Motion Coordination for Multi-Agent Networks

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Francesco Bullo



Mechanical Engineering Center for Control, Dynamical Systems and Computation University of California at Santa Barbara http://motion.mee.ucsb.edu

Ack: Anurag Ganguli, Sara Susca, Ketan Savla Jorge Cortés (UCSC), Emilio Frazzoli (UCLA), Sonia Martínez (UCSD) AFOSR MURI Appers, ARO MURI Swarms, NSF Sensors, ONR YIP

What kind of systems?

Groups of systems with control, sensing, communication and computing

Individual members in the group can

- sense its immediate environment
- communicate with others
- process the information gathered
- take a local action in response

Example networks from biology and engineering

Biological populations and swarms



Wildebeest herd in the Serengeti

Geese flying in formation

Atlantis aquarium, CDC Conference 2004

Multi-vehicle and sensor networks embedded systems, distributed robotics

Distributed information systems, large-scale complex systems intelligent buildings, stock market, self-managed air-traffic systems Useful engineering through small, inexpensive, limited-comm vehicles/sensors

Problem	lack of understanding of how to assemble and co- ordinate individual devices into a coherent whole		
Distributed feedback	rather than "centralized computation for known and static environment"		
Approach	integration of control, comm, sensing, computing		

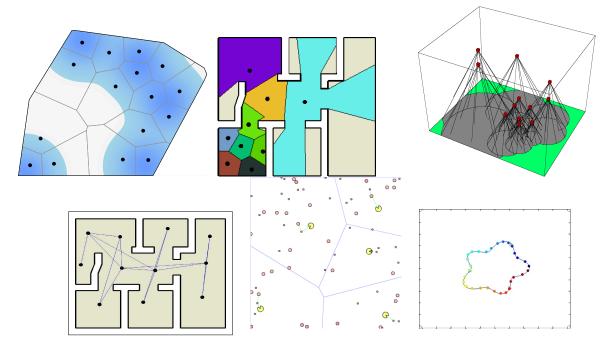
Research in Animation

(i) elementary motion tasks

deployment, rendezvous, flocking, self-assembly

(ii) sensing tasks

detection, localization, visibility, vehicle routing, search, plume tracing



- I: Models for Multi-Agent/Robotic Networks: tools and modeling results
- II: Motion Coordination: algorithms for multiple tasks rendezvous, deployment
- III: Sensing Tasks: sensing problems

target servicing, boundary estimation

Part I: Models for Multi-Agent Networks

References

- (i) I. Suzuki and M. Yamashita. Distributed anonymous mobile robots: Formation of geometric patterns. *SIAM Journal on Computing*, 28(4):1347–1363, 1999
- (ii) N. A. Lynch. Distributed Algorithms. Morgan Kaufmann Publishers, San Mateo, CA, 1997. ISBN 1558603484
- (iii) D. P. Bertsekas and J. N. Tsitsiklis. Parallel and Distributed Computation: Numerical Methods. Athena Scientific, Belmont, MA, 1997. ISBN 1886529019
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Objective

- (i) meaningful + tractable model
- (ii) feasible operations and their cost
- (iii) control/communication tradeoffs

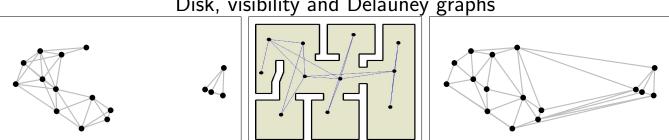
Part I: Synchronous robotic network

A uniform/anonymous robotic network S is

(i) $I = \{1, \dots, N\}$; set of unique identifiers (UIDs)

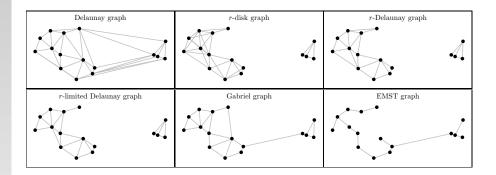
(ii) $\mathcal{A} = \{A_i\}_{i \in I}$, with $A_i = (X, U, X_0, f)$ is a set of identical control systems; set of physical agents

