

THE COMPUTATION OF β_Q

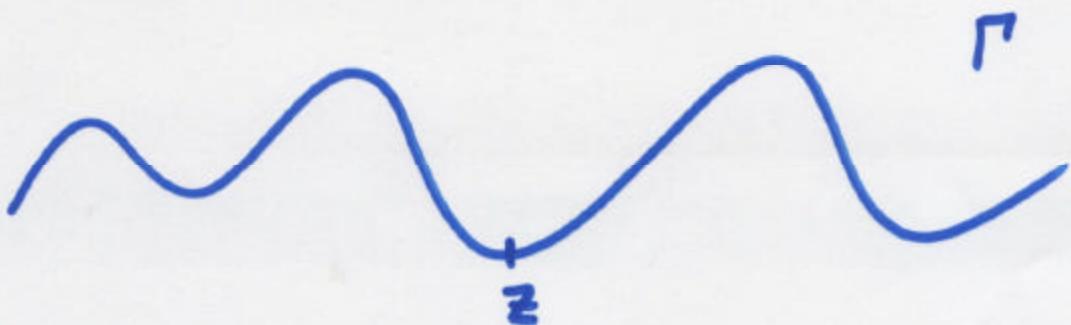
- Let $N = |K \cap Q|$.
- Let A be an $N \times n$ matrix whose rows are the points of $K \cap Q$.
- The axes of the ellipsoid are the singular vectors of A .
- β_Q is determined by the singular values of A .
- β_Q is computed in $O(n^2 \cdot N + n^3)$ operations.

Model Problems Related to β

Wolff Schmitz,
Bogomolov BineG

1. Cauchy Integrals on Lipschitz Curves
2. Analytic Capacity
3. Estimates for Brownian Motion
4. Estimates (Non-Sharp!) for Dimension
5. Mumford - Shah Conjecture
(Work of David - Semmes on $K, \mathcal{D}^1(K), \dots$)

Model Problems: Cauchy Integ. on Lip Curves, Analyt. Cap. (Tolsa, NTV),
Sharp estimates for w (Schw, Brownian Motion)



$$\mathcal{E}f(z) = \frac{1}{2\pi i} \int\limits_{\Gamma} \frac{f(w) dw}{w - z}$$

Theorem (Calderón, C-M-M)

If Γ is a Lipschitz curve, then

$$\|\mathcal{E}f\|_{L^2(\Gamma)} \leq A \|f\|_{L^2(\Gamma)}$$

A depends on Lipschitz Const.

$\mathcal{E} \approx H(\text{diagonal}) + \text{Off Diagonal}$

Off Diag. Controlled by $\sum \beta(Q) l(Q)$



Def. $\gamma(K) > 0 \iff$

$\exists F \not\equiv \text{Constant},$
 F bounded, analytic on K^c

Surprise: This is an L^2 problem!

Critical Case: $\dim(K) = 1$

$$\ell(K) = 0 \Rightarrow \gamma(K) = 0$$

1. (David) If $0 < \ell(K) < +\infty$,

$$\gamma(K) = 0 \iff \ell(K \cap \Gamma) = 0$$

for all curves
 $\Gamma, \ell(\Gamma) < +\infty$

Melnikov's Philosophy

M. Conjecture: $\gamma(K) > 0 \iff \exists u \in \text{Pr}(K)$ s.t.

$$(1) \quad u(D(x, r)) \leq Ar$$

$$(2) \quad \iiint c^2(x, y, z) du(x) du(y) du(z) < \infty$$

Note: C.Z. theory shows \leftarrow
 Hard Part is to show \rightarrow

Thm 1. says MC is true when
 K is a continuum.

(Then $\gamma(K) > 0$ because
 of Riemann mapping)

$$F: S^2 \setminus K \rightarrow \mathbb{D}$$

\uparrow unit disk
 unbounded component

Tolsa's Theorem: $\gamma(K) > 0$

$\iff \exists \mu \in \text{Pr}(K) :$

$$\mu(D(x,r)) \leq Ar$$

$$\iiint_{KKK} c^2(x,y,z) d\mu(x) d\mu(y) d\mu(z) < \infty$$



Problem: Given $\gamma(K) > 0$
(or μ as above),

CONSTRUCT $f \in H^\infty(K^c)$

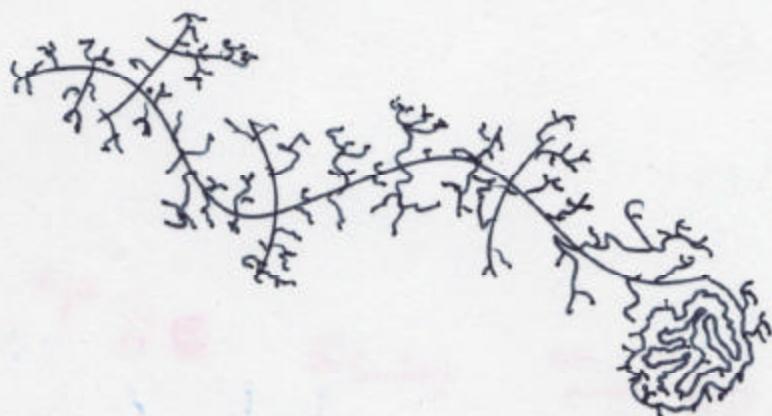
Theorem 2. (\mathcal{T})

$K \subset \mathbb{C}$ a continuum,

Diameter(K) = 1

$\Rightarrow \exists u \in \text{Pr}(K),$

$$\| u * \frac{1}{z} \|_{H^\infty(\mathbb{C} \setminus K)} \leq C_0$$



$H^\infty \leftrightarrow$ Holomorphic, Bounded
Sup Norm

Nazarov - Treil - Volberg '98

Found the correct extension
of the $\{\Psi_I\}$ method (C-J-S)
for non-doubling measures.

A. Volberg will give
a lecture series on this
plus more recent advances such
as:

$$K \subset \mathbb{R}^n$$

When is there $\Delta u = 0$
on K^c , $\nabla u \in L^\infty(K^c)$?

Answer: $\iff \exists \mu \in \text{Pr}(K)$

$$\mu * K \rightarrow \in L^\infty$$

\nwarrow this is a kernel!
no relation to the set.

Low-dimensional data sets

- Mixed dimensionalities.
- Non-smooth structure.
- Noise.
- Dependence on the distribution of samples.

distribution?

