#### Robust Uncertainty Principles and Optimally Sparse Decompositions

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# **Uncertainty Principles**

- Heisenberg (1927) Uncertainty principle for continuous-time signals  $\hat{f}$  $\boldsymbol{f}$  $\sigma_{t}$ σ  $\boldsymbol{\sigma_t \sigma_\omega} \geq \frac{1}{4\pi}$
- Limits *joint resolution* in time and frequency

# **Uncertainty Principles**

• Donoho and Stark (1989) Discrete uncertainty principle for  $\mathbb{C}^N$ 



- Implications: recovery from partial information, unique sparse decompositions
- Generalization to pairs of bases  $B_1, B_2$ [Donoho,Huo,Elad,Bruckstein,Gribonval,Nielsen]

# **Dirac Comb**

The discrete uncertainty principle is *exact*.



- $\sqrt{N}$  spikes spaced  $\sqrt{N}$  apart
- Invariant under Fourier transform  $(f = \hat{f})$
- $|T| + |\Omega| = 2\sqrt{N}$

#### **Sparse Representations**

• Decompose *f* as a superposition of spikes and sinusoids

$$f(s) = \sum_{t \in T} \alpha_t \delta(t - s) + \sum_{\omega \in \Omega} \alpha_\omega e^{i\omega s}$$

T = locations of spikes,  $\Omega$  = frequencies of sinusoids

• Matrix form:

$$\Phi oldsymbol{lpha} = oldsymbol{f} \quad \Leftrightarrow \quad igg( oldsymbol{I} \quad oldsymbol{F}^st igg) igg( oldsymbol{lpha}_t \ oldsymbol{lpha}_{oldsymbol{\omega}} igg) = oldsymbol{f}$$

• Many solutions exist, we want the *sparsest*.

# Example



spikes

t 
ightarrow

# Image Processing and Sparse Representations



- Geometrical structure and texture are separate phenomena
- Each is sparse in a certain "basis"
- See also Geometric Separation of Donoho et. al, edge/texture separation of Starck et. al.



## $\ell_0$ -minimization

• Dictionary  $\Phi$ : columns are spikes and sinusoids

$$\Phi oldsymbol{lpha} = oldsymbol{f} \quad \Leftrightarrow \quad igg( oldsymbol{I} \quad oldsymbol{F}^* igg) igg( oldsymbol{lpha}_t \ oldsymbol{lpha}_{\omega} igg) = oldsymbol{f}$$

- Observed signal f
- Sparsest representation: solve (combinatorial problem)

$$\min \|\boldsymbol{\alpha}\|_{\boldsymbol{\ell}_0}$$
( $\boldsymbol{\ell}_0$ -min)  
s.t.  $\Phi \boldsymbol{\alpha} = \boldsymbol{f}$ 

- $\|\alpha\|_{\ell_0} = |\operatorname{supp} \alpha|$  = number of non-zero terms
- When is  $\ell_0$ -min unique?

#### Uncertainty Principles and Sparsity [Donoho,Huo (2001)]

**Fact:** Say  $\alpha$  is a sparse decomposition of f  $\|\alpha\|_{\ell_0} = |\operatorname{supp} \alpha| < \sqrt{N}$ .

$$f = \Phi \alpha$$

then  $\alpha$  is the *only* sparse decomposition.

**Reason:** Uncertainty Principle

 $|\operatorname{supp} f| + |\operatorname{supp} \hat{f}| \ge 2\sqrt{N}$ 

Say  $\Phi \alpha' = f$  as well:  $\alpha' = \alpha + \gamma$  with  $\Phi \gamma = 0$ 

Consequence:  $\alpha$  is the *unique*  $\ell_0$  minimizer.

## **Dirac Comb**

The "Dirac comb" has two representations of size  $\sqrt{N}$ .



#### Sharp Uncertainty Principle: N Prime [Tao (2004)]

- For N prime, evenly spaced signals like the Dirac comb are impossible.
- The uncertainty principle is *significantly* more relaxed:

 $|\operatorname{supp} f| + |\operatorname{supp} \hat{f}| > N$ 

• Key: minors of the Fourier matrix have full rank

$$oldsymbol{A} = oldsymbol{R}_\Omega oldsymbol{F} oldsymbol{R}_T^st$$

 $R_T, R_\Omega$  are restriction operators.

