

Efficient quantum algorithms for nonlinear ODEs and PDEs

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Reference

- **Jin-Peng Liu**, Herman Øie Kolden, Hari K. Krovi, Nuno F. Loureiro, Konstantina Trivisa and Andrew M. Childs, *Efficient quantum algorithm for dissipative nonlinear differential equations*, Proceedings of the National Academy of Sciences 118, 35 (2021).

(Presented by Andrew Childs on Monday)

- Dong An, Di Fang, Stephen Jordan, **Jin-Peng Liu**, Guang Hao Low and Jiasu Wang, *Efficient quantum algorithm for nonlinear reaction-diffusion equations and energy estimation*, in preparation.

Quantum Numerical Linear Algebra

Quantum computers can offer potential **exponential** speedup for producing the quantum-encoding solutions of

- quantum simulations[†];
- linear systems[‡];
- linear differential equations[§].

[†][Lloyd 96; Berry et al. 15; Low, Chuang 17]

[‡][Harrow, Hassidim, Lloyd 08; Ambainis 12; Childs, Kothari, Somma 15]

[§][Berry 14; Berry et al. 17; Childs, Liu 19]

Quantum Scientific Computation

For real-world problems in aeronautics, climate, epidemic, finance... a challenge is the difference between the linearity of quantum mechanics and **nonlinearities** in complex systems.

Questions

- Can we provide exponential speedup for nonlinear dynamics?
- Can we postprocess the quantum states to provide meaningful classical information?

Quantum Algorithm for Nonlinear ODEs

Fundamental obstacles

- Direct simulation: represent nonlinearities by multiple copies of quantum states. Need to maintain overall $2^{O(T)}$ copies[†].
- Quantum mechanics with strong nonlinearities can imply poly-time solution for NP-complete and $\#P$ problems[‡].

Solutions[§]

- Carleman linearization for dissipative ODEs with $R < 1$, where R is a ℓ_2 ratio of nonlinearities and linear dissipation.
- When $R \geq \sqrt{2}$, there is a worst case such that any quantum algorithm must suffer from $\exp(T)$.

[†][Leyton, Osborne 08]

[‡][Abrams, Lloyd 98; Aaronson 05; Childs, Young 16]

[§][Liu et al. 21]

Quantum Scientific Computation

Questions

- Can we provide exponential speedup for **stronger** nonlinear dynamics?
- Can we postprocess the quantum states to provide meaningful classical information?

Inspirations

- Carleman linearization practically works for well-behaved solutions even $R > 1$, e.g. laminar flows[†].
- Analog of energy estimation in quantum physics/chemistry.

[†][Liu et al. 21]

Reaction-diffusion Equations

$$\frac{\partial u}{\partial t}(\mathbf{x}, t) = D\Delta u(\mathbf{x}, t) + f(u(\mathbf{x}, t)), \quad \mathbf{x} \in [0, 1]^d. \quad (1)$$

- $f(u) = u - u^2$: biology, ecology, social networks.
- $f(u) = u - u^3$: phase separation, data/image processing.
- $f(u) = \sum_k a_k u^k$: disorder systems, credit valuation adjustment.

$u(\mathbf{x}, t)$ is the L^2 gradient flow of minimizing the energy functional

$$E(u) = D \int |\nabla u|^2 dx + \int F(u) dx, \quad (2)$$

where $\frac{\partial F}{\partial u}(u) = f(u)$. E.g. $f(u) = u - u^3$, $F(u) = \frac{1}{4}(1 - u^2)^2$.

Reaction-diffusion Equations

Spatial discretization

$$\frac{dU_i}{dt} = D \sum_{k=1}^{n^d} (\Delta_h)_{ij} U_k + f(U_i), \quad i \in [n^d], \quad (3)$$

$U_i(t) = u(\mathbf{x}_i, t)$ on grids, Δ_h : Laplacian matrix with Dirichlet/mixed BCs. Assume polynomial f has roots $\gamma_1 < \dots < \gamma_M$, $M \geq 2$.

Proposition

Comparison principle. *If $U_i(0)$ lies in $[\gamma_k, \gamma_{k+1}]$, then $\forall t$, $U_i(t)$ always stays in $[\gamma_k, \gamma_{k+1}]$.*

Maximum principle. *If $U_i(0)$ lies in $[\gamma_1, \gamma_M]$, then $\forall t$, $U_i(t)$ always stays in $[\gamma_1, \gamma_M]$. It indicates that $\|U(t)\|_\infty \leq \gamma := \max\{|\gamma_1|, |\gamma_M|\}$.*

Carleman Linearization

Consider $\frac{dU_i}{dt} = \frac{1}{h^2}(U_{i-1} - 2U_i + U_{i+1}) + U_i - U_i^2$.