MOTIVATION

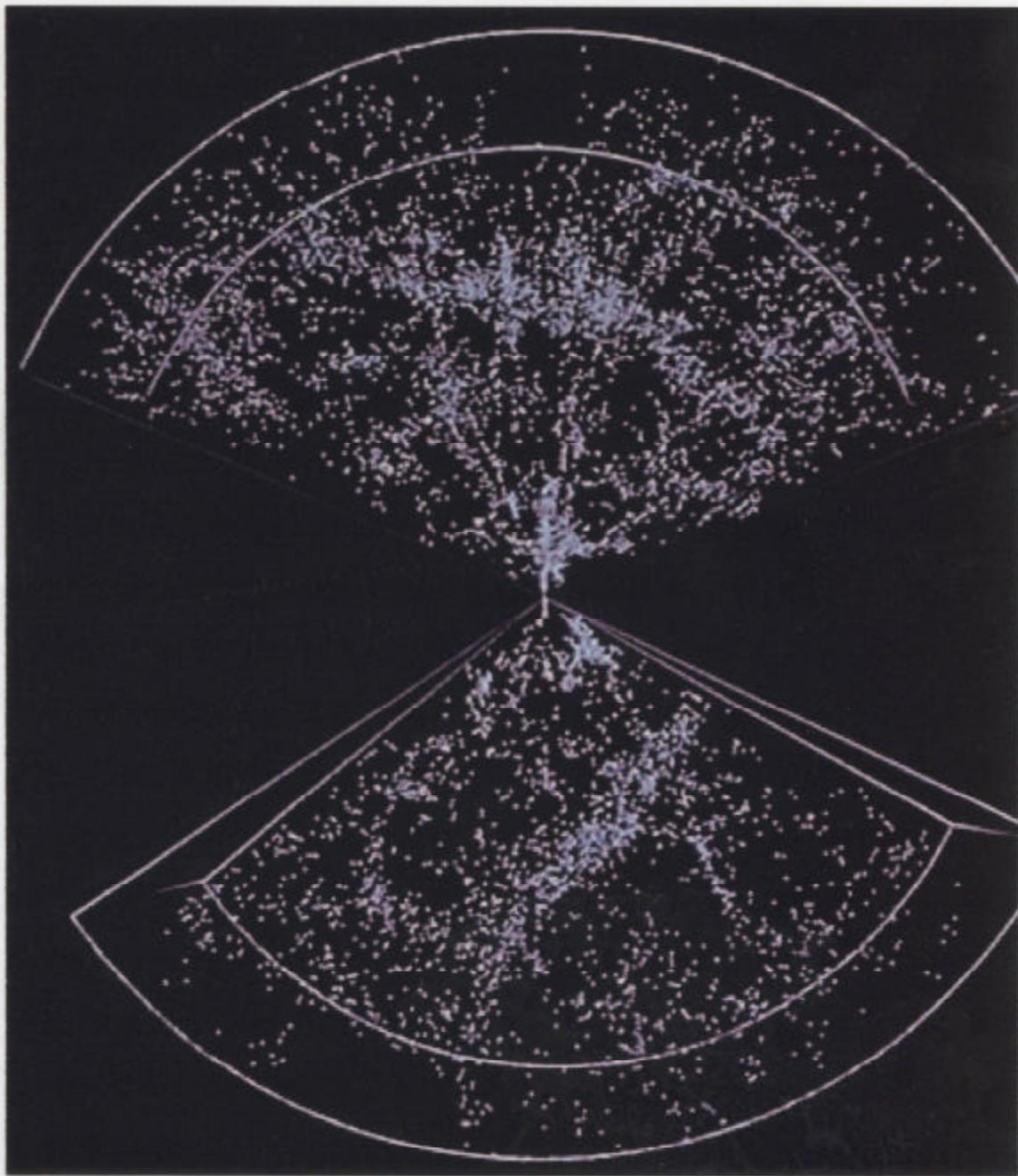
Goal

Efficient description of low dimensional

data sets in \mathbb{R}^n :

- Fast to compute.
- Insensitive to noise.
- Useful in analysis and synthesis of data sets.

arranged in thin sheets or long filaments. The longest sheet detect the "Great Wall," extends hundreds of millions of light years across maps.



North and South, Sheets and Voids

(Courtesy: Margaret J. Geller and Emilio E. Falco, Harvard-Smithsonian Astrophysics. Credits: Geller, da Costa, Huchra, and Falco.)
Large-scale structure in the universe in the northern and south

CURVE-LIKE SETS

Heuristic definition

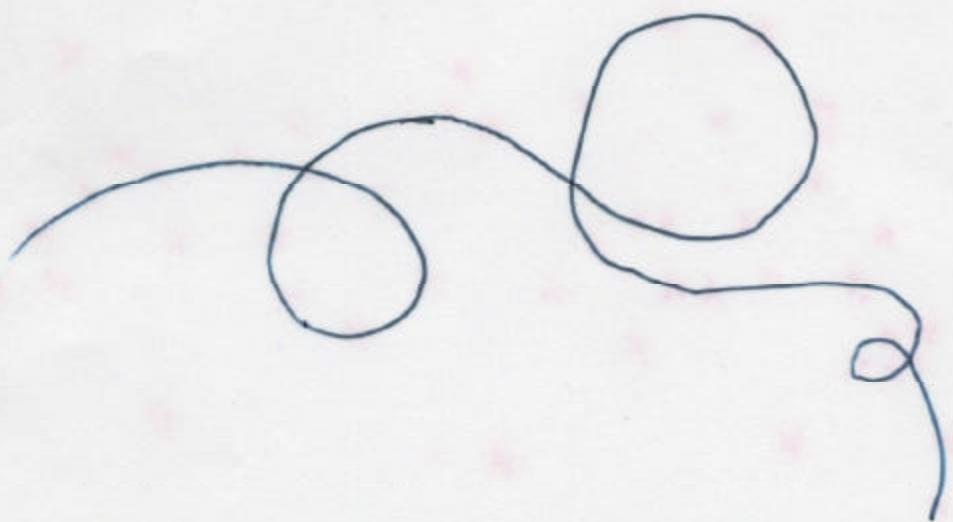
There exists a curve of short length containing a large fraction of the set.

Related question: approximate TSP

Construct a nearly shortest curve containing the given set.

Another definition

The set is well approximated by lines at different scales and locations.



Ǝ Short Curve Through
Many Points?



Approximate TSP

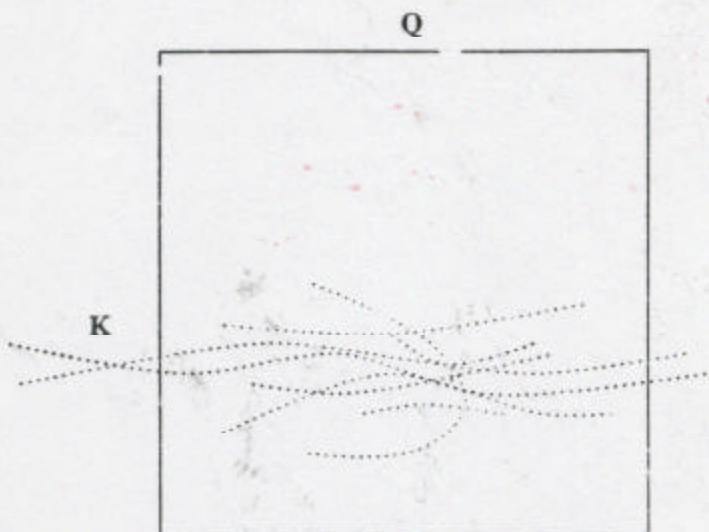
ONE-DIMENSIONAL GEOMETRIC TRANSCRIPTIONS

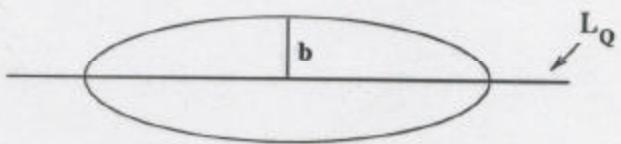
Consider a data set $K \subseteq \mathbb{R}^n$

