

# Efficiency of Quasi-Newton methods on strictly positive functions

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October 12, 2010, UCLA

IPAM Workshop II: Numerical methods for Continuous  
Optimization

# Outline

- 1 Problem classes for absolute and relative accuracy
- 2 Strictly positive functions
- 3 Quasi-Newton method
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# Standard approach: Absolute accuracy

**Problem:**  $f^* \stackrel{\text{def}}{=} \min_{x \in Q} f(x)$ , where  $Q \subseteq R^n$  is a closed convex set.

**Definition:**

For  $\epsilon > 0$ , find  $\bar{x} \in Q$  satisfying  $f(\bar{x}) \leq f^* + \epsilon$ .

**Problem classes**

**1** *Bounds on the growth.* (Strong) convexity with  $\mu \geq 0$ :  
$$f(y) \geq f(x) + \langle f'(x), y - x \rangle + \frac{1}{2}\mu\|y - x\|^2, \quad x, y \in Q.$$

**2** *Bounds on derivatives.* For example,  
$$\|f'(x)\|_* \leq M, \quad \text{or,} \quad \|f''(x)\| \leq L, \quad \text{etc.}$$

**Important:** operation  $f \Rightarrow f + \text{const}$  does not change complexity.

# Relative accuracy (RA)

**Problem:**  $f^* \stackrel{\text{def}}{=} \min_{x \in Q} f(x) > 0$ , where  $Q$  is a closed convex set.

**Definition:**

For  $\delta \in (0, 1)$ , find  $\bar{x} \in Q$  satisfying  $(1 - \delta)f(\bar{x}) \leq f^* \leq f(\bar{x})$ .

Condition  $f^* > 0$  must be guaranteed. *How?*

1. Homogeneous model [N.08]: Let  $0 \notin Q$ . For  $f(x) = \max_{s \in B} \langle s, x \rangle$  with  $0 \in \text{int } B$  define  $\gamma_0, \gamma_1$ :  $\gamma_0 \|x\| \leq f(x) \leq \gamma_1 \|x\|$ ,  $\alpha = \frac{\gamma_0}{\gamma_1}$ . Then, by smoothing technique, we get complexity  $O^*\left(\frac{1}{\alpha\delta}\right)$ .

2. Polyhedral model [N.09]: if  $B = \text{Conv}(\pm a_i, i = 1 \dots m) \subset R^n$ , then we need  $O^*\left(\frac{n^{1/2}}{\delta}\right)$  iterations. (An appropriate norm is constructed by preprocessing.)

**Question:** Can we address RA in Black-Box Framework?  $\Rightarrow$  Need new problem classes. (Invariant with respect to multiplication.)

# Barrier subgradient method [N.10]

**Problem:**  $\max_{x \in Q} \phi(x),$

- $Q$  is a closed convex set endowed with a  $\nu$ -self-concordant barrier  $F(x)$ .
- $\phi$  is a concave function, which is *non-negative* on  $Q$ .

**Method:** (potential  $f = \ln \phi$ )

$$x_{k+1} = \arg \max_{x \in Q} \left[ \sum_{i=0}^k \left\langle \frac{\phi'(x_i)}{\phi(x_i)}, x - x_i \right\rangle - \left( 1 + \sqrt{\frac{k+1}{\nu}} \right) F(x) \right].$$

**Convergence:**  $\max_{0 \leq i \leq k} \phi(x_i) \geq \phi^* \left( 1 - O^* \left( \sqrt{\frac{\nu}{k+1}} + \frac{\nu}{k+1} \right) \right).$

**Complexity:**  $O^* \left( \frac{\nu}{\delta^2} \right)$  iterations.

# Strictly positive functions

## Definition

Convex function  $f$  is called strictly positive on  $Q$  if

$$f(y) + f(x) + \langle f'(x), y - x \rangle \geq 0, \quad x, y \in Q.$$

**Corollary:**  $f(y) \geq |f(x) + \langle f'(x), y - x \rangle|, \quad x, y \in Q.$

## Simple properties

- $f(x) \equiv \text{const} > 0$  is strictly positive.
- Strict positivity is an *affine-invariant* property.
- Class of strictly positive functions is a convex cone.

# Simple examples

Lemma 1. Let  $B$  be bounded, closed, and centrally symmetric.

Then  $f(x) = \max_{x \in B} \langle s, x \rangle$  is strictly positive on  $R^n$ .

