Finding counterexamples via reinforcement learning

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- Agent tries to improve his total score (sum of all rewards) through some optimization algorithm.

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- Try to avoid using human insights as much as possible
- Would like a general setup: use the same program for every problem, only change reward function
- Throw this setup at 100 open conjectures and hope for the best.

Conjecture

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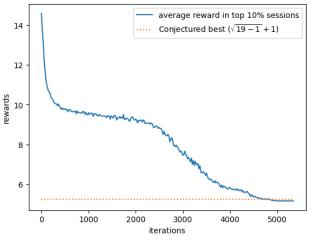
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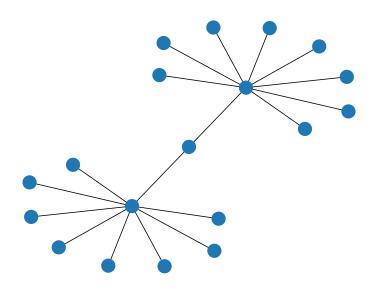
Reward: $\lambda_1 + \mu$ (minimize).

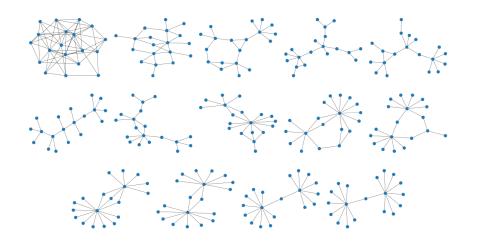
Run a reinforcement learning algorithm for n = 19:

Conjecture

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- This was the dream scenario: there is an obvious way to phrase the conjecture as a game, there is an obvious choice for the score function, we plug these into the RL program and it spits out a counterexample.
- When this happens, there is not much to talk about. But often it is not that simple.
- We will see 5 more examples. In each of them we will succeed in refuting an open conjecture, but each example will illustrate a unique thing that could "go wrong" and how to overcome it.

Example 2 – What if we don't succeed?

Conjecture (Auchiche-Hansen, 2016)

Let G be a connected graph with diameter D, proximity π and distance spectrum $\partial_1 \geq \ldots \geq \partial_n$. Then

$$\pi + \partial_{\left\lfloor \frac{2D}{3} \right\rfloor} > 0.$$

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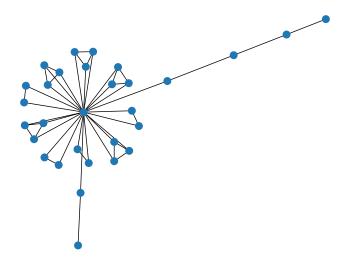
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Reward: $\pi + \partial_{\lfloor \frac{2D}{3} \rfloor}$ (minimize).

Run it for n = 30:



This is not quite a counterexample $(\pi + \partial_{\lfloor \frac{2D}{3} \rfloor} \approx 0.4)$, but it tells us very clearly what counterexamples could look like.

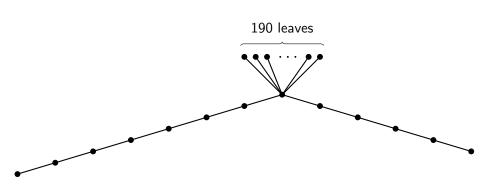


Figure: A counterexample to the conjecture

Example 3 - Not just graphs

Question (Brualdi-Cao)

How large can the permanent of a 312-pattern avoiding 0-1 matrix be?



Figure: The pattern 312

$$per(A) = \sum_{\sigma \in S_n} \prod_{i=1}^n A_{i,\sigma(i)}$$

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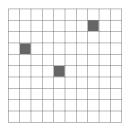
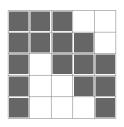


Figure: This is also not allowed

More precisely: we are not allowed to have three ones (dark squares) (x_i, y_i) : $i \in \{1, 2, 3\}$ such that $y_1 < y_2 < y_3$ and $x_2 < x_1 < x_3$.

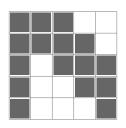
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The best you can do is $Fib_{n+2} - 1$.

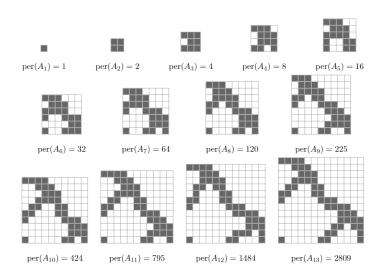


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Reward: per(A) - penalty(# of 312-s)



These are best possible for $n \le 8$ (computer proof). So the sequence starts with 1, 2, 4, 8, 16, 32, 64, 120.

Example 4 - Problems on trees

Conjecture (Collins, 1989)

Given a tree T, let p(T) and q(T) be the characteristic polynomials of the adjacency and the distance matrices of T, respectively. The coefficients of p and q are both unimodal, and their peaks are asymptotically at the same place.

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Reward: distance of the peaks.



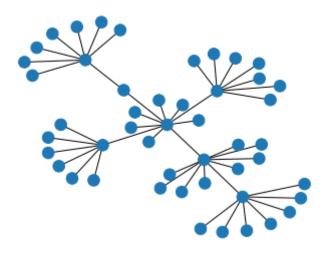


Figure: Best construction found for n = 48

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Say that a graph property \mathcal{P} is preserved under $\mathcal{D}^L(G)$ -cospectrality if $\operatorname{spec}_{\mathcal{D}^L}(G) = \operatorname{spec}_{\mathcal{D}^L}(H)$ implies $\mathcal{P}(G) = \mathcal{P}(H)$.