β number for a cube Q

L_Q = best l_2 approximating line for $K \cap Q$

$$\beta_Q = \frac{\text{ l_2 -average distance of } K \cap Q \text{ from } L_Q}{l(Q)}$$





$$\beta_Q = \frac{b}{l(Q)}$$

$C \leq \beta \leq "1"$

9-a

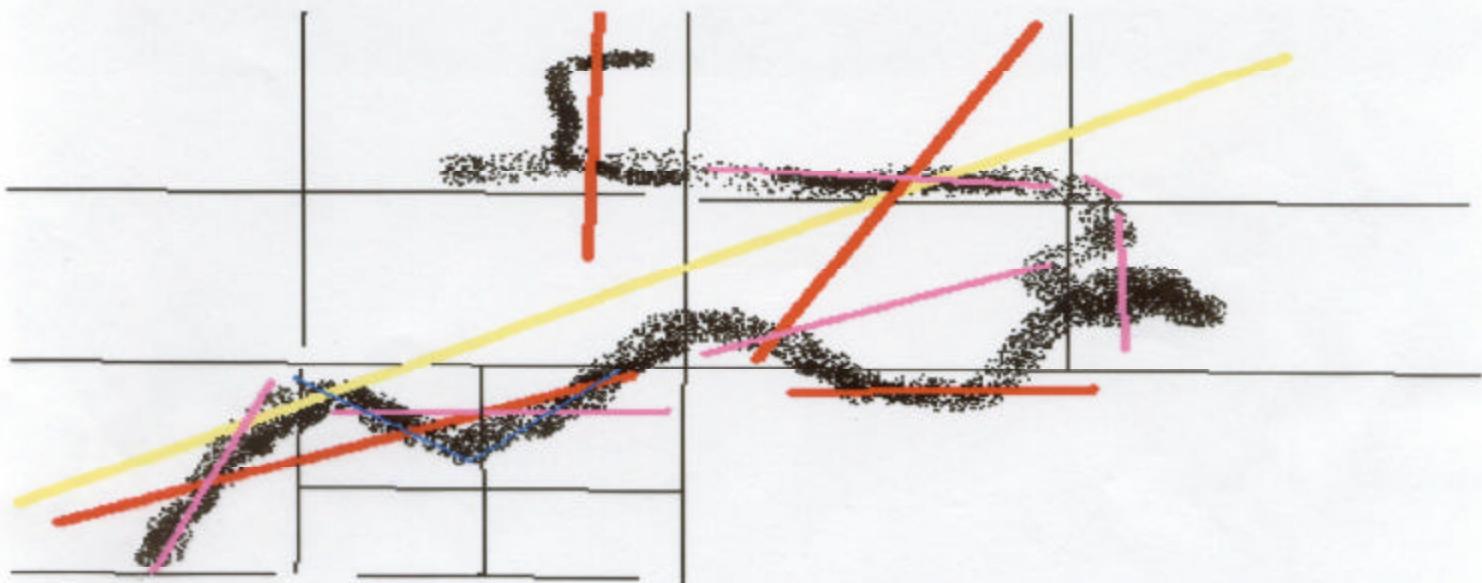
d-DIMENSIONAL GEOMETRIC TRANSCRIPTIONS

β number for a cube Q

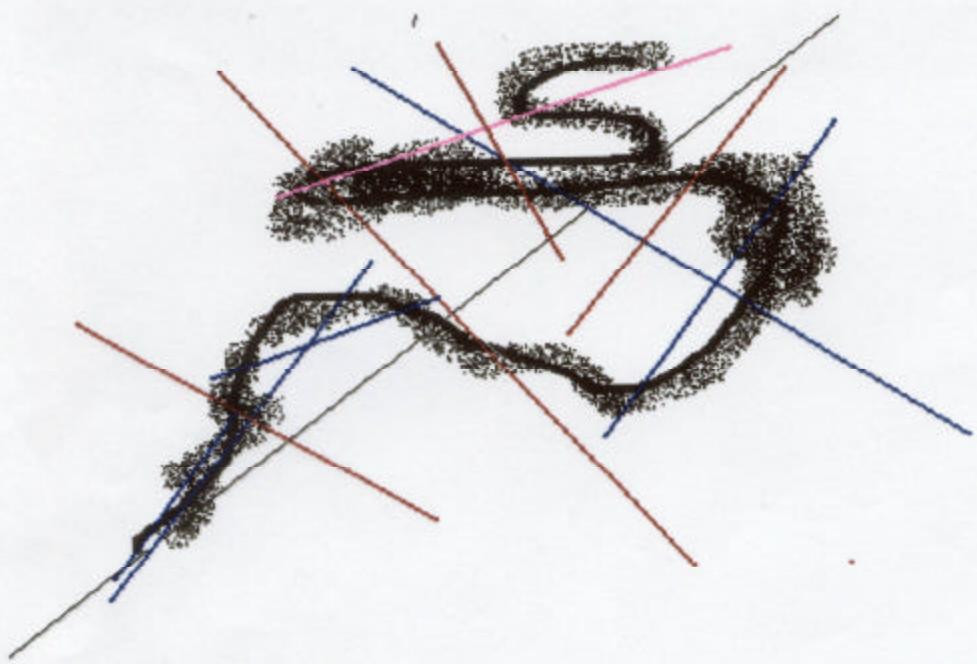
D_Q = best l_2 approximating d -plane for $K \cap Q$

$$\beta_Q = \frac{l_2\text{-average distance of } K \cap Q \text{ from } D_Q}{l(Q)}$$

- β_Q is determined by the singular values of the data matrix.



A multiscale approximation of a set in the plane (P.Jones)
The max of ratios between the side of a square and the sum
of deviations on all subsquares controls the bi Lipshitz
constant of a one dimensional parametrization of most of
the set



A data driven binary transcription of a set

$$J(x) = \sum_j \beta_{Q_j}^2$$

Q_j : dyadic cube containing x , $l(Q_j) = 2^{-j}$

J(x) measures rectifiability of K around x:

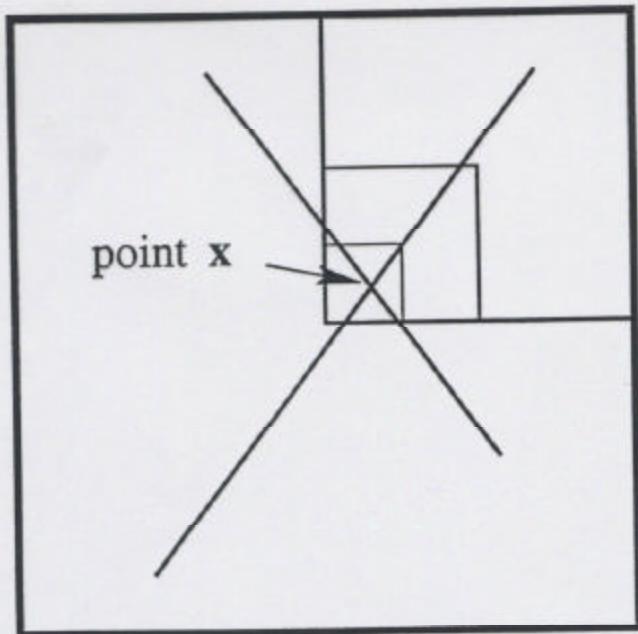
If $J(x) \leq M$ for all $x \in K$, then a portion of

K is contained in a curve, whose length is

controlled by M .

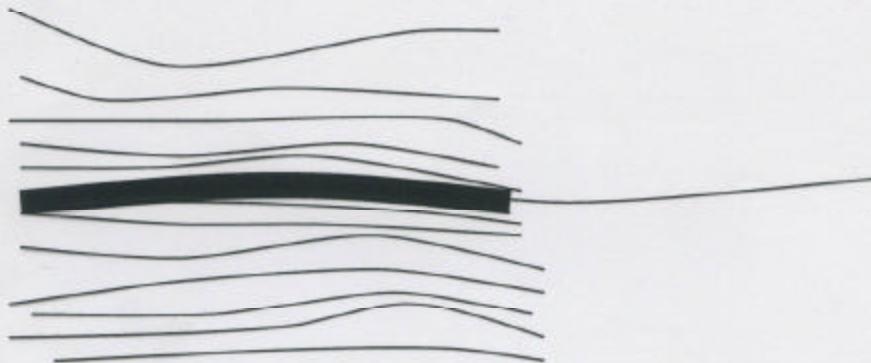
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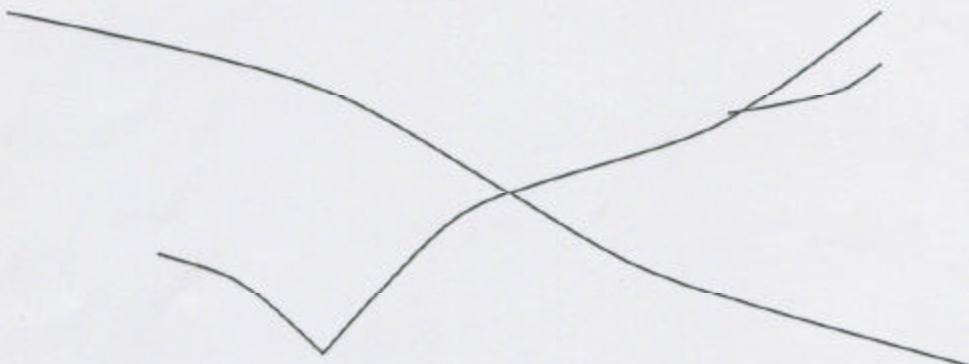


Where is $J(x)$ large ?

