Material Flows in Production Networks

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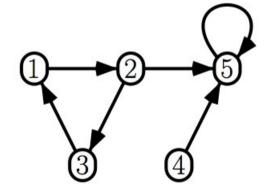
IPAM, Los Angeles, May 2009

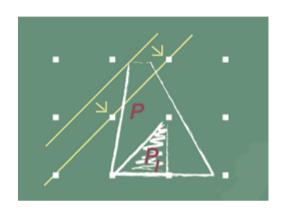
Continuous and Discrete models

Motivation and Application:

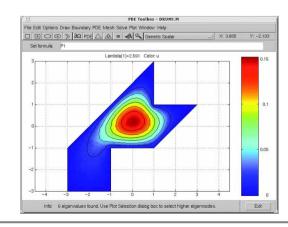
- Water Networks
- Gas Networks
- Traffic Flow Networks
- Supply Chains

- ...





MIP meets PDE





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Traffic Flow Networks

See Fügenschuh, Herty, Klar, Martin (2006): Combinatorial and Continuous Models for the Optimization of Traffic Flows on Networks



Minimize driving time subject to

$$\partial_t \rho_j(x,t) + \partial_x f_j(\rho_j(x,t)) = 0$$

$$f_j(\rho) = \rho u^e(\rho)$$

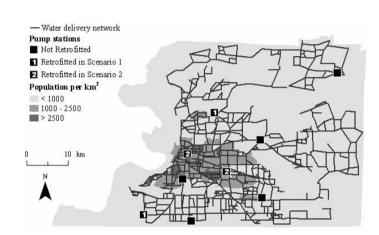
$$\sum_{j=1}^n f(\rho_j(b_j,t)) = \sum_{j=n+1}^{n+m} f(\rho_j(a_j,t))$$



Protection of Drinking Water Systems

See Fügenschuh, Göttlich, Herty (2007): Water Contamination Detection





Objective

Try to identify sources such that the time evolved concentrations coincide with given measurements

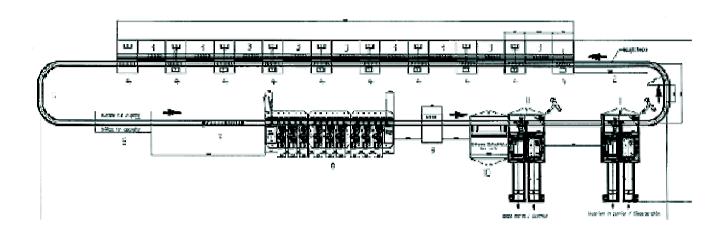
$$\sum_{j \in \mathcal{A}_{meas}} \max_{t \in (0,T)} (c^{j}(\bar{x}^{j},t) - \bar{c}^{j}(\bar{x}^{j},t)) + \rho \sum_{v \in \mathcal{V}} \max_{t \in (0,T)} q^{v}(t).$$

Subject to

Water network model (advection, decay, coupling)