**Proof:** Since  $f(x) = \langle f'(x), x \rangle$  and  $-f'(x) \in B$ , we have

$$f(y) \geq \langle -f'(x), y \rangle = -f(x) - \langle f'(x), y - x \rangle. \quad \square$$

The simplest examples of strictly positive functions are *norms*.

Lemma 2. Let  $f_1(x)$  and  $f_2(x)$  be strictly positive on  $Q$ .

Then  $f(x) = \max\{f_1(x), f_2(x)\}$  is also strictly positive.

**Proof:** For arbitrary  $x \in Q$ , assume  $f_1(x) \geq f_2(x)$ . Then,

$$\begin{aligned} f(y) &\geq f_1(y) \geq -f_1(x) - \langle f_1'(x), y - x \rangle \\ &= -f(x) - \langle f'(x), y - x \rangle. \quad \square \end{aligned}$$

# Particular examples

All functions below are strictly positive:

$$f(x) = \max_{1 \leq i \leq m} \|A_i x - b_i\|,$$

$$f(x) = \sum_{i=1}^m \|A_i x - b_i\|,$$

$$f(x) = \sigma_{\max} \left( \sum_{i=1}^n A_i x^{(i)} \right),$$

$$f(x) = \sum_{j=1}^m \sigma_j \left( \sum_{i=1}^n A_i x^{(i)} \right),$$

where  $A_i \in R^{m \times n}$ , and  $b_i \in R^m$ ,  $i = 1 \dots n$ .



# General convex functions

Theorem 1. Let  $\phi$  be convex function on  $Q$  with uniformly bounded subgradients:  $\|\phi'(x)\|^* \leq L, \quad x \in Q.$

Then  $f(x) = \max\{\phi(x), L\|x\|\}$  is strictly positive on  $Q$ .

**Proof:** Clearly,  $\|f'(x)\|^* \leq L$ . Therefore,

$$\begin{aligned} f(y) + f(x) + \langle f'(x), y - x \rangle &\geq L\|y\| + L\|x\| + \langle f'(x), y - x \rangle \\ &\geq L\|y\| + L\|x\| - L\|y - x\| \geq 0. \end{aligned}$$



# Shifted general optimization problem

Consider the problem:  $\min_{x \in Q} \phi(x)$ , where  $\phi$  has bounded subgradients. Let  $x^* \in Q$  be its optimal solution.

Lemma 3. For  $x_0 \in Q$  define

$$f(x) = \max\{\phi(x) - \phi(x_0) + 2LR, L\|x - x_0\|\}.$$

It is strictly positive. If  $\|x - x_0\| \leq R$  then  $f(x) \equiv \phi(x) + \text{const}$ .

If  $\|x_0 - x^*\| \leq R$ , then the optimal value  $f^*$  of the equivalent problem  $\min_{x \in Q} f(x)$  satisfies  $LR \leq f^* \leq 2LR$ .

**Proof:** If  $\|x - x_0\| \leq R$ , then

$$\phi(x) - \phi(x_0) + 2LR \geq 2LR - L\|x - x_0\| \geq L\|x - x_0\|.$$

Further,  $f^* \leq f(x_0) = 2LR$ , and

$$f(x) \geq \max\{2LR - L\|x - x_0\|, L\|x - x_0\|\} \geq LR. \quad \square$$

# Optimization problem with squared objective

**Problem:**  $\min_{x \in Q} f(x)$ , where  $f$  is strictly positive on  $Q$ .

**New objective:**  $\hat{f}(x) = \frac{1}{2}f^2(x)$ ,  $\hat{f}'(x) = f(x) \cdot f'(x)$ .

**Equivalent problem:**  $\min_{x \in Q} \hat{f}(x)$ .

Lemma 4. Let  $f$  be strictly positive on  $Q$ . Then for  $x, y \in Q$

$$\hat{f}(y) \geq \hat{f}(x) + \langle \hat{f}'(x), y - x \rangle + \frac{1}{2} \langle f'(x), y - x \rangle^2.$$

**Proof:** Indeed,

$$\begin{aligned} \hat{f}(y) &= \frac{1}{2}f^2(y) \geq \frac{1}{2}[f(x) + \langle f'(x), y - x \rangle]^2 \\ &= \hat{f}(x) + \langle \hat{f}'(x), y - x \rangle + \frac{1}{2} \langle f'(x), y - x \rangle^2. \quad \square \end{aligned}$$