- Higher dimensional structure.



- Non-smooth structure (crossings, corners).



THEOREM (Lerman)

Let K be a set in \mathbb{R}^n with $\tilde{J}(x) \leq M$ for all $x \in K$. Then there exists a curve Γ such that

$$l(\Gamma) \leq c_1 e^{c_2 M} \cdot \text{diam}(K)$$

and

$$|\Gamma \cap K| \geq \frac{1}{c_1} e^{-c_2 M} \cdot |K|.$$

Remarks:

- The theorem can be formulated for an infinite set with a probability measure.
- $c_1 = c_0 \sqrt{n}$, where c_0 is a universal constant.
Worst case scenario: $c_2 = c_0 n$. If there is “roughly” a D -dimensional set carrying K , then $c_2 = c_0 D$.
- Empirical bounds are better.

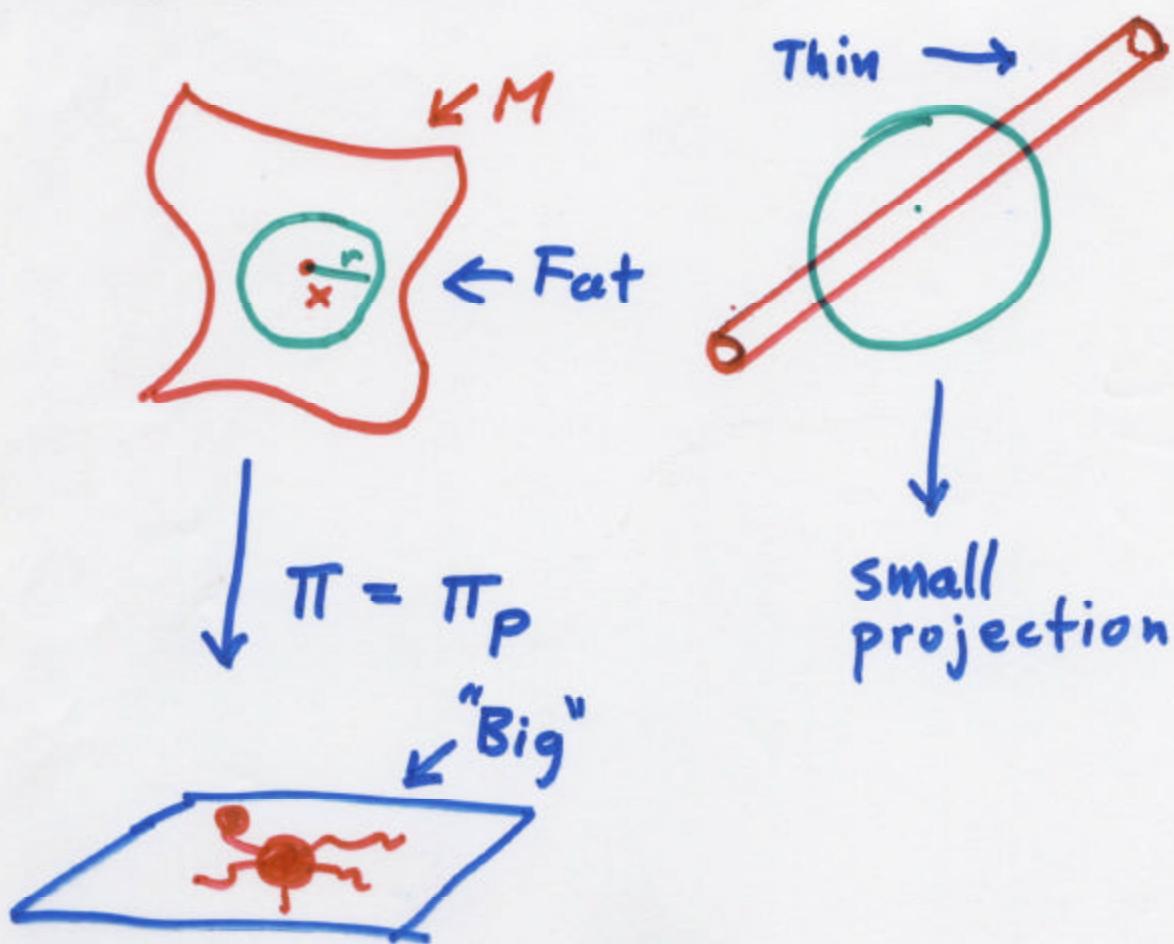
"Fat" d sets

M is "fat" for dimension d if for all $x \in M$

\exists a (d) hyperplane P such that

$$\pi(M \cap B(x, r))$$

contains a Ball $B(y, \varepsilon r)$ on P



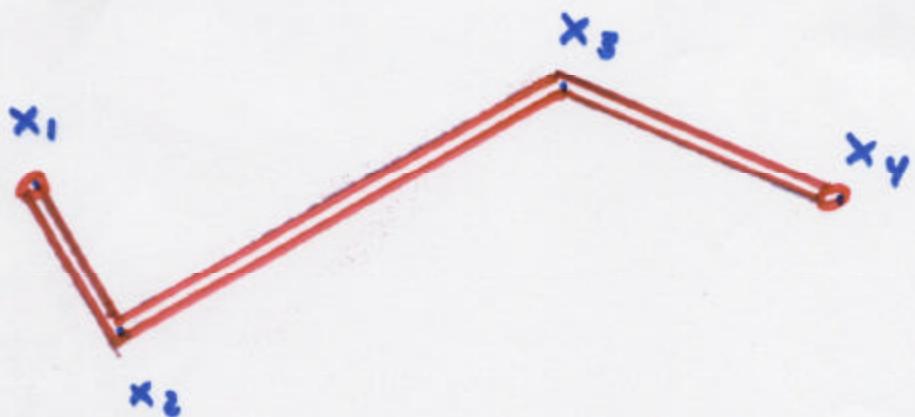
Why Demand Fat?

Let $\{x_1, \dots, x_N\} \subset \mathbb{R}^n$,

$$2 \leq d \leq n-1$$

\exists a difgeo S^d ($= M$) containing all points,

$$\text{Area}_d(M) < \varepsilon$$



Thin hairs attached to surfaces are the bane of GMT.

Th. (J, Lerman to appear)

Lerman CPAM for $d = 1$

μ on $[0,1]^n$

$$1 \leq d \leq n-1$$

$$\beta = \beta_2 = \beta_{2,d}$$

Approximate by d Planes

$$\int \exp \left\{ c_1 \sum_{Q \ni x} \hat{\beta}(Q) \right\} d\mu \leq A$$

$\Rightarrow \exists$ "Good" d set M s.t.

$$(1) \quad \mu(M) \geq A^{-1}$$

$$(2) \quad \text{Area}_d(M) \leq c_2 A$$

Comments: Best Possible (except c_1, c_2 !!)

"Good" = Fat "Glueing" of Lipschitz Graphs (+ regularity)

Given: A measure $\mu \in \mathcal{P}_r$
 or points $\{x_j\} \equiv K \subseteq \mathbb{R}^n$
 (Here set $\mu = \frac{1}{|K|} \sum \delta_{x_j}$)

Sample Problem: Which
 pieces of μ (points in K)
 are "one dimensional",
 "two dimensional", etc.?