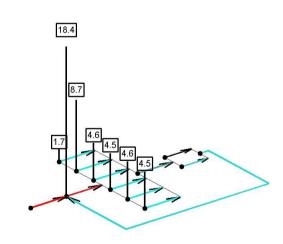


Production Systems



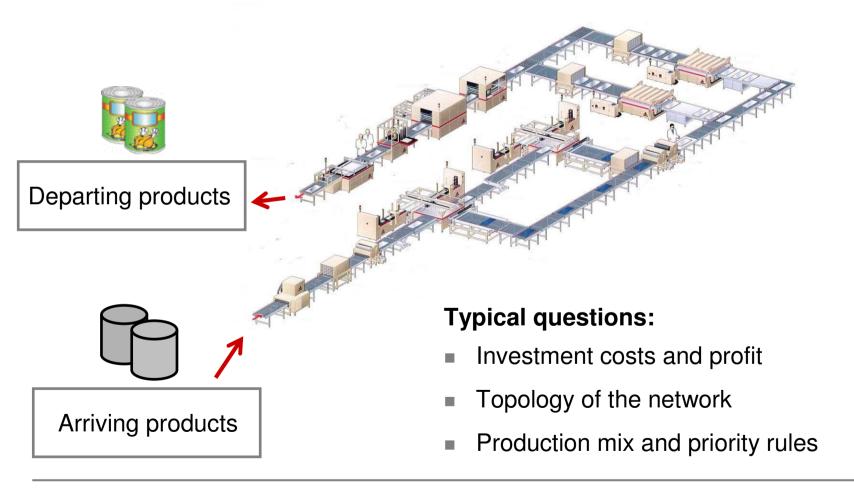
Maximize output subject to

$$\begin{aligned} \partial_t \rho^e(x,t) + \partial_x f^e(\rho^e(x,t)) &= 0 \\ \partial_t q^e(t) &= \alpha_v^e(t) \sum_{\bar{e} \in \delta_v^-} f^{\bar{e}}(\rho^{\bar{e}}(b^{\bar{e}},t)) - f^e(\rho^e(a^e,t)) \\ f^e(\rho^e(a^e,t)) &= \min\{\mu^e, \frac{q^e(t)}{\epsilon}\} \end{aligned}$$



Motivation

Planning of a Production Network





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Literature

System Dynamics

Forrester (1961): Industrial Dynamics

■ Baumol (1970): Economic Dynamics

Discrete Event Simulation

■ Banks et al. (1996): Discrete-Event System Simulation

■ Fishman (2001): Discrete-Event Simulation

Discrete
Optimization

■ Voß, Woodruff (2003): Introduction to Computational Optimization Models for Production Planning in a SC

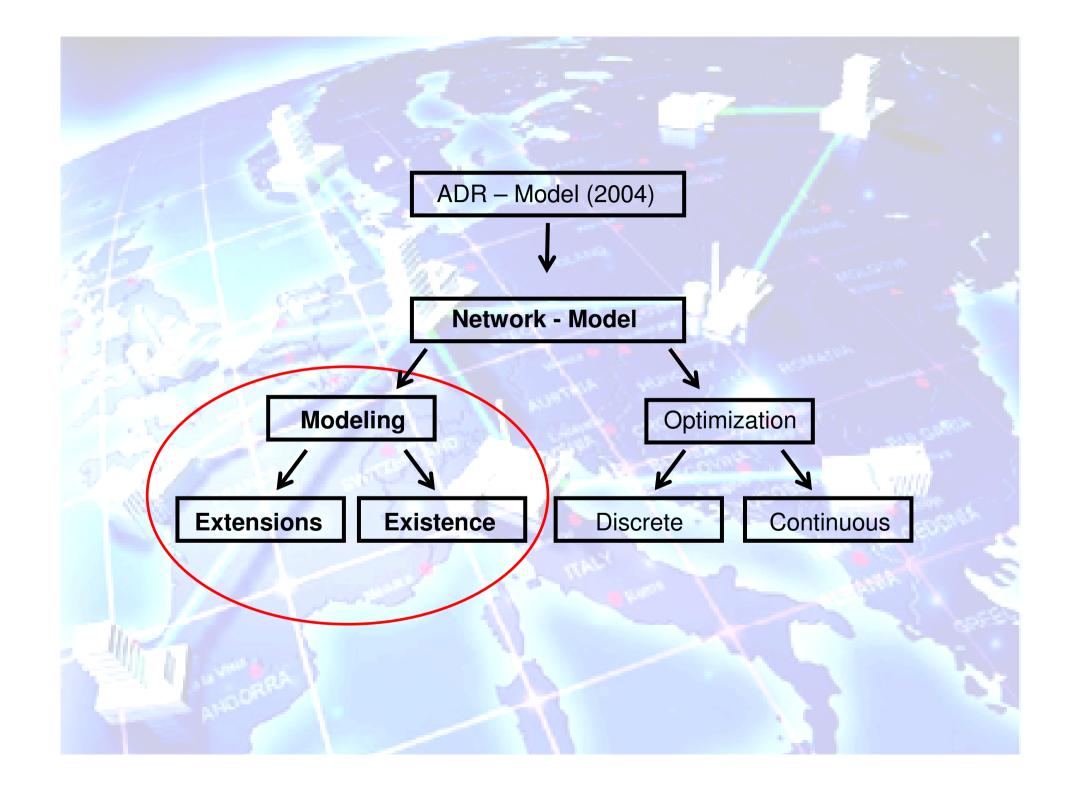
■ Wolsey, Pochet (2006): Production Planning by Mixed Integer Programming