**Important:** We have *nonlinear support function!*

# Quasi-Newton Method

Let us fix  $G_0 \succ 0$ , starting point  $x_0 \in Q$ , and accuracy  $\delta \in (0, 1)$ . Define  $\psi_0(x) = \frac{1}{2} \|x - x_0\|_{G_0}^2$ . For  $k \geq 0$ , consider the process:

$$x_k = \arg \min_{x \in Q} \psi_k(x),$$
$$\psi_{k+1}(x) = \psi_k(x) + a_k \left[ \hat{f}(x_k) + \langle \hat{f}'(x_k), x - x_k \rangle + \frac{1}{2} \langle f'(x_k), x - x_k \rangle^2 \right],$$

where

$$a_k = \frac{\delta}{1-\delta} \cdot \frac{1}{(\|f'(x_k)\|_{G_k}^*)^2}, \quad G_k = \psi_k''(x), \quad k \geq 0,$$

and  $\|h\|_G = \langle Gh, h \rangle^{1/2}$ ,  $\|g\|_G^* = \langle g, G^{-1}g \rangle^{1/2}$ .

Denote  $A_k = \sum_{i=0}^{k-1} a_i$ . Clearly,  $\psi_k(x) \leq A_k \hat{f}(x) + \psi_0(x)$ ,  $x \in Q$ .

We can use the technique of estimate sequences!

# Evolution of the Hessians

Since  $\psi_k(x)$  are quadratic, their Hessians  $G_k \succ 0$  are updated as

$$G_{k+1} = G_k + a_k \cdot f'(x_k) f'(x_k)^T = G_k + \frac{\delta}{1-\delta} \cdot \frac{f'(x_k) f'(x_k)^T}{(\|f'(x_k)\|_{G_k}^*)^2}, \quad k \geq 0.$$

Therefore,  $G_{k+1}^{-1} = G_k^{-1} - \delta \cdot \frac{G_k^{-1} f'(x_k) f'(x_k)^T G_k^{-1}}{(\|f'(x_k)\|_{G_k}^*)^2}$ .

**Important:**  $\det G_{k+1} = \frac{1}{1-\delta} \det G_k = \frac{1}{(1-\delta)^{k+1}} \det G_0$ .

Moreover,

$$\begin{aligned} \frac{1}{2} a_k^2 (\|\hat{f}'(x_k)\|_{G_{k+1}}^*)^2 &= a_k^2 \cdot \hat{f}(x_k) \cdot (\|f'(x_k)\|_{G_{k+1}}^*)^2 \\ &= a_k^2 \cdot \hat{f}(x_k) \cdot (1-\delta) \cdot (\|f'(x_k)\|_{G_k}^*)^2 \\ &= \delta \cdot a_k \cdot \hat{f}(x_k). \end{aligned}$$

# Main Lemma

For any  $k \geq 0$ ,

$$\psi_k^* \stackrel{\text{def}}{=} \min_{x \in Q} \psi_k(x) \geq (1 - \delta) \sum_{i=0}^{k-1} a_i \hat{f}(x_i).$$

**Proof:** Assume this is true for some  $k \geq 0$ . Since  $\psi_k(x)$  quadratic,

$$\psi_k(x) = \psi_k^* + \langle \psi'_k(x_k), x - x_k \rangle + \frac{1}{2} \|x - x_k\|_{G_k}^2 \geq \psi_k^* + \frac{1}{2} \|x - x_k\|_{G_k}^2.$$

Therefore,

$$\psi_{k+1}^* \geq \psi_k^* +$$

$$\min_{x \in Q} \left\{ \frac{1}{2} \|x - x_k\|_{G_k}^2 + a_k [\hat{f}(x_k) + \langle \hat{f}'(x_k), x - x_k \rangle + \frac{1}{2} \langle f'(x_k), x - x_k \rangle^2] \right\}$$

$$= \psi_k^* + a_k \hat{f}(x_k) + \min_{x \in Q} \left\{ \frac{1}{2} \|x - x_k\|_{G_{k+1}}^2 + a_k \langle \hat{f}'(x_k), x - x_k \rangle \right\}$$

$$\geq \psi_k^* + a_k \hat{f}(x_k) - \frac{1}{2} a_k^2 \|\hat{f}'(x_k)\|_{G_{k+1}}^2 = \psi_k^* + (1 - \delta) \cdot a_k \hat{f}(x_k). \quad \square$$

# Rate of convergence

Denote  $\tilde{x}_k = \frac{1}{A_k} \sum_{i=0}^{k-1} a_i x_i$ . Recall:  $G_{k+1} = G_k + a_k \cdot f'(x_k) f'(x_k)^T$ .