(iii) communication graph



Disk, visibility and Delauney graphs

Communication models for robotic networks



Relevant graphs

- (i) fixed, balanced
- (ii) geometric or state-dependent
- (iii) switching
- (iv) random, random geometric

Message model message, packet, bits; absolute or relative positions

Control and communication law

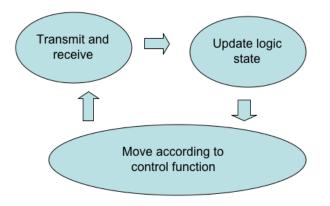
(i) communication schedule(ii) communication language(iii) set of values for logic variables

(iv) message-generation function(v) state-transition functions(vi) control function

 $\mathbb{T} = \{t_\ell\}_{\ell \in \mathbb{N}_0} \subset \overline{\mathbb{R}}_+$

L including the null message \ensuremath{W}

$$\begin{split} \mathsf{msg} \colon \mathbb{T} \times X \times W \times I &\to L \\ \mathsf{stf} \colon \mathbb{T} \times W \times L^N &\to W \\ \mathsf{ctrl} \colon \overline{\mathbb{R}}_+ \times X \times W \times L^N \to U \end{split}$$



- Coordination task is $(\mathcal{W}, \mathcal{T})$ where $\mathcal{T}: X^N \times \mathcal{W}^N \to \{\texttt{true}, \texttt{false}\}$
- For {S, T, CC}, define costs/complexity: control effort, communication packets, computational cost
- \bullet time complexity to achieve ${\mathcal T}$ with ${\mathcal C}{\mathcal C}$

$$\begin{aligned} \mathrm{TC}(\mathcal{T},\mathcal{CC},x_0,w_0) &= \inf \left\{ \ell \mid \mathcal{T}(x(t_k),w(t_k)) = \texttt{true}, \text{ for all } k \geq \ell \right\} \\ \mathrm{TC}(\mathcal{T},\mathcal{CC}) &= \sup \left\{ \mathrm{TC}(\mathcal{T},\mathcal{CC},x_0,w_0) \mid (x_0,w_0) \in X^N \times \mathcal{W}^N \right\} \\ \mathrm{TC}(\mathcal{T}) &= \inf \left\{ \mathrm{TC}(\mathcal{T},\mathcal{CC}) \mid \mathcal{CC} \text{ achieves } \mathcal{T} \right\} \end{aligned}$$

Example tasks / control objective

Motion: deploy, gather, flock, reach pattern Logic-based: achieve consensus, synchronize, form a team Sensor-based: search, estimate, identify, track, map (i) complexity analysis (time/energy)

(ii) models/algorithms for asynchronous networks with agent arrival/departures

(iii) parallel, sequential, hierarchical composition of behaviors

Part II: Motion Coordination

Scenarios examples of networks, tasks, ctrl+comm laws

- (i) rendezvous
- (ii) deployment

Rendezvous

- (i) H. Ando, Y. Oasa, I. Suzuki, and M. Yamashita. Distributed memoryless point convergence algorithm for mobile robots with limited visibility. *IEEE Transactions on Robotics and Automation*, 15(5):818–828, 1999
- (ii) J. Lin, A. S. Morse, and B. D. O. Anderson. The multi-agent rendezvous problem. In *IEEE Conf. on Decision and Control*, pages 1508–1513, Maui, HI, December 2003
- (iii) J. Cortés, S. Martínez, and F. Bullo. Robust rendezvous for mobile autonomous agents via proximity graphs in arbitrary dimensions. *IEEE Transactions on Automatic Control*, 51(6), 2006. To appear