Queueing Theory ■ Bolch et al. (1998): Queueing Networks and Markov Chains

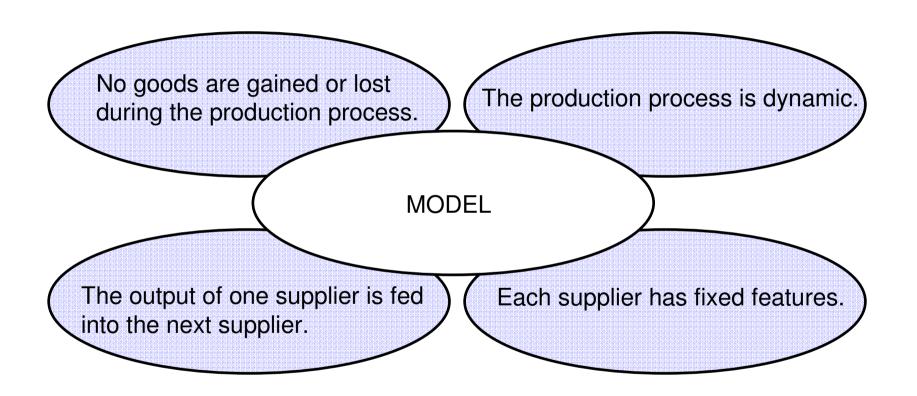
■ Chen, Yao (2001): Fundamentals of Queueing Networks



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Assumptions

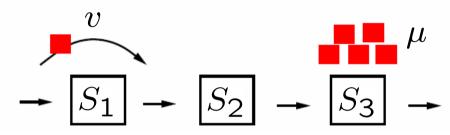


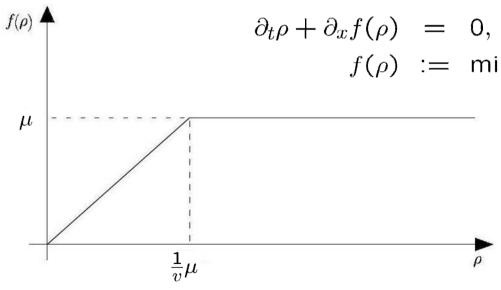
Continuous models are valid for large quantity production!



ADR - Model

See Armbruster, Degond and Ringhofer (2004): A model for the dynamics of large queuing networks and supply chains





 $\partial_t \rho + \partial_x f(\rho) = 0, \quad \forall x \in [a, b], \ t \ge 0$ $f(\rho) := \min \{v\rho, \mu\}$

 ρ : density of parts

maximum capacity

processing velocity

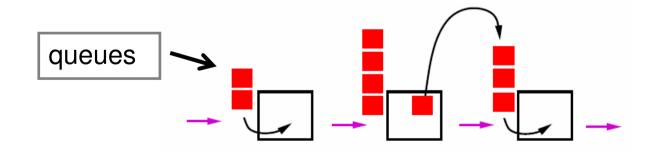
Network Model

Idea:

- Each processor is described by one arc
- Use ADR Model for dynamics inside the processor
- Add equations for queues in front of the processor

Advantage:

- Standard treatment of equations
- Straightforward definitions for complex networks





Network Model

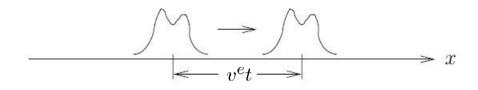
See Göttlich, Herty, Klar (2005): Network models for supply chains

Definition

A production network is a finite directed graph $(\mathcal{A}, \mathcal{V})$ where each arc $e \in \mathcal{A}$ corresponds to a processor on the intervall $[a^e, b^e]$. Each processor e has an associated queue q^e in front.