Property	\mathcal{D}^L	
# Edges	No	
Diameter	No	
Girth	No	
Planarity	No	
Wiener index	Yes	
Degree sequence	No	
Transmission sequence	No	
Transmission regularity	? <	
$\#$ connected components in \overline{G}	Yes	

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Is transmission regularity preserved under \mathcal{D}^L -cospectrality?

Task: find two graphs G and H such that they have the same \mathcal{D}^L -eigenvalues, but G is transmission regular and H is not.



A graph is transmission regular, if for each vertex, the sum of distances to all other vertices is the same. So if $\sum_w d(v,w) = \sum_w d(u,w)$ for all vertices u,v.

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Idea:

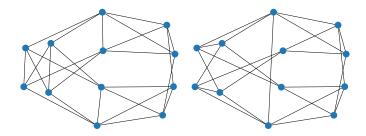
$$score(G, H) = f_1(G, H) + f_2(G) + f_3(H),$$

where

- f_1 measures how close the \mathcal{D}^L -spectrum of G and H is,
- \bullet f_2 measures how close G is to being transmission regular, and
- f_3 gives a penalty if H is transmission regular.

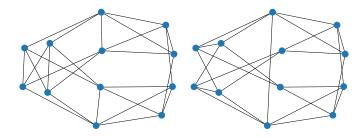


Example 5



The graph on the left is transmission regular, whereas the graph on the right is not. The characteristic polynomials of their distance Laplacians are the same ($x^{12}-216x^{11}+21188x^{10}-1245904x^9+48797440x^8-1336652544x^7+26129121472x^6-364516883456x^5+3556516628224x^4-23113129559040x^3+90045806284800x^2-159318669312000x$), so they are \mathcal{D}^L -cospectral.

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So transmission regularity is \mathbf{not} preserved under \mathcal{D}^L -cospectrality.

Many interesting problems can not have finite counterexamples.

Conjecture (Erdős, 1962)

The function

$$K_4(G)+K_4(\bar{G})$$

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Figure: Gwenaël Joret

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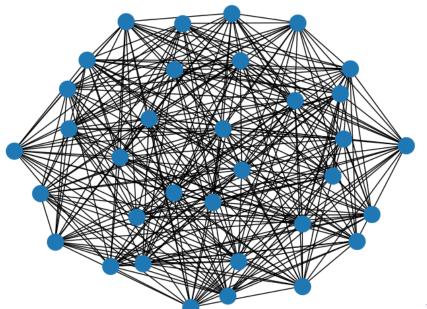
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Solution: "blowing up"! Construct a finite graph G, so that $G \times K_m$ is a counterexample as $m \to \infty$.

 $\lim_{m \to \infty} \frac{K_4(G \times K_m) + K_4(G \times K_m)}{m^4}$ depends only on G, and there is an easy formula for it. This will be our reward function.

Run RL for $n = 34 \longrightarrow \text{find a counterexample}$.

Example 6



Which RL algorithm to use?

- Value-based methods
 - Learn a value function: "I know how good this chess position is for black".
- Policy-based methods
 - Do not learn a value function: "I have no idea how good this chess position is for black, but I know the best move is c4".
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Task: Find the best algorithm for our problems!

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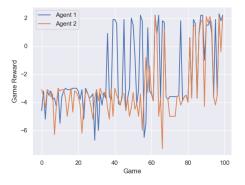
- Do we offer the edges one after the other, starting with $(1,2),(1,3),\ldots(n-1,n)$?
- Or do we generate a graph vertex by vertex?
- We could repeatedly ask the neural network to pick one edge to add, out of all remaining edges?
- Do we offer edges in random order? Do we offer multiple edges at a time, is it beneficial to offer an edge again after it was rejected once?
- What neural network architecture to use? Dense layers may not be the best (doesn't understand symmetry).

Reasons an RL algorithm might not work

- Sparse rewards problem: we give rewards only at the end of a game.
- Credit assignment problem: which of my $\frac{n(n-1)}{2}$ moves was responsible for getting a bad score?
- Bad reward design
- Explore-exploit dilemma

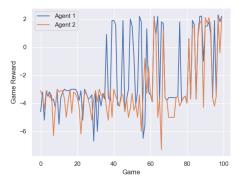
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Learning rate too high? Some other hyperparameter is wrong? Not enough training? Coincidence? Maybe this algorithm is unsuited for this specific problem, but good for others?

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- Simplest possible policy-based method.
- Fast convergence, very stable.
- This is the algorithm used in all the examples.

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Implementation details

- Offer edges one by one: (1,2), then (1,3),
- Input: adjacency matrix we have created so far, and a matrix indicating which edge we consider now.
- Each game lasts $\frac{n(n-1)}{2}$ steps (if we generate a graph).
- The input is two vectors of length n(n-1)/2 (or equivalent). The first contains 1-s for each edge we have taken, and 0-s for each edge rejected, or not considered yet. The second is all zeros, except for one place, corresponding to the next edge offered.
- Architecture: dense net, three layers of sizes 128, 64, 4.
- Learn from top 10%, but keep the top 5% for the next iteration.

Improvements to the method

Don't do a generalist approach, focus on one conjecture only, and find the best setup for this one problem!

- Pick an architecture that takes the symmetries of the problem into account (transformers, GNNs, canonicalizing the data)
- We might know that the counterexample must have a specific structure → use it to massively restrict the search space

Future goals:

- Improve or find a better RL algorithm for these purposes
- Refute conjectures with a different RL algorithms, or with human + RL collaboration
- Find other ways to use ML in mathematics

Thank you!