Embedding and truncation

- $\frac{dU_i}{dt} = \frac{1}{h^2}(U_{i-1} - 2U_i + U_{i+1}) + U_i - U_i^2$.
- $\frac{dU_i^2}{dt} = 2U_i \frac{dU_i}{dt} = \frac{2}{h^2}(U_i U_{i-1} - 2U_i^2 + U_i U_{i+1}^2) + 2U_i^2 - 2U_i^3$.
-
- $\frac{dU_i^N}{dt} \approx \frac{N}{h^2}(U_i^{N-1} U_{i-1} - 2U_i^N + U_i^{N-1} U_{i+1}) + N U_i^N$.
- Give a linear ODE with variables $y_j \approx U^{\otimes j} \in \mathbb{R}^{n^{jd}}$ for $k \in [N]$.

A system of n^d -dim **nonlinear** ODEs is embedded to a system of **linear** ODEs with truncation order N , with dimension $n^d + n^{2d} + \dots + n^{Nd}$.

Carleman Linearization

We give a linear ODEs $\frac{d\hat{y}}{dt} = A\hat{y}$ with $\hat{y}(0) = \hat{y}_{\text{in}}$, by

$$\frac{d}{dt} \begin{pmatrix} \hat{y}_1 \\ \hat{y}_2 \\ \vdots \\ \hat{y}_N \end{pmatrix} = \begin{pmatrix} A_1^1 & \cdots & A_M^1 & & & \\ A_1^2 & A_2^2 & & \ddots & & \\ & \ddots & \ddots & & & \\ & & \ddots & \ddots & & \\ & & & A_{N-1}^N & & \\ & & & & A_N^N & \end{pmatrix} \begin{pmatrix} \hat{y}_1 \\ \hat{y}_2 \\ \vdots \\ \hat{y}_N \end{pmatrix}, \quad (4)$$

where $\hat{y}_j \approx U^{\otimes j} \in \mathbb{R}^{n^{jd}}$, $\hat{y}_{\text{in}} = [U_{\text{in}}; U_{\text{in}}^{\otimes 2}; \dots; U_{\text{in}}^{\otimes N}]$, and A_{j+k-1}^j encodes k -th order polynomial.

We denote the error from the truncation as $\eta_j(t) := U^{\otimes j}(t) - \hat{y}_j(t)$.

Carleman Linearization

$$\frac{dU_i}{dt} = D \sum_{k=1}^{n^d} (\Delta_h)_{ij} U_k + a_1 U_i + a_2 U_i^2, \quad i \in [n^d]. \quad (5)$$

Assume the largest eigenvalue of $\Delta_h + a_1 I$: $\lambda_1 < 0$, and U_{in} satisfies Maximum Principle, then $\forall t, \|U(t)\|_\infty \leq \gamma$.

Lemma (convergence analysis)

(i) Assume $R = \frac{|a_2|}{|\lambda_1|} \|U_{\text{in}}\| < 1$. The ℓ_2 error bound satisfies[†]

$$\|\eta_j(t)\| \leq \|U_{\text{in}}\|^j R^{N-1}. \quad (6)$$

(ii) Assume $R_D = \frac{|a_2|}{|\lambda_1|} \gamma C < 1$, where $C = O(d)$ and independent of n . The ℓ_∞ error bound satisfies

$$\|\eta_j(t)\|_\infty \leq \gamma^j R_D^{N-1}. \quad (7)$$

[†][Liu et al. 21]

Carleman Linearization

$$\frac{dU_i}{dt} = D \sum_{k=1}^{n^d} (\Delta_h)_{ij} U_k + a_1 U_i + g(U_i), \quad i \in [n^d]. \quad (8)$$

$g(u) = a_0 + \sum_{k=2}^M a_k u^k$. Given $\lambda_1 < 0$, and $\forall t, \|U(t)\|_\infty \leq \gamma$.