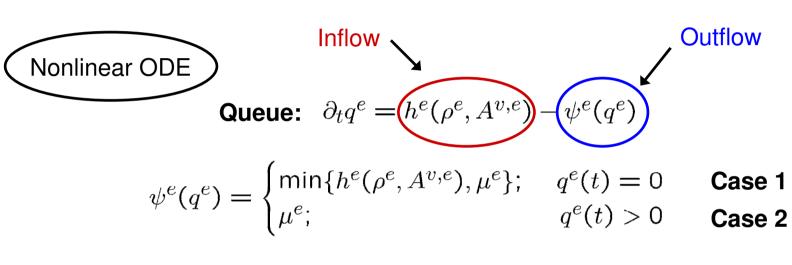
PDE: Linear transport equation

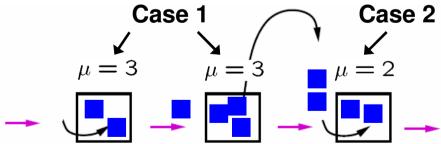
Processor: $\partial_t \rho^e(x,t) + v^e \partial_x \rho^e(x,t) = 0, \quad \forall x \in [a^e, b^e]$



Network Model

See Göttlich, Herty, Klar (2006): Modeling and optimization of supply chains on complex networks

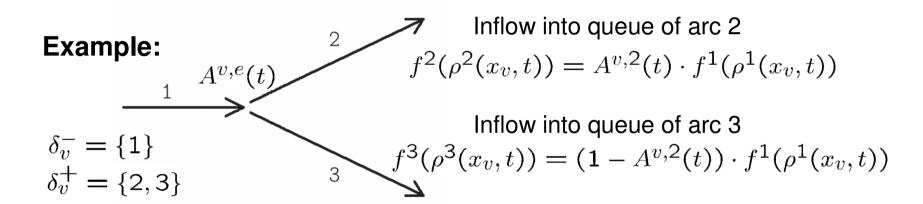




Network Coupling

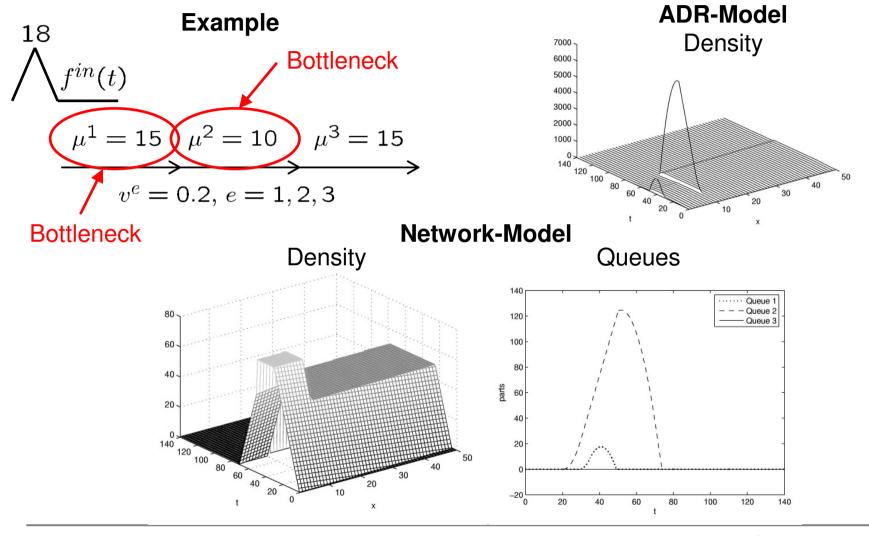
Definition

We define time-dependent distribution rates $A^{v,e}(t)$ for each vertex with multiple outgoing arcs. The functions $A^{v,e}(t)$ are required to satisfy $\sum_{e \in \delta_v^+} A^{v,e}(t) = 1$ and $A^{v,e}(t) \in [0,1]$.