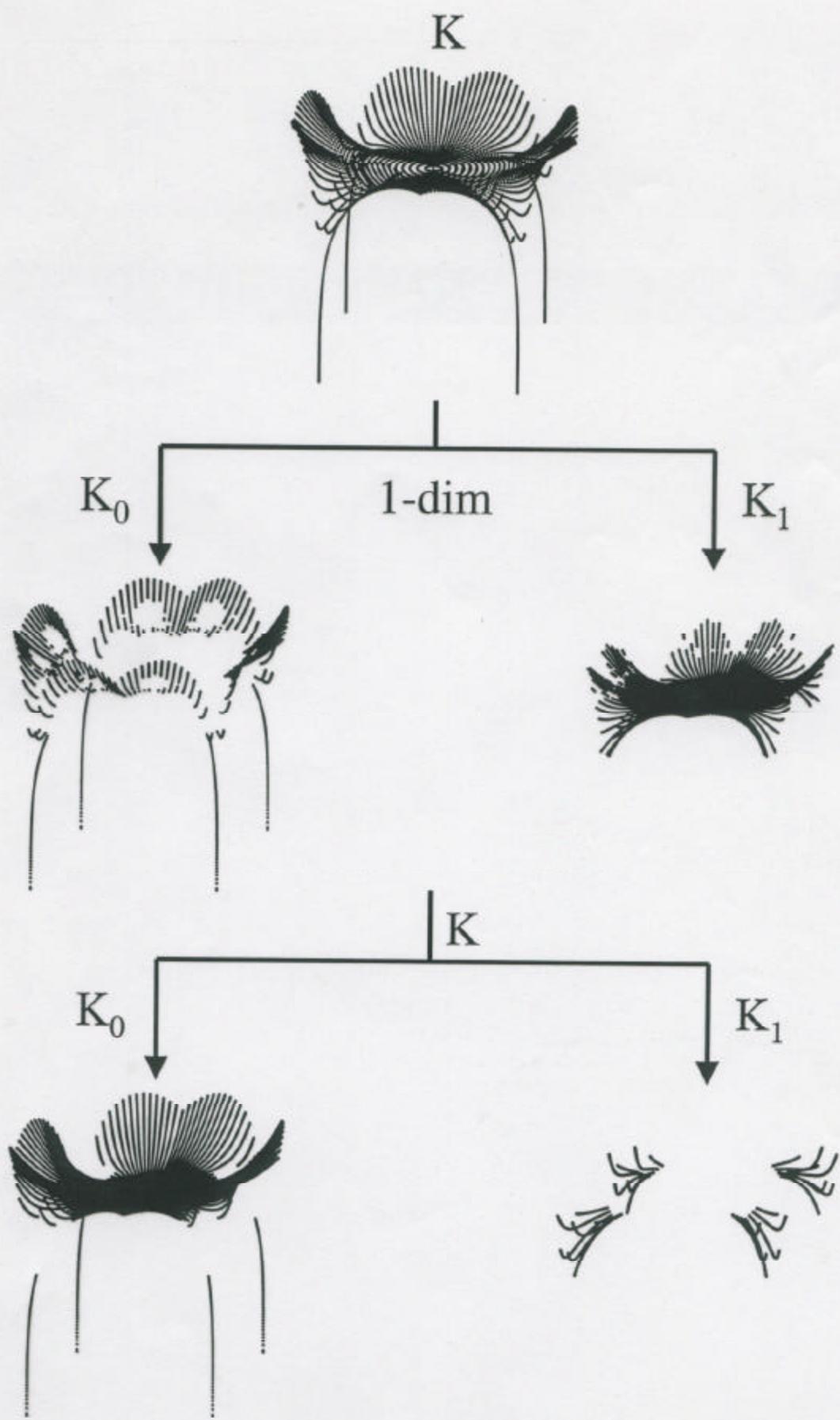
Answers: Find "short" curve
 (or "small" surface, etc.) hitting
 μ in large measure.

Philosophy (for Answer):

points in the "short" curve
 are "one dimensional"

DESCRIPTION OF THE ALGORITHM

1. If K has structure - stop at K .
2. Compute $J(x)$ for all $x \in K$.
3. Compute $M = \max_{x \in K} J(x)$
 $m = \min_{x \in K} J(x)$
4. Let $K_0 \equiv \{ x \mid m \leq J(x) < \frac{m+M}{2} \}$
 $K_1 \equiv \{ x \mid \frac{m+M}{2} \leq J(x) \leq M \}$
5. Repeat steps 1 – 5 for $K = K_0$
and $K = K_1$.