Lemma (convergence analysis)

(i) Assume $R = \frac{1}{|\lambda_1| \|U_{\text{in}}\|} g(\|U_{\text{in}}\|) < 1$. The ℓ_2 error bound satisfies

$$\|\eta_j(t)\| \leq \|U_{\text{in}}\|^j R^{\frac{N-1}{M-1}}. \quad (9)$$

(ii) Assume $a_0 = 0$, $R_D = \frac{C}{|\lambda_1| \gamma} g(\gamma) < 1$, where $C = O(d)$ and independent of n . The ℓ_∞ error bound satisfies

$$\|\eta_j(t)\|_\infty \leq \gamma^j R_D^{\frac{N-1}{M-1}}. \quad (10)$$

Carleman Linearization

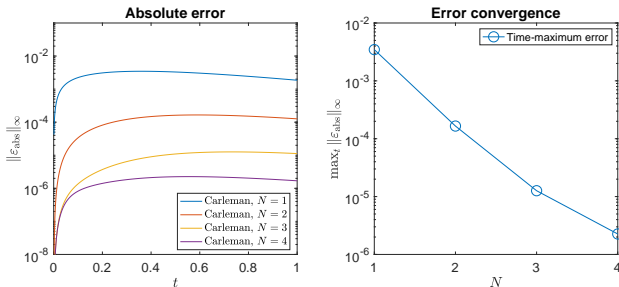


Figure: $\frac{\partial u}{\partial t} = 0.2\Delta u + 0.2u - u^2$, $u(x, 0) = 0.1(1 - \cos(2\pi x))$ with homogenous Dirichlet boundary condition. Spatial grid number $n = 16$. $C = 4.4620$, $R_D = 0.5047$. Left: l_∞ norm of the absolute error between the Carleman solutions. Right: convergence of the time-maximum error.

Quantum Carleman Linearization

We linearize the equation with the N -th order truncation, which we solve using the Euler method with step h and QLSA.

Theorem (quantum algorithm)

Assume $R_D < 1$. Let $q := \frac{\|U_{\text{in}}\|}{\|U(T)\|}$. There is a quantum algorithm that

(i) produces $|U\rangle \sim \sum_k U(kh)|k\rangle$ with $\tilde{O}(T^2 d^2 \|U_{\text{in}}\|^{2N} / \epsilon)$;

(ii) produces $|U(T)\rangle$ with $\tilde{O}(T^2 d^2 q \|U_{\text{in}}\|^{2N} / \epsilon)$.

Corollary

Assume $R < 1$ (or assume $R \leq 1$ with $R_D < 1$). The above results can be reduced to $\tilde{O}(T^2 d^2 q / \epsilon)$ and $\tilde{O}(T^2 d^2 / \epsilon)$.

Quantum Carleman Linearization

Why better convergence

- If there is a blow-up solution, the truncation error is generally unbounded. The exponentially increasing error is also used to show the worst-case complexity exponential in time[†].
- Maximum Principle: $\forall t, \|U(t)\|_{\infty} \leq \gamma$.

Comparison

- $R_D \ll R$ if $\gamma \ll \|U_{\text{in}}\|$ for large n and d .
- Assume $R_D < 1$, error decays in R_D^N and cost scales in $\|U_{\text{in}}\|^{2N}$: a trade-off between the approximation and the cost; while there is no dependence of N when $R < 1$.

[†][Liu et al. 21]

Applications

Postprocess the history or final state

We have developed efficient quantum algorithm for producing

$$|U\rangle = \frac{1}{Z_0} \sum_{k \in [m], \mathbf{l} \in [n^d]} u(\mathbf{x}_{\mathbf{l}}, t_k) |l_1\rangle \dots |l_d\rangle |k\rangle. \quad (11)$$

Postselecting $|U(T)\rangle$ ($T = mh$) relies on $q = \frac{\|U_{\text{in}}\|}{\|U(T)\|}$.

- For homogeneous systems with $R < 1$, e.g. $f(u) = u - u^3$, the solution decays exponentially in time, i.e. $q = \exp(T)$.
- Given external forces or inhomogeneous BCs, the solution can remain nonzero, decay slowly, or be oscillatory i.e. $q = \text{poly}(T)$.

