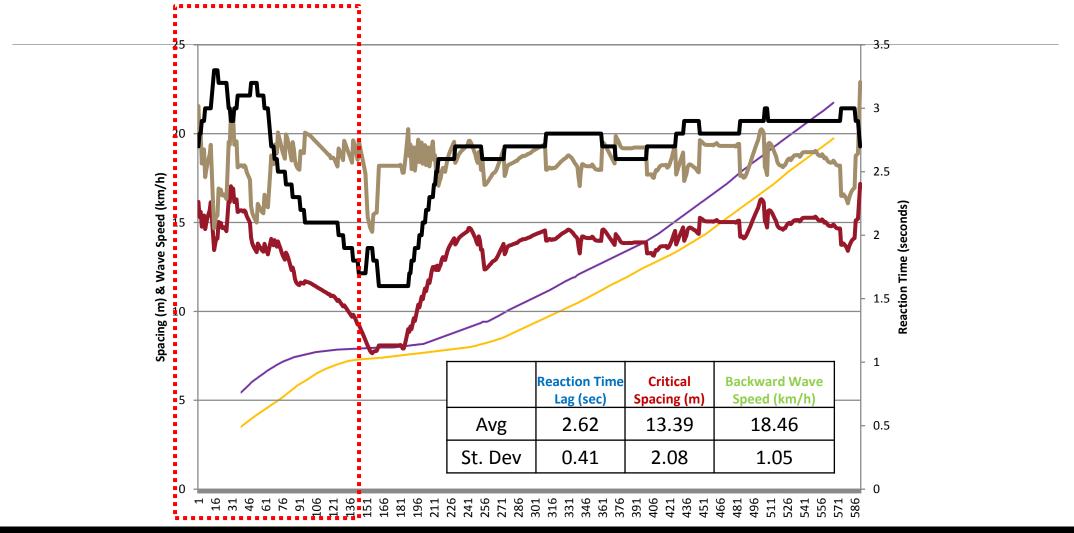
Algorithm finds optimal  $\tau_n$  (time lag) for best velocity match

 Calculate d<sub>n</sub> for all time steps along the trajectory



 $x_n(t+\tau_n) = x_{n-1}(t) - d_n$ 

### Calibrated Parameters: Car 1737



Critical Jam Spacing — Backward Wave Speed — Reaction Time

## NGSIM Data: I-80 Lane 4

Taylor, J., Zhou, X. Rouphail, N., Porter, R.J. (2015) Method for investigating intradriver heterogeneity using vehicle trajectory data: A Dynamic Time Warping approach. Transportation Research Part B, 73, 59-80 Reaction Time Distribution

### Application 3: Phase-time network for Signal Optimization

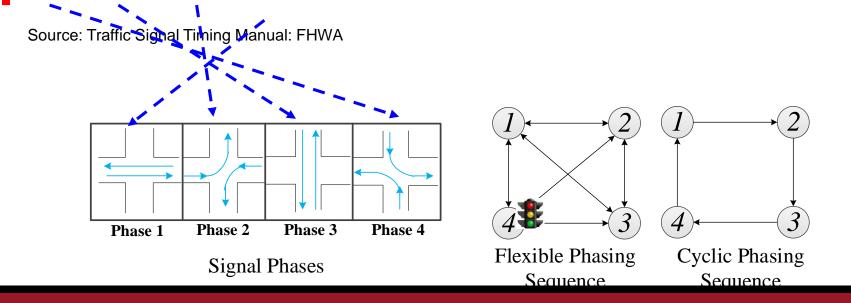
Traffic signal phasing sequence representation

• NEMA ring structure signal phase (North America)



P. Li., P. Mirchandani, X. Zhou, Solving Simultaneous Route Guidance and Traffic Signal Optimization Problem Using Coupled Space-time and Phase-time Networks. Transportation Research Part B. 81, 103-130.

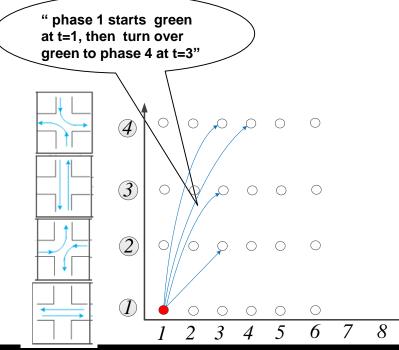
• Movement-based signal phase (same control flexibility, fewer variables)