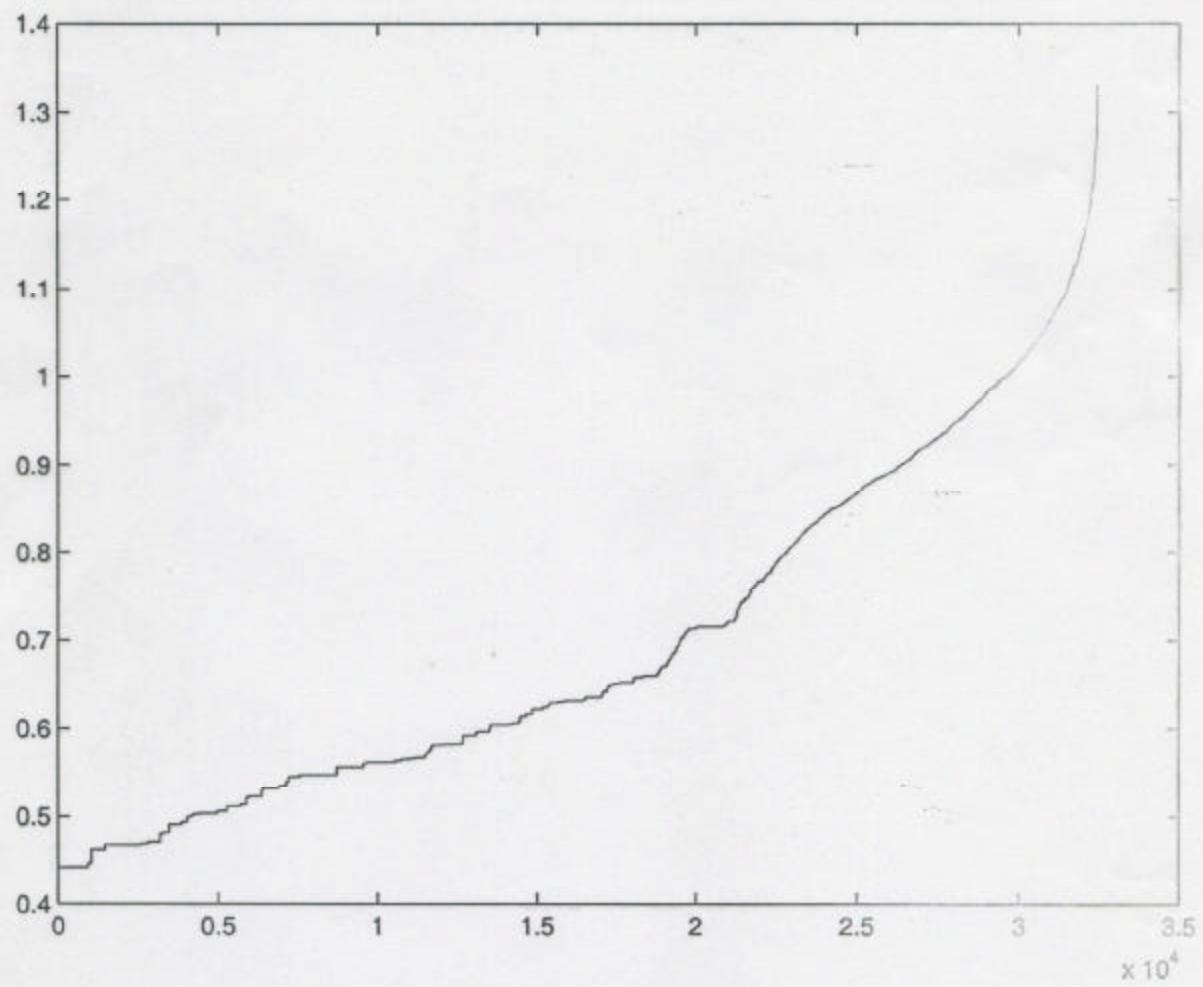
EXAMPLES OF SETS AND ANALYSES

EXAMPLE 1

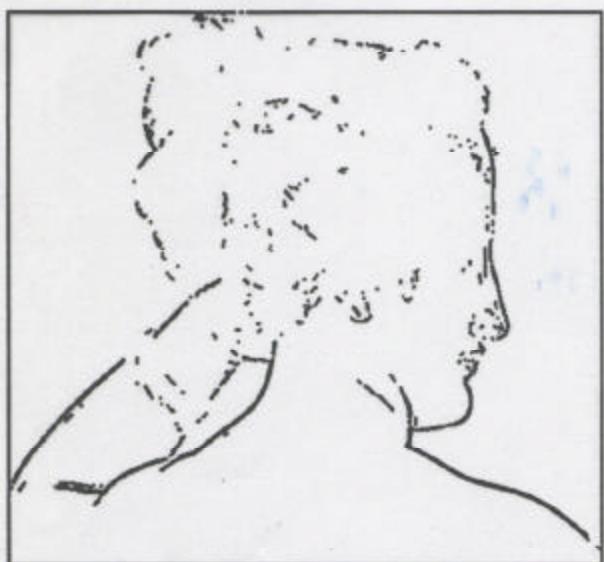
Input: $\approx 80,000$ points in the plane given by the edge map of the image "Paolina".

Output: Partition into subsets with different geometric structures.