Deployment

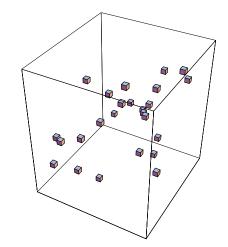
- (i) J. Cortés, S. Martínez, T. Karatas, and F. Bullo. Coverage control for mobile sensing networks. *IEEE Transactions* on *Robotics and Automation*, 20(2):243–255, 2004
- J. Cortés, S. Martínez, and F. Bullo. Spatially-distributed coverage optimization and control with limited-range interactions. ESAIM. Control, Optimisation & Calculus of Variations, 11:691–719, 2005

Scenario 1: aggregation laws for rendezous

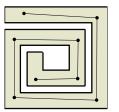
Aggregation laws

At each comm round:

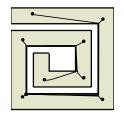
- 1: acquire neighbors' positions
- 2: compute connectivity constraint set
- 3: move towards circumcenter of neighbors (while remaining connected)



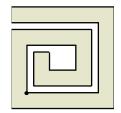
Initial position of the agents



Evolution of the network



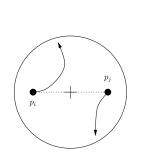
Final position of the agents

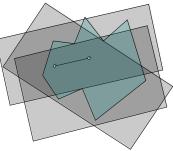


Task: rendezvous with connectivity constraint

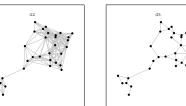
Scenario 1: aggregation laws for rendezous, cont'd

Pair-wise motion constraint set for connectivity maintenance

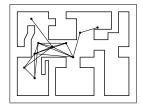


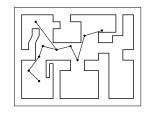


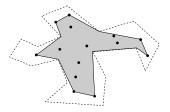
Reducing number of constraints



Lyapunov function







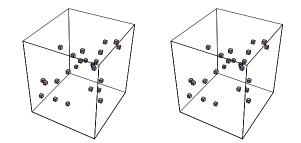
Scenario 1: Example complexity analysis

(i) first-order agents with disk graph, for d = 1,

$$\Gamma C(\mathcal{T}_{\text{rendezvous}}, \mathcal{CC}_{\text{circumcenter}}) \in \Theta(N)$$

(ii) first-order agents with limited Delaunay, for d = 1,

$$\mathrm{TC}(\mathcal{T}_{(r\epsilon)\text{-rendezvous}}, \mathcal{CC}_{\mathsf{circumcenter}}) \in \Theta(N^2 \log(N\epsilon^{-1}))$$



Tridiagonal Toeplitz and circulant systems

Let $N \geq 2$, $\epsilon \in]0, 1[$, and $a, b, c \in \mathbb{R}$. Let $x, y \colon \mathbb{N}_0 \to \mathbb{R}^N$ solve:

$$\begin{aligned} x(\ell+1) &= \operatorname{Trid}_N(a,b,c) \, x(\ell), & x(0) = x_0, \\ y(\ell+1) &= \operatorname{Circ}_N(a,b,c) \, y(\ell), & y(0) = y_0. \end{aligned}$$

(i) if $a = c \neq 0$ and |b| + 2|a| = 1, then $\lim_{\ell \to +\infty} x(\ell) = 0$, and the maximum time required for $||x(\ell)||_2 \le \epsilon ||x_0||_2$ is $\Theta(N^2 \log \epsilon^{-1})$;

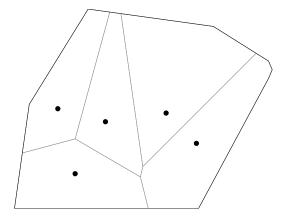
- (ii) if $a \neq 0$, c = 0 and 0 < |b| < 1, then $\lim_{\ell \to +\infty} x(\ell) = 0$, and the maximum time required for $||x(\ell)||_2 \le \epsilon ||x_0||_2$ is $O(N \log N + \log \epsilon^{-1})$;
- (iii) if $a \ge 0$, $c \ge 0$, b > 0, and a + b + c = 1, then $\lim_{\ell \to +\infty} y(\ell) = y_{\mathsf{ave}} \mathbf{1}$, where $y_{\mathsf{ave}} = \frac{1}{N} \mathbf{1}^T y_0$, and the maximum time required for $\|y(\ell) - y_{\mathsf{ave}} \mathbf{1}\|_2 \le \epsilon \|y_0 - y_{\mathsf{ave}} \mathbf{1}\|_2$ is $\Theta(N^2 \log \epsilon^{-1})$.