Theorem: Assume that for SP-function  $f$ ,  $\|f'(\cdot)\|_{G_0}^* \leq L$ .

Then,  $(1 - \delta) \hat{f}(\tilde{x}_k) \leq \hat{f}(x^*) + \frac{L^2 \|x_0 - x^*\|_{G_0}^2}{2n[e^{\delta(k+1)/n} - 1]}$ .

**Proof:** We have  $(1 - \delta) \hat{f}(x_k^*) \leq \hat{f}(x^*) + \frac{1}{2A_{k+1}} \|x_0 - x^*\|_{G_0}^2$ .

Let us estimate the growth of  $A_k$ . Denote  $\bar{G}_k = G_0^{-1/2} G_k G_0^{-1/2}$ .

$$\begin{aligned} A_k &= \sum_{i=0}^{k-1} a_i \geq \frac{1}{L^2} \sum_{i=0}^{k-1} a_i \|f'(x_i)\|_{G_0}^2 = \frac{1}{L^2} [\text{Trace } \bar{G}_k - n] \\ &\geq \frac{n}{L^2} \left[ \frac{1}{(1-\delta)^{k/n}} - 1 \right] \geq \frac{n}{L^2} [e^{\delta k/n} - 1]. \quad \square \end{aligned}$$

Definition: point  $\bar{x} \in Q$  is a solution with *mixed*  $(\epsilon, \delta)$ -accuracy if

$$(1 - \delta)\hat{f}(\bar{x}) \leq \hat{f}(x^*) + \epsilon.$$

- $\epsilon > 0$  serves as an absolute accuracy.
- $\delta \in (0, 1)$  represents the relative accuracy.

**Complexity:**  $N_n(\epsilon, \delta) \stackrel{\text{def}}{=} \frac{n}{\delta} \ln \left( 1 + \frac{L^2 R^2}{2n \cdot \epsilon} \right)$  iterations of Q-N scheme.

**Note:**

- High absolute accuracy is *easy* to achieve.
- High relative accuracy is *difficult*. (No need?)
- # of iterations is proportional to  $\frac{n}{\delta}$ . (Compare with BSM.)
- We have a uniform bound:  $N_n(\epsilon, \delta) < N_\infty(\epsilon, \delta) \stackrel{\text{def}}{=} \frac{L^2 R^2}{2\epsilon\delta}$ .



# Relative accuracy

**Our goal:** generate  $\bar{x} \in Q$  satisfying  $(1 - \delta)f(\bar{x}) \leq f^*$ .

After  $k$  iterations of Q-N method, we have

$$\begin{aligned}(1 - \delta)(f(x_k^*) - f^*)f^* &\leq (1 - \delta)(\hat{f}(x_k^*) - \hat{f}(x^*)) \\ &\leq \delta\hat{f}(x^*) + \frac{L^2R^2}{2n[e^{\delta(k+1)/n} - 1]} \stackrel{(?)}{\leq} \delta(f^*)^2.\end{aligned}$$

Thus, we need  $k = R_n(\delta) \stackrel{\text{def}}{=} \frac{n}{\delta} \ln \left( 1 + \frac{L^2R^2}{n\delta(f^*)^2} \right)$  iterations.

**Note:**

- The main factor  $\frac{n}{\delta}$  does not depend on the data. (*Fully polynomial approximation scheme.*)
- Dependence in  $n$  is the same as for optimal methods.
- Each iteration of Q-N method is very simple, same as in Ellipsoid Method (complexity  $O(n^2 \ln \frac{LR}{\delta f^*})$  iterations).
- We have a uniform bound  $R_n(\delta) < R_\infty(\delta) \stackrel{\text{def}}{=} \frac{L^2R^2}{\delta^2(f^*)^2}$ .

# Absolute accuracy

**Our goal:** for problem  $\min_{x \in Q} \phi(x)$  find  $\bar{x} \in Q : \phi(\bar{x}) \leq \phi^* + \epsilon$ .

Assume  $\phi$  has bounded subgradients and  $\|x - x_0\| \leq R, \forall x \in Q$ .