• Compare to general UP: N vs.  $2\sqrt{N}$ 

#### $\ell_0$ uniqueness: *N* Prime

**Theorem:** Let *N* be prime. Say  $\alpha^{\sharp}$  is supported on  $T \cup \Omega$  with

$$\Phi lpha^{\sharp} = f \qquad |T| + |\Omega| = \| lpha^{\sharp} \|_{\ell_0} < N/2.$$

Then for all  $\alpha'$ ,  $\Phi \alpha' = f$ ,  $\alpha' \neq \alpha$ 

 $\|\alpha'\|_{\boldsymbol{\ell}_0} > N/2$ 

 $\Rightarrow \alpha^{\sharp}$  is the *unique*  $\ell_0$  minimizer.

- Follows directly from relaxed UP
- Compare to general  $\ell_0$  uniqueness result: N/2 vs.  $\sqrt{N}$

# "Typical" Signals

• Robust UP: For an overwhelming percentage of  $T, \Omega$  with

$$|T| + |\Omega| \sim rac{N}{\sqrt{\log N}}$$

it is *impossible* to find an f with supp f = T, supp  $\hat{f} = \Omega$ .

•  $\ell_0$ -uniqueness: Given "generic"  $\alpha$  on  $T \cup \Omega$ , solving

$$\min \|\boldsymbol{\alpha'}\|_{\boldsymbol{\ell}_0} \qquad \text{s.t.} \qquad \Phi \boldsymbol{\alpha'} = \boldsymbol{f}$$

will recover  $\alpha$  *exactly*.

• Tractability: If  $T, \Omega$  are slightly smaller

$$|T| + |\Omega| \sim \frac{N}{\log N}$$

then there is a *tractable* algorithm to recover  $\alpha$  from f.

• Similar results hold for pairs of bases  $B_1, B_2$ .

#### **Probabilistic Framework**

1. Generate T,  $\Omega$  with Bernoulli sequences

$$oldsymbol{\chi}(oldsymbol{t}),oldsymbol{\chi}(oldsymbol{\omega}) = egin{cases} 1 & extsf{w/prob} & au \ 0 & extsf{w/prob} & 1- au \ 0 & extsf{m/prob} & 1- au \end{cases}$$

$$egin{array}{rcl} m{T}&=&\{m{t}:m{\chi}(m{t})=1\}\ \ \Omega&=&\{m{\omega}:m{\chi}(m{\omega})=1\} \end{array}$$

We will have  $|T|, |\Omega| \sim \tau N$ .

2. Generate  $\alpha$  on  $T \cup \Omega$  from a continuous distribution.

We derive bounds that hold with very high probability.

#### A Quantitative Robust Uncertainty Principle

Draw sets  $T, \Omega$  at random with  $au = C_M / \sqrt{\log N}$ :

$$|T|, |\Omega| pprox C_M rac{N}{\sqrt{\log N}}$$

Then with probability  $1 - O(N^{-M})$ , for any f with

 $\operatorname{supp} \boldsymbol{f} = \boldsymbol{T}$ 

most of the energy of  $\hat{f}$  lies outside of  $\Omega$ :

$$\sum_{\boldsymbol{\omega}\in\Omega}|\hat{f}(\boldsymbol{\omega})|^2 \leq \frac{1}{2}\|\hat{f}\|^2.$$

• No  $f \in \mathbb{C}^N$  exists with supp f = T, supp  $\hat{f} = \Omega$ .

For relatively large sets T,Ω, it is impossible to concentrate
 f on T and f on Ω.

#### Key Estimate

• Again, 
$$A = R_{\Omega}FR_T^*$$
 plays a *key* role.

Since  $\operatorname{supp} f = T$ ,  $R_T^* f|_T = f$ , where  $f|_T = R_T f$ .

$$egin{aligned} &\sum_{oldsymbol{\omega}\in\Omega} |\hat{f}(oldsymbol{\omega})|^2 &= & \langle R_\Omega F^* R_T^* f|_T, R_\Omega F R_T^* f|_T 
angle \ &= & \langle f|_T, A^* A f|_T 
angle \ &= & \|A^* A\| \cdot \|f\|^2 \end{aligned}$$

• Need to show  $||A^*A|| \le 1/2$  (bound the largest eigenvalue).

#### Generic $\ell_0$ Uniqueness

**Theorem:** Let  $T, \Omega$  be sets for which the robust uncertainty principle holds,

$$|T|+|\Omega| pprox C_M rac{N}{\sqrt{\log N}}.$$

Choose  $\alpha$  on  $T \cup \Omega$  from a continuous distribution. Then given  $f = \Phi \alpha$ ,  $\alpha$  is the unique minimizer to

$$\min \|\boldsymbol{\alpha'}\|_{\boldsymbol{\ell}_0} \qquad \text{s.t.} \qquad \Phi \boldsymbol{\alpha'} = \boldsymbol{f}$$

with probability 1.