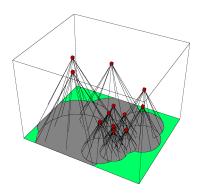
Scenario 2: dispersion laws for deployment

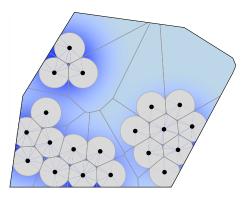
Dispersion laws

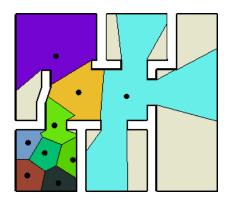
At each comm round:

- 1: acquire neighbors' positions
- 2: compute own dominance region
- 3: move towards incenter / circumcenter / centroid of own dominance region









Scenarios: optimal deployment

ANALYSIS of cooperative distributed behaviors

(i) how do animals share territory?what if every fish in a swarm goestoward center of own dominance region?



Barlow, Hexagonal territories. Anim. Behav. '74

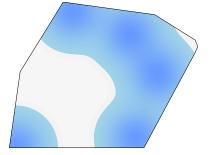
(ii) what if each vehicle moves toward center of mass of own Voronoi cell?(iii) what if each vehicle moves away from closest vehicle?

DESIGN of performance metric

(iv) how to cover a region with n minimum radius overlapping disks?
(v) how to design a minimum-distorsion (fixed-rate) vector quantizer? (Lloyd '57)
(vi) where to place mailboxes in a city / cache servers on the internet?

Scenario 2: general multi-center function

Objective: Given agents (p_1, \ldots, p_n) in convex environment Q unspecified comm graph, achieve optimal coverage



Expected environment coverage

- let ϕ be distribution density function
- let *f* be a performance/penalty function

$$f(\|q-p_i|)$$
 is price for p_i to service q

define multi-center function

$$\mathcal{H}_{\mathsf{C}}(p_1,\ldots,p_n) = E_{\phi}\left[\min_i f(\|q-p_i\|)\right]$$

Scenario 2: distributed gradient result

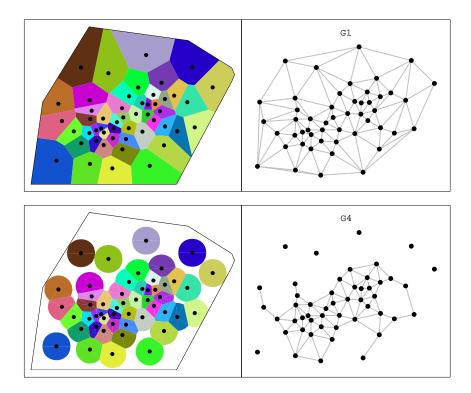
For a general non-decreasing $f : \overline{\mathbb{R}}_+ \to \mathbb{R}$ piecewise differentiable with finite-jump discontinuities at $R_1 < \cdots < R_m$

Thm:

$$\begin{split} \frac{\partial \mathcal{H}_{\mathsf{C}}}{\partial p_{i}}(p_{1},\ldots,p_{n}) &= \int_{V_{i}} \frac{\partial}{\partial p_{i}} f(\|q-p_{i}\|) \phi(q) dq \\ &+ \sum_{\alpha=1}^{m} \Delta f_{\alpha}(R_{\alpha}) \Big(\sum_{k=1}^{M_{i}(2R_{\alpha})} \int_{\operatorname{arc}_{i,k}(2R_{\alpha})} n_{B_{R_{\alpha}}(p_{i})} d\phi \Big) \\ &= \operatorname{integral over} V_{i} + \operatorname{integral along arcs inside} V_{i} \end{split}$$