### Phase-time network for Signal Optimization

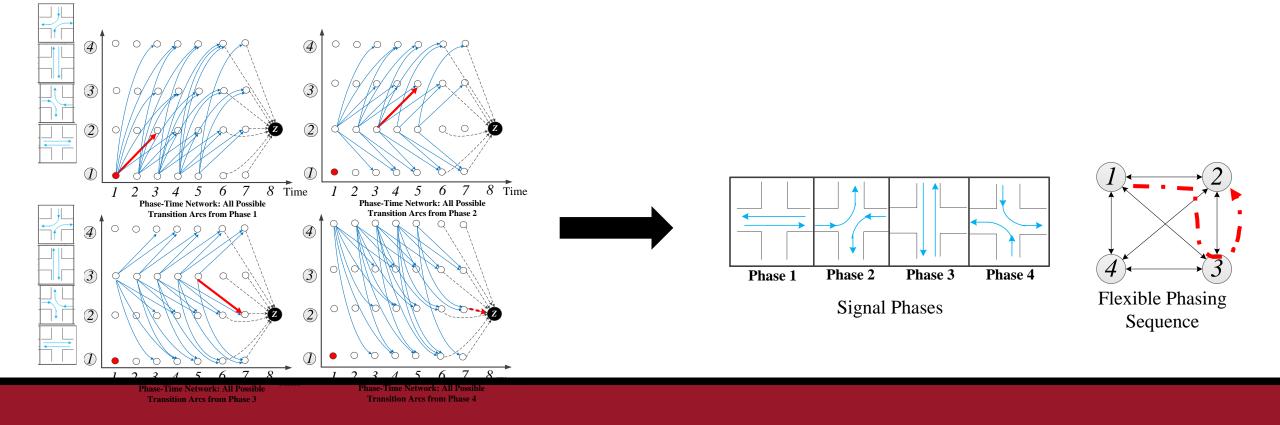
- A signal timing plan is composed of:
  - Phasing sequence
  - Phase duration
- Signal timings can also be represented with a series of phase nodes in the 2-D phase-time network
   "phase 1 starts green at t=1 then turn over
- Similar structure with a space-time network!
- Solution is provided like:

 $w_{(m,k,\tau,h)} = \begin{cases} 1, \text{ if signal phase } m \text{ is green at } \tau \\ \text{and then turns green to signal phase } k \text{ at } h \\ 0, \text{ otherwise} \end{cases}$ 



### Optimize traffic signal in phase-time network

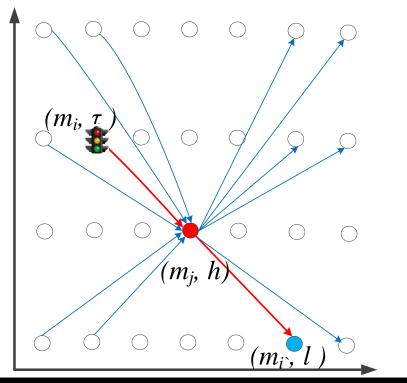
- Each green phase will generate cost on other phases (e.g., delays)
- Find a least-cost path from origin (starting phase) to destination (end of horizon)



### Intersection control constraints

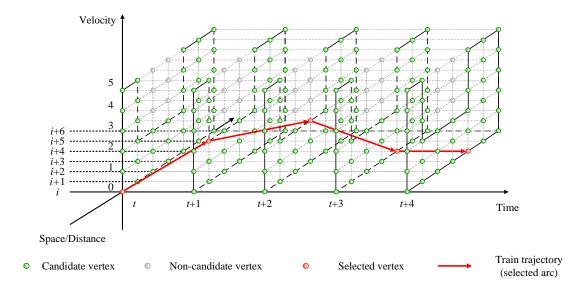
Mutual exclusiveness of signal phases

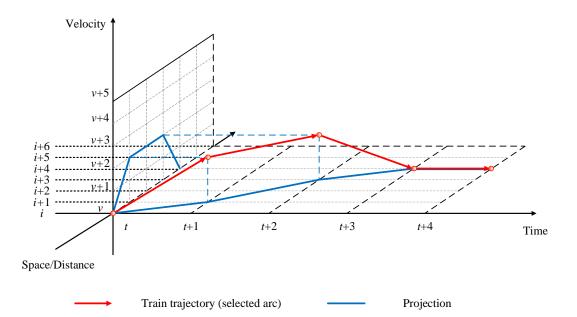
• Any signal phase has one and only one predecessor phase and successor phase at one time.





# Application 4: Speed-space-time network for high-speed train timetabling and speed control





# Conclusions

Present a state-space-time (SST) based modeling framework

By adding additional state dimensions, we prebuild many complex constraints into a multidimensional network

- Computationally efficient solution algorithm: forward dynamic programming + Lagrangian relaxation
- Wide range of applications
- (i) how to estimate macroscopic and microscopic freeway traffic states from heterogeneous measurements,
- (ii) how to optimize transportation systems and ride-sharing services involving vehicular routing decisions with pickup and delivery time windows (VRPPDTW).

#### Challenges

- Selecting state is an art...
- How to overcome curse of dimensionality
- Dynamics and uncertainty (demand/supply)
- Smart search space reduction and metaheuristics algorithms for real-time applications