• The uncertainty principle is fundamental, but the uniqueness is not immediate.

# $\boldsymbol{\ell}_1$ minimization

•  $\ell_0$ -min is highly non-convex

 $\rightarrow$  combinatorial algorithm required to solve it in general.

• Instead, consider the convex program

$$\min \|\boldsymbol{\alpha}\|_{\boldsymbol{\ell}_1}$$
 ( $\boldsymbol{\ell}_1$ -min) s.t.  $\Phi \boldsymbol{\alpha} = \boldsymbol{s}$ 

- $\ell_1$  norm also a sparsity measure
- $\ell_1$ -min *convex*
- When are the solutions to min- $\ell_0$  and min- $\ell_1$  the same?

## **Previous Work**

- Donoho, Stark: uncertainty principles, sparse recovery via  $\ell_1$ -min
- Donoho, Huo:  $\ell_1$ -min for  $|\operatorname{supp} \alpha| \leq \frac{\sqrt{N}}{2}$
- Bruckstein, Elad: general basis pairs,  $\ell_1$ -min for  $|\operatorname{supp} \alpha| \leq .9\sqrt{N}$
- Gribonval, Nielsen, Donoho, Elad: general dictionaries
- Tropp: greedy algorithms, robust recovery
- Gilbert, Strauss et al. (sparse Fourier approx, random algorithm)

## Generic $\ell_1$ recovery

Theorem: Draw sets  $T, \Omega$  at random with  $au = C_M / \log N$ 

$$|T| + |\Omega| pprox C_M rac{N}{\log N}.$$

Choose  $\alpha$  on  $T \cup \Omega$  from a continuous distribution with uniform phase. Then given  $f = \Phi \alpha$ ,  $\alpha$  is the unique minimizer to

$$\min \|\boldsymbol{\alpha'}\|_{\boldsymbol{\ell}_1} \qquad \text{s.t.} \qquad \Phi \boldsymbol{\alpha'} = \boldsymbol{f}$$

with probability  $1 - O(N^{-M})$ .

We can recover "sparse" decompositions with

 $|T| + |\Omega| \sim N/\log N$ 

(compare to  $\sim \sqrt{N}$ )

## Duality

$$\min \|\boldsymbol{\alpha}\|_{\boldsymbol{\ell}_1} \qquad (\boldsymbol{\ell}_1 \text{-min})$$
  
s.t.  $\Phi \boldsymbol{\alpha} = \boldsymbol{f}$ 

•  $\alpha^{\sharp}$  solves  $\ell_1$ -min  $\Leftrightarrow \exists \lambda$  such that for  $P(k) = (\Phi^*\lambda)(k)$ 

$$P(k) = \operatorname{sgn}(\alpha^{\sharp})(k)$$
 on  $T \cup \Omega$  (1)  
 $|P(k)| < 1$  on  $(T \cup \Omega)^{c}$ . (2)

- Sufficient: construct min-energy  $\lambda$  s.t. (1) is satisfied. Check:
  - if it exists
  - if (2) holds
- See also Fuchs (04)

# Example



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#### **Constructing the Dual Vector**

$$\begin{aligned} \boldsymbol{P}(\boldsymbol{k}) &\in \operatorname{Range}(\Phi^*) \\ \boldsymbol{P}(\boldsymbol{k}) &= \operatorname{sgn}(\boldsymbol{\alpha}^{\sharp})(\boldsymbol{k}) \quad \text{on} \quad \Gamma := \boldsymbol{T} \cup \Omega \\ \boldsymbol{P}(\boldsymbol{k})| &< 1 \qquad \text{on} \quad (\boldsymbol{T} \cup \Omega)^c. \end{aligned}$$

• Minimum energy  $\lambda$  with  $R_{\Gamma} \Phi^* \lambda = R_{\Gamma} \operatorname{sgn}(\alpha^{\sharp})$  ("least squares"):

$$\boldsymbol{\lambda} = \Phi \boldsymbol{R}_{\Gamma}^{*} (\boldsymbol{R}_{\Gamma} \Phi^{*} \Phi \boldsymbol{R}_{\Gamma}^{*})^{-1} \boldsymbol{R}_{\Gamma} \operatorname{sgn}(\boldsymbol{\alpha}^{\sharp})$$

• Need to invert

$$(\boldsymbol{R}_{\Gamma}\Phi^{*}\Phi\boldsymbol{R}_{\Gamma}^{*})^{-1} = \begin{pmatrix} \boldsymbol{I} & \boldsymbol{A}^{*} \\ \boldsymbol{A} & \boldsymbol{I} \end{pmatrix}^{-1}, \qquad \|(\boldsymbol{R}_{\Gamma}\Phi^{*}\Phi\boldsymbol{R}_{\Gamma}^{*})^{-1}\| = \frac{1}{1-\sqrt{\|\boldsymbol{A}^{*}\boldsymbol{A}\|}}$$

with  $A = R_{\Omega}FR_T^*$ .