Gradient depends on information contained in V_i

On Voronoi and limited-Voronoi partitions



 $\frac{\partial \mathcal{H}_{C}}{\partial p_{i}}$ is distributed over Delaunay graph, but not disk graph

Scenario 2: truncation

problem $\partial \mathcal{H}_C$ distributed over Delaunay graph, but comm. is disk graph

approach truncate $f_{\frac{r}{2}}(x) = f(x) \ 1_{[0,\frac{r}{2})}(x) + (\sup_Q f) \cdot 1_{[\frac{r}{2},+\infty)}(x)$,

$$\mathcal{H}_{\frac{r}{2}}(p_1,\ldots,p_n) = E_{\phi}\left[\min_i f_{\frac{r}{2}}(\|q-p_i\|)\right]$$

Result 1: \mathcal{H}_{C} constant-factor approximation

$$\beta \mathcal{H}_{\frac{r}{2}}(P) \leq \mathcal{H}_{\mathsf{C}}(P) \leq \mathcal{H}_{\frac{r}{2}}(P), \quad \beta = \left(\frac{r}{2\operatorname{diam}(Q)}\right)^2$$

Result 2 Gradient of $\mathcal{H}_{\frac{r}{2}}$ is distributed over limited-range Delaunay

$$\frac{\partial \mathcal{H}_{\frac{r}{2}}}{\partial p_{i}} = 2M_{V_{i}(P)\cap B_{\frac{r}{2}}(p_{i})}(C_{V_{i}(P)\cap B_{\frac{r}{2}}(p_{i})}-p_{i}) - \left(\left(\frac{r}{2}\right)^{2} - \operatorname{diam}(Q)^{2}\right)\sum_{k=1}^{\mathcal{M}_{i}(r)}\int_{\operatorname{arc}_{i,k}(r)} n_{B_{\frac{r}{2}}(p_{i})}\phi$$

Aggregate objective functions

design of aggregate network-wide cost/objective/utility functions

- objective functions to encode motion coordination objective
- objective functions as Lyapunov functions
- objective functions for gradient flows

	\mathcal{H}_{C}	\mathcal{H}_{area}	\mathcal{H}_{diam}
DEFINITION	$E\left[\min d(q,p_i) ight]$	${\sf area}_\phi(\cup_i B_{r/2}(p_i))$	$\max_{i,j} \ p_i - p_j\ $
SMOOTHNESS	C^1	globally Lipschitz	continuous, locally Lipschitz
CRITICAL	Centroidal	r-limited	common
POINTS	Voronoi con-	Voronoi	location
MINIMA	figurations	configurations*	for p_i
HEURISTIC	expected distortion	area covered	diameter connected component
DESCRIPTION			

(i) general pattern formation problem
(ii) static and dynamic motion patterns
(iii) algorithms for line-of-sight 3D networks
(iv) connectivity and collision avoidance algorithms

Problems of interest

- optimal sensor placement
- localization, estimation
- distributed sensing tasks:

search, exploration, map building, target identification

References

Target Servicing

- (i) R. W. Beard, T. W. McLain, M. A. Goodrich, and E. P. Anderson. Coordinated target assignment and intercept for unmanned air vehicles. *IEEE Transactions on Robotics and Automation*, 18(6):911–922, 2002
- (ii) A. E. Gil, K. M. Passino, and A. Sparks. Cooperative scheduling of tasks for networked uninhabted autonomous vehicles. In *IEEE Conf. on Decision and Control*, pages 522–527, Maui, Hawaii, December 2003
- (iii) W. Li and C. G. Cassandras. Stability properties of a cooperative receding horizon controller. In *IEEE Conf. on Decision and Control*, pages 492–497, Maui, HI, December 2003
- (iv) E. Frazzoli and F. Bullo. Decentralized algorithms for vehicle routing in a stochastic time-varying environment. In *IEEE Conf. on Decision and Control*, pages 3357–3363, Paradise Island, Bahamas, December 2004