Define now a new SP-objective

$$\begin{aligned} f(x) &= \max\{\phi(x) - \phi(x_0) + 2LR, L\|x - x_0\|\} \\ &= \phi(x) - \phi(x_0) + 2LR \quad \forall x \in Q. \end{aligned}$$

Applying now Q-N method to  $\hat{f}$ , we get

$$\begin{aligned} \phi(x_k^*) - \phi^* &= f(x_k^*) - f^* \leq \frac{\delta f^*}{2(1-\delta)} + \frac{L^2 R^2}{2n[e^{\delta(k+1)/n} - 1] \cdot (1-\delta) f^*} \\ &\leq LR \left[ \frac{\delta}{1-\delta} + \frac{1}{2n[e^{\delta(k+1)/n} - 1] \cdot (1-\delta)} \right]. \end{aligned}$$

# Choice of parameters

Let us find  $\delta = \delta(\epsilon)$  from equation

$$\frac{\delta}{1-\delta} = \frac{\epsilon}{2LR} \Rightarrow \delta(\epsilon) = \frac{\epsilon}{\epsilon+2LR}.$$

Then, we need at most

$$\begin{aligned} k &= R_n(\epsilon) \stackrel{\text{def}}{=} \frac{n}{\delta(\epsilon)} \ln \left( 1 + \frac{LR}{n\epsilon(1-\delta(\epsilon))} \right) \\ &= n \left( 1 + 2\frac{LR}{\epsilon} \right) \cdot \ln \left( 1 + \frac{\epsilon+2LR}{2n\epsilon} \right) \end{aligned}$$

iterations. At the same time,

$$R_n(\epsilon) < R_\infty(\epsilon) = \frac{1}{2} \left( 1 + 2\frac{LR}{\epsilon} \right)^2.$$

**Note:** worst-case complexity bound of Q-N method is always better than that of the standard subgradient scheme.

Schemes of Q-N methods look very natural.

**1. Minimization in relative scale:** (No parameters!)

$$x_{k+1} = \arg \min_{x \in Q} \left[ \|x - x_0\|_{G_0}^2 + \frac{\delta}{1-\delta} \sum_{i=0}^k \frac{(f(x_i) + \langle f'(x_i), x - x_i \rangle)^2}{(\|f'(x_i)\|_{G_i}^*)^2} \right].$$

**2. Minimization in absolute scale:**

$$x_{k+1} = \arg \min_{x \in Q} \left[ \|x - x_0\|_{G_0}^2 + \frac{\epsilon}{LR} \sum_{i=0}^k \frac{(\phi(x_i) - \phi(x_0) + \langle \phi'(x_i), x - x_i \rangle + 2LR)^2}{(\|\phi'(x_i)\|_{G_i}^*)^2} \right].$$

**Compare:** Dual gradient method

$$x_{k+1} = \arg \min_{x \in Q} \left[ \|x - x_0\|_{G_0}^2 + \frac{2}{L_k} \sum_{i=0}^k (\phi(x_i) + \langle \phi'(x_i), x - x_i \rangle) \right].$$

for  $C_L^{1,1}$  :  $L_k \equiv L$ , and for  $C_M^{1,0}$  :  $L_k \approx \sqrt{k} \cdot \frac{M}{R}$ .

# Revival of old questions

1. Roles of *dimension* and *accuracy* in complexity estimates.  
Denote by  $\mathcal{T}$  complexity of the oracle.

Available methods	Complexity
Subgradient*	$\frac{L^2 R^2}{\epsilon^2} \cdot (n + \mathcal{T})$
Quasi-Newton	$\frac{nLR}{\epsilon} \ln \left(1 + \frac{LR}{n\epsilon}\right) \cdot (n^2 + \mathcal{T})$
Ellipsoids	$n^2 \ln \frac{LR}{\epsilon} \cdot (n^2 + \mathcal{T})$
Inscribed ellipsoids*	$n \ln \frac{LR}{\epsilon} \cdot (n^3 + \mathcal{T})$

If  $\frac{1}{\epsilon} < n \ln \frac{1}{\epsilon}$  and  $\mathcal{T} \approx n^2$ , then QN is the best.

## Questions:

- 1 What are the best methods for all spectrum of  $n$ ,  $\epsilon$ , and  $\mathcal{T}$ ?
- 2 Are  $n$ ,  $\epsilon$ , and  $\mathcal{T}$  really independent?  $\Rightarrow$  Accuracy of the model?  
(Finite elements, Truss topology, Optimal Control, PDE, etc.)

## 2. Do we have a future for Quasi-Newton methods?

- Two decades of intensive research (60's, 70's).  
Good computational performance.
- $r$ -algorithm by Shor for nonsmooth minimization.
- Excellence and failure of Ellipsoid Method.  
(Wrong application?)
- 25 years of complete silence.

*Can we finally do a proper global complexity analysis for these schemes?*