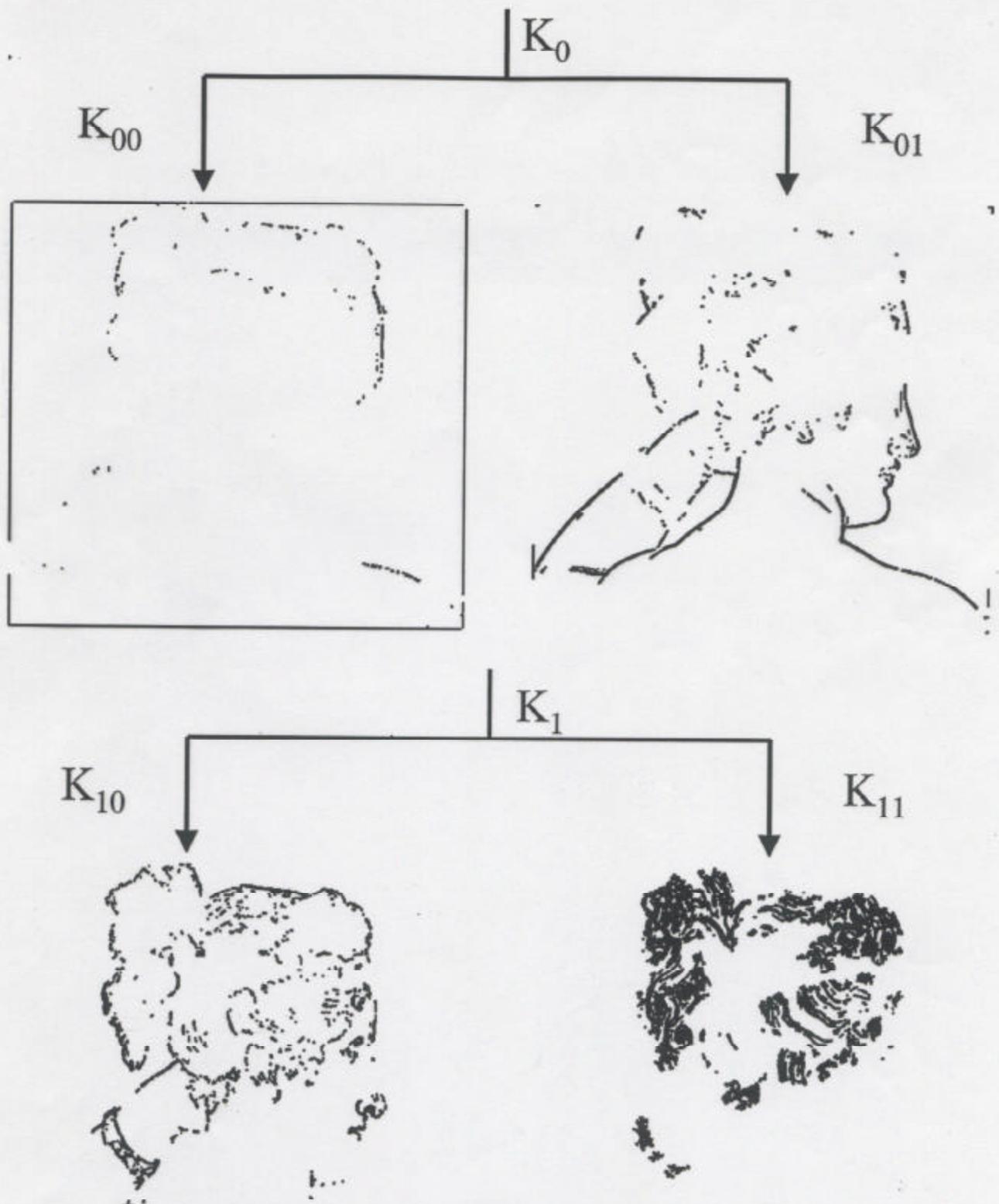


K

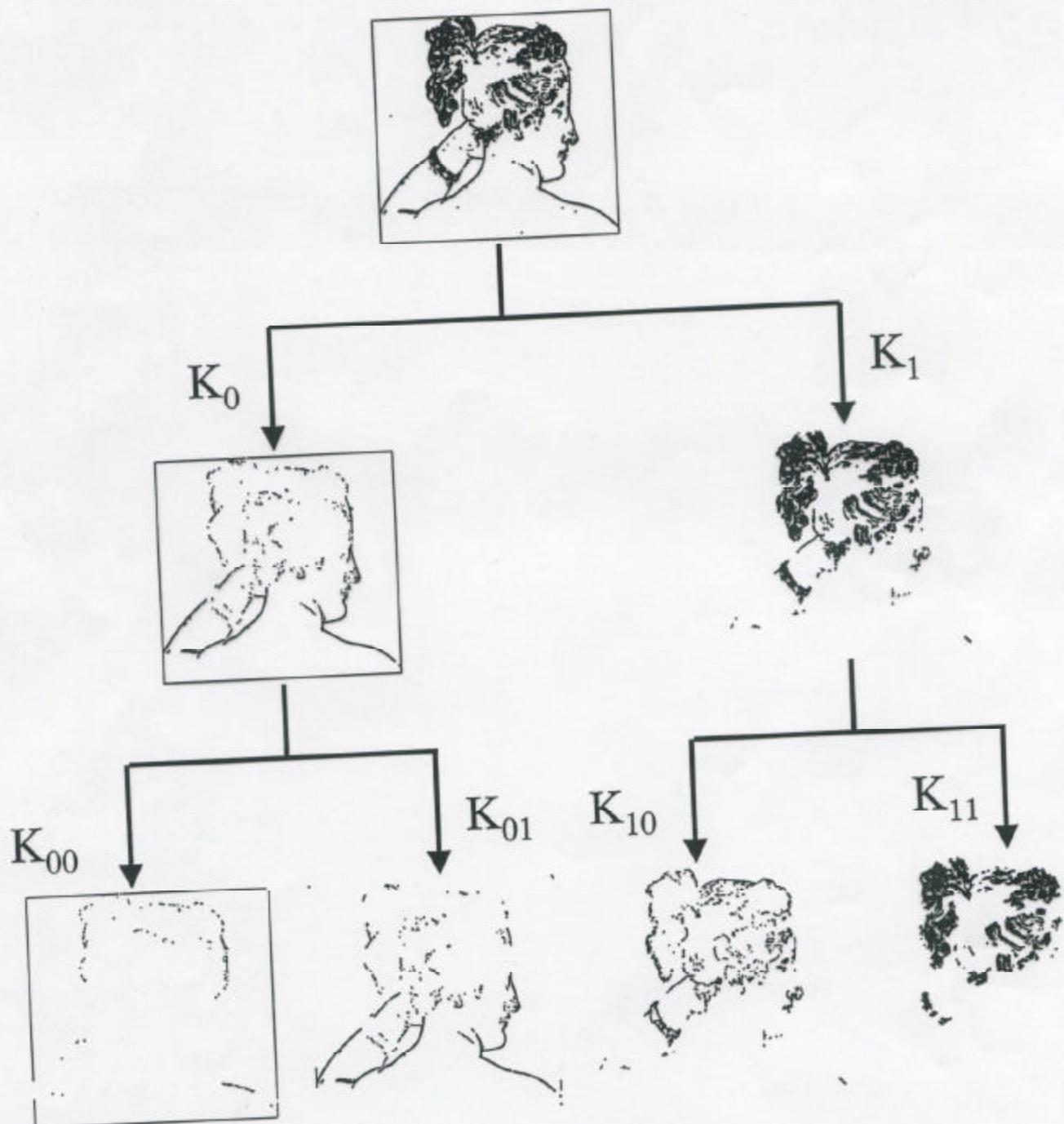
K₀K₁

$$\sum_{Q \ni x} \beta(Q) \leq M,$$

- The new sets are partitioned again in the same way.



K

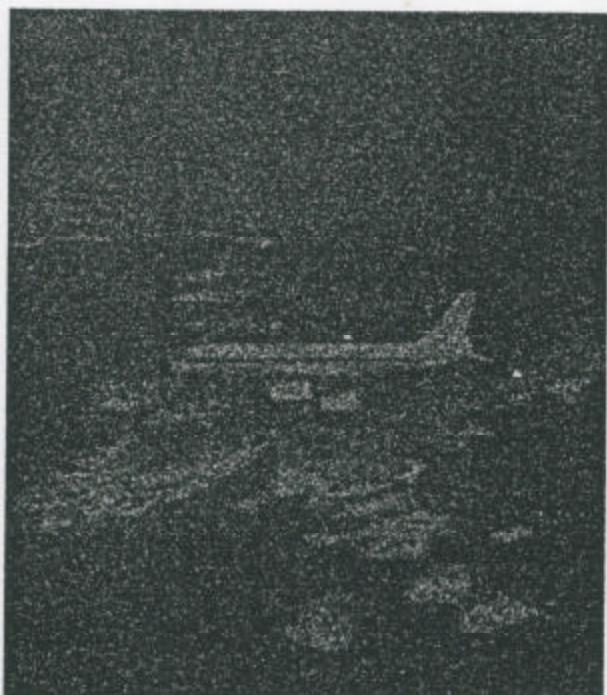


The Airplane Images

512*511 original
image (jpeg, NASA
AILS):

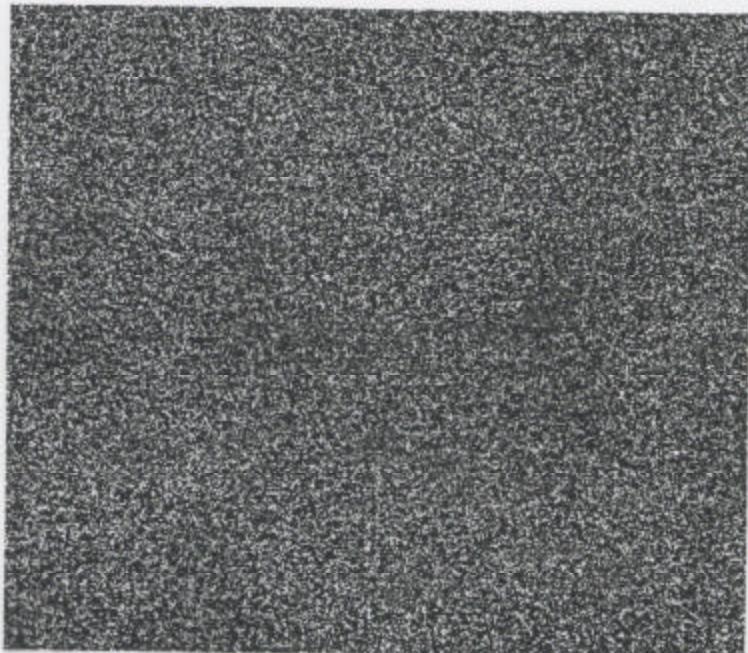


same image +
noise $\sim n(0,100)$:

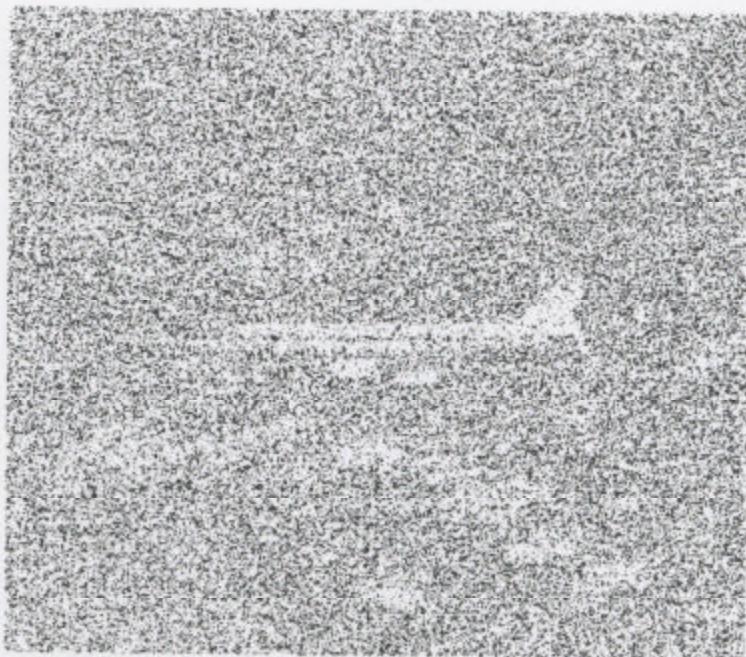


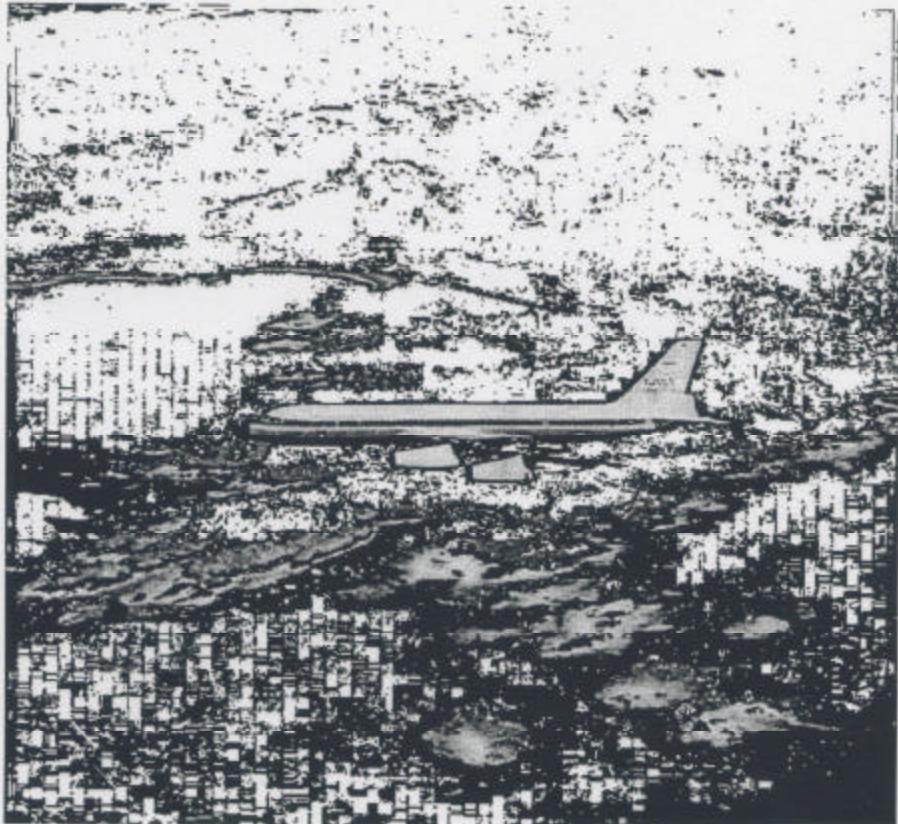
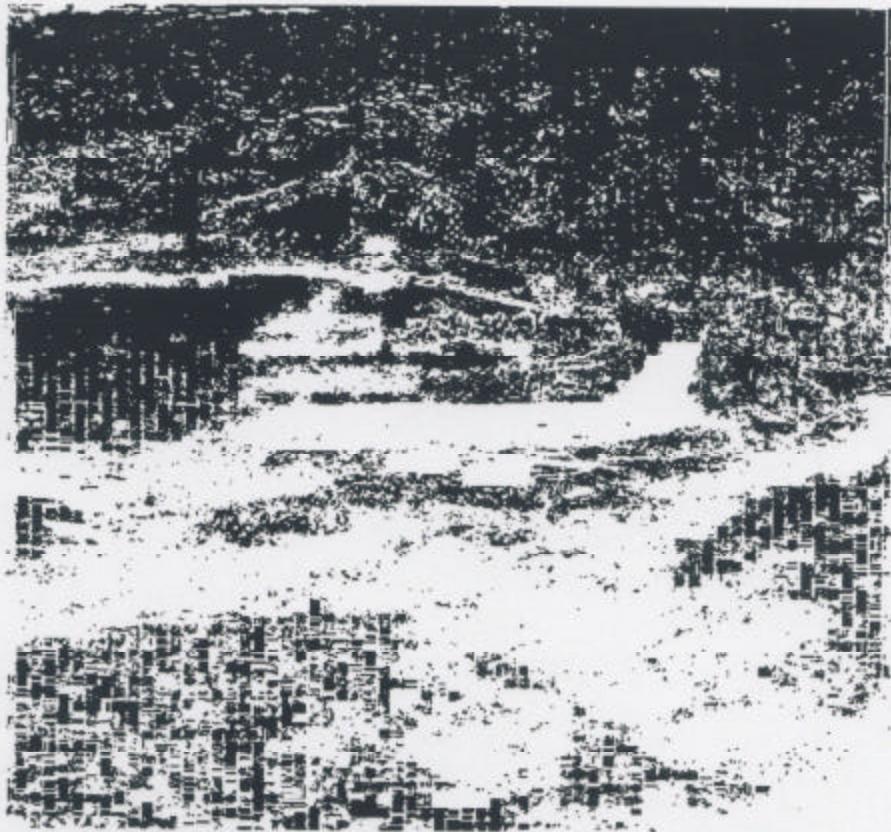
Noisy image partition:

K_0



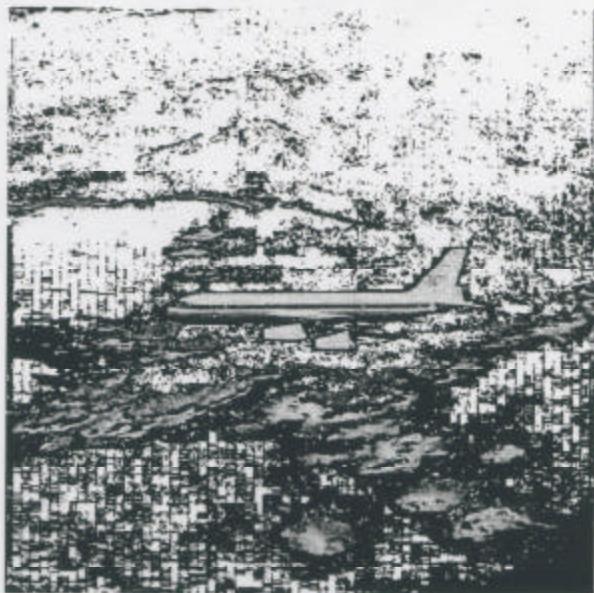
K_1



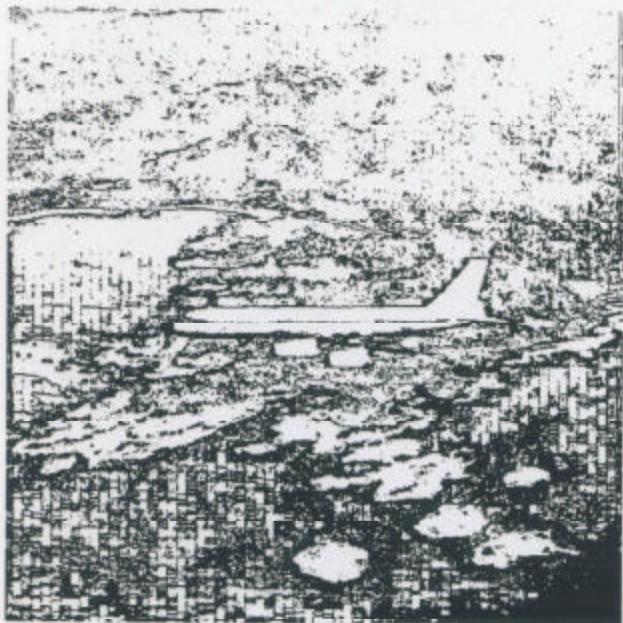
K_0  K_1 

80 11

K_0



K_{00}



K_{01}



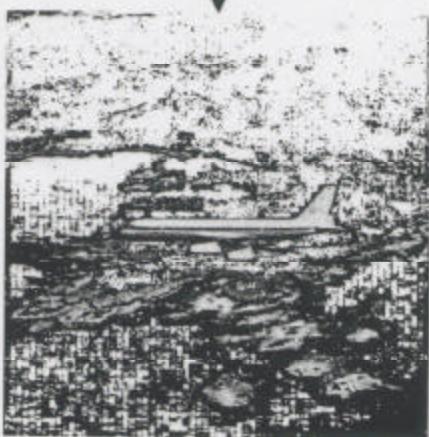
81

12

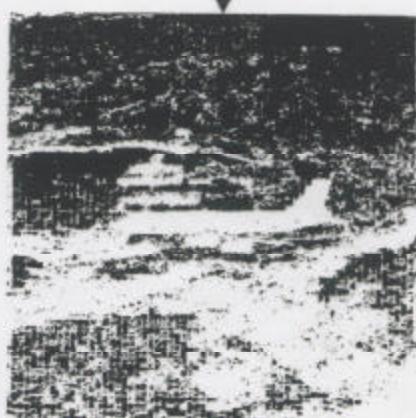
K



K_0



K_1



K_{00}



K_{01}