Boundary Estimation

- (i) D. Marthaler and A. L. Bertozzi. Tracking environmental level sets with autonomous vehicles. In *Proc. of the Conference on Cooperative Control and Optimization*, Gainesville, FL, December 2002
- (ii) J. Clark and R. Fierro. Cooperative hybrid control of robotic sensors for perimeter detection and tracking. In American Control Conference, pages 3500–3505, Portland, OR, June 2005
- (iii) D. W. Casbeer, S.-M. Li, R. W. Beard, R. K. Mehra, and T. W. McLain. Forest fire monitoring with multiple small UAVs. In *American Control Conference*, pages 3530–3535, Portland, OR, June 2005
- (iv) F. Zhang and N. E. Leonard. Generating contour plots using multiple sensor platforms. In *IEEE Swarm Intelligence Symposium*, pages 309–316, Pasadena, CA, June 2005

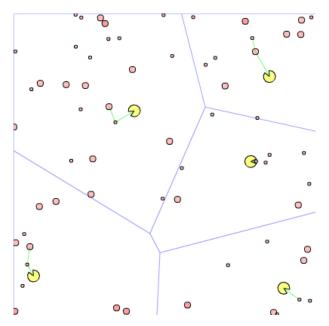
Scenario 3: Vehicle Routing

Objective: Given agents (p_1, \ldots, p_n) moving in environment Q service targets in environment

Model:

- targets arise randomly in space/time
- vehicle know of targets arrivals

Scenario 3 — min expected waiting time



Scenario 3: receding-horizon TSP algorithm, I

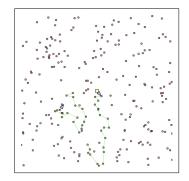
Name: (Single Vehicle) Receding-horizon TSP

For $\eta \in (0, 1]$, single agent performs:

- 1: while no targets, dispersion/coverage algorithm
- 2: while targets waiting,

(i) compute optimal TSP tour through all targets

(ii) service the η -fraction of tour with maximal number of targets



Asymptotically optimal in light and high traffic

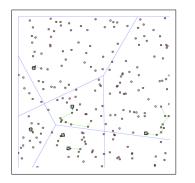
Scenario 3: receding-horizon TSP algorithm, II

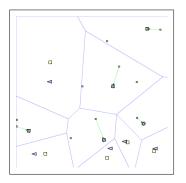
Name: Receding-horizon TSP

For $\eta \in (0, 1]$, agent *i* performs:

- 1: compute own Voronoi cell V_i
- 2: apply Single-Vehicle RH-TSP policy on V_i

Asymptotically optimal in light and high traffic (simulations only)

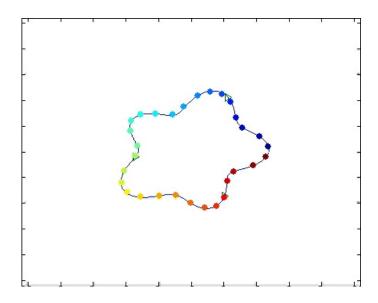




Scenario 4: Boundary Estimation

Objective: estimate/interpolate moving boundary **Model:**

- UAVs can locally sense and track
- comm. graph is ring
- interpolation via waypoints



(i) network modeling

network, ctrl+comm algorithm, task, complexity

coordination algorithm

optimal deployment, rendezvous, vehicle routing scalable, adaptive, asynchronous, agent arrival/departure

(ii) Systematic algorithm design

- meaningful aggregate cost functions
- class of (gradient) algorithms local, distributed
- geometric graphs
- stability theory for networked hybrid systems