• Eigenvalues of  $A^*A$  again play a key role

$$A^*A = R_T F^* R^*_\Omega R_\Omega F R^*_T \qquad |T| imes |T|$$

Rewrite

$$oldsymbol{A}^{*}oldsymbol{A} = rac{\Omega}{oldsymbol{N}}oldsymbol{I} - rac{1}{oldsymbol{N}}oldsymbol{H}$$

where *H* is the *random matrix* 

$$m{H}(m{t},m{t'}) = egin{cases} 0 & m{t} = m{t'} \ -\sum_{m{\omega}\in\Omega} e^{m{i}m{\omega}(m{t}-m{t'})} & m{t} 
eq m{t'} \end{cases}$$

How big can we make |T| and  $|\Omega|$  and keep the eigenvalues under control?

 $\|\boldsymbol{H}\|<<|\Omega|$ 

# |P(k)| < 1

$$\boldsymbol{P} = \Phi^* \Phi \boldsymbol{R}^*_{\Gamma} (\boldsymbol{R}_{\Gamma} \Phi^* \Phi \boldsymbol{R}^*_{\Gamma})^{-1} \boldsymbol{R}_{\Gamma} \operatorname{sgn}(\boldsymbol{\alpha}^{\sharp})$$

- $(\mathbf{R}_{\Gamma}\Phi^{*}\Phi\mathbf{R}_{\Gamma}^{*})^{-1}$  is very well conditioned
- For  $|T|, |\Omega| \sim N/\log N$ , the vector  $(R_{\Gamma} \Phi^* \Phi R_{\Gamma}^*)^{-1} R_{\Gamma} \operatorname{sgn}(\alpha^{\sharp})$  has small norm
- Use simple large deviation theory for  $\Phi^*\Phi$  acting on a "small" vector

# Example



• In practice, we can recover "sparse" decompositions of size  $\sim N/2$  (!!)

# Example



#### **Recovery Curves**



# Pairs of Bases

Orthogonal bases  $B_1, B_2, \Phi = \begin{pmatrix} B_1^* & B_2^* \end{pmatrix}$ . Define

$$oldsymbol{\mu} = \sqrt{N} \cdot \max \, \left| \langle oldsymbol{b}_1, oldsymbol{b}_2 
ight| \qquad oldsymbol{b}_1 \in oldsymbol{B}_1, \ oldsymbol{b}_2 \in oldsymbol{B}_2.$$

With high probability:

• RUP: For  $oldsymbol{f}=oldsymbol{B}_1 oldsymbol{lpha}_1, oldsymbol{B}_2 oldsymbol{lpha}_2$ , we have

$$|\operatorname{supp} lpha_1| + |\operatorname{supp} lpha_2| \sim rac{N}{\mu^2 (\log N)^p}$$

•  $\ell_0$  uniqueness: Given  $\alpha := \begin{pmatrix} \alpha_1 & \alpha_2 \end{pmatrix}$  supported on  $\Gamma = \Gamma_1 \cup \Gamma_2$ , set  $f = \Phi \alpha$ . For  $|\Gamma| \sim \frac{N}{\mu^2 (\log N)^p}$ 

$$\alpha$$
 will be the unique  $\ell_0$  minimizer.

•  $\ell_1$  equivalence: For  $\Gamma$  as above,  $\alpha$  will also be the unique minimizer of the convex  $\ell_1$  program.

## **Other Examples**

- $B_1 = spikes$ ,  $B_2 = sinusoids$
- B<sub>1</sub> = wavelets, B<sub>2</sub> = sinusoids at fine scales, wavelets look like spikes
- $B_1$  = spikes,  $B_2$  = random basis performance is very close to spikes and sinusoids

# Conclusions

- Sparse decompositions are *unique* 
  - we can recover them from an observation of the signal
- Key: uncertainty principles
- The recovery is performed via convex optimization (min- $\ell_1$ )
- "Typical" case bounds (  $\sim N/\log N$  ) much more lenient than worst case (  $\sim \sqrt{N}$  )
- Derivation: show when the dual vector exists
- Bounds can be generalized to pairs of bases

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