Applications

Ratio of mean square amplitude

$$\frac{\int_{\Omega_t} \int_{\Omega_x} |u(\mathbf{x}, t)|^2 dt d\mathbf{x}}{\int_0^T \int_{[0,1]^d} |u(\mathbf{x}, t)|^2 dt d\mathbf{x}} \sim \frac{\sum_{k \in I_t, \mathbf{l} \in I_x} |u(\mathbf{x}_l, t_k)|^2}{\sum_{k \in [m], \mathbf{l} \in [n^d]} |u(\mathbf{x}_l, t_k)|^2}. \quad (12)$$

Let the projector P associate with indices $I_t \subset [m]$ and $I_x \subset [n^d]$:
 $P = \sum_{k \in I_t, \mathbf{l} \in I_x} (|l_1\rangle\langle l_1|) \otimes \cdots \otimes (|l_d\rangle\langle l_d|) \otimes (|k\rangle\langle k|).$

Amplitude estimation: perform $I - 2P$ with $\tilde{O}(1/\epsilon)$ to estimate $\langle U|P|U\rangle$.

Example: diffusive Lotka-Volterra equations

$$\begin{aligned} \frac{\partial u}{\partial t} &= D\Delta u + \alpha u - \beta uv \\ \frac{\partial v}{\partial t} &= D\Delta v + \delta uv - \gamma v. \end{aligned} \quad (13)$$

Traveling waves in predator-prey, economic cycle, and disease models.

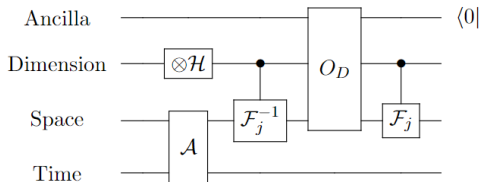
Applications

Ratio of diffusive energy

We can produce

$$|\nabla_{\mathbf{x}} U\rangle = \frac{1}{Z_1} \sum_{k \in [m], \mathbf{l} \in [n^d]} \nabla_{x_j} u(\mathbf{x}_{\mathbf{l}}, t_k) |j\rangle |l_1\rangle \dots |l_d\rangle |k\rangle \quad (14)$$

by performing QFT/IQFT on $|U\rangle$ with cost $\tilde{O}(d)$.



Quantum circuit for preparing a quantum state encoding partial derivatives

Applications

Ratio of diffusive energy

$$\frac{\int_{\Omega_t} \int_{\Omega_x} |\nabla_{\mathbf{x}} u(\mathbf{x}, t)|^2 dt d\mathbf{x}}{\int_0^T \int_{[0,1]^d} |\nabla_{\mathbf{x}} u(\mathbf{x}, t)|^2 dt d\mathbf{x}} \sim \frac{\sum_{k \in I_t, \mathbf{l} \in I_x} |\nabla_{\mathbf{x}} u(\mathbf{x}_{\mathbf{l}}, t_k)|^2}{\sum_{k \in [m], \mathbf{l} \in [n^d]} |\nabla_{\mathbf{x}} u(\mathbf{x}_{\mathbf{l}}, t_k)|^2}. \quad (15)$$

Perform $I - 2P$ with $\tilde{O}(1/\epsilon)$ to estimate $\langle \nabla_{\mathbf{x}} U | P | \nabla_{\mathbf{x}} U \rangle$.

Example: Allen-Cahn equation

- $\frac{\partial u}{\partial t} = D\Delta u + u - u^3 + F(t)$: phase separation and transition.
- $\frac{\partial u}{\partial t} = D\Delta u + u - u^3 + F(u - u_0)$: data classification, graph cuts, signal or image denoising and reconstruction.

u is the L^2 gradient flow of minimizing a regularized energy functional.

$\|\nabla_{\mathbf{x}} u\|^2$ measures the L^2 total variation distance.

Summary

Takeaways

- Quantum computer can efficiently characterize weak gradient flows in $\tilde{O}(T^2 d^2 / \epsilon)$ when $R_D < 1$ or $R < 1$. l_∞ aprior estimate is used to improve the linearization and rule out the worst case.
- Nonlinear ODEs/PDEs exhibit rich phenomena. Ratio of energy proportion to amplitude squared and total variation can be estimated in $\tilde{O}(d/\epsilon)$.

Outlook

Quantum algorithm

- Advection term $f(u, \nabla u)$: N-S equation, Boltzmann equation.
- Non-dissipative systems: nonlinear Schrödinger equations.
- Discrete gradient flows for optimization and control.
- Hermitian/skew-Hermitian linearization.

Postprocessing

- Scattering cross section.
- Time frequency analysis.
- Free energy estimation: $D \int |\nabla u|^2 dx + \int F(u) dx$.