 $A^{v,e}(t)$ will be obtained as solutions of the optimization problem

Numerical Results

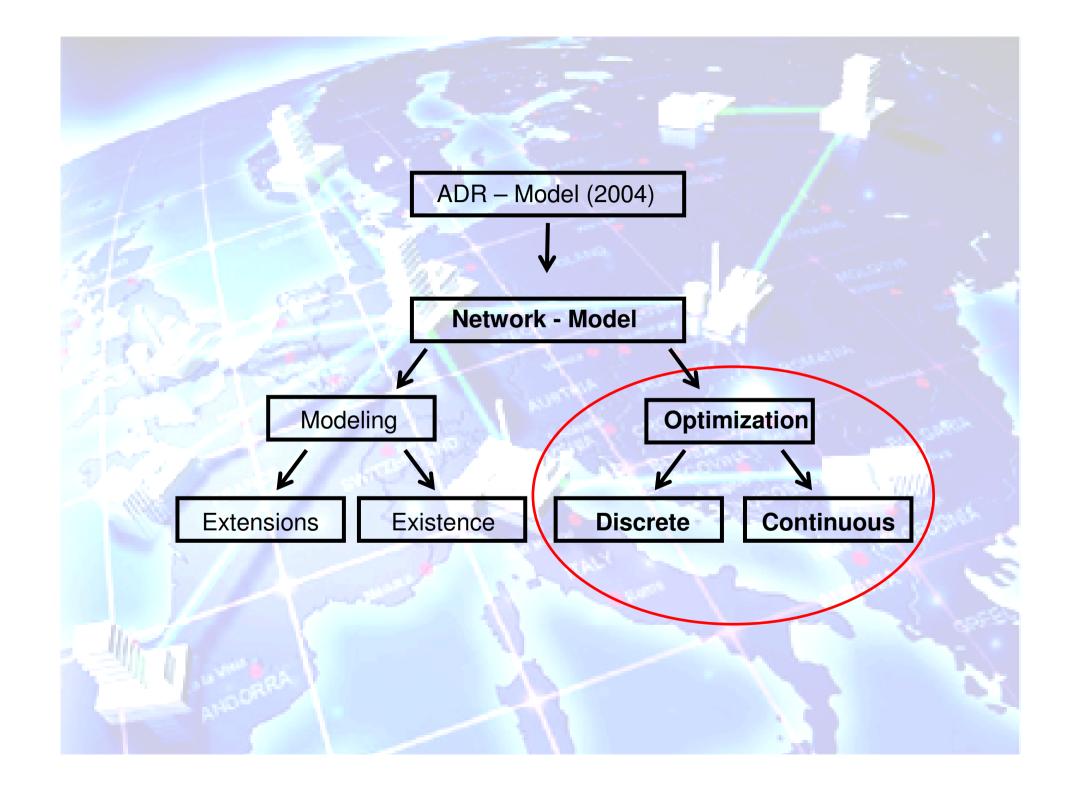




Continuous Production Models

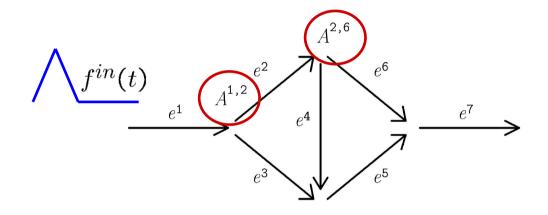
Pros	Cons	
 Accurate description of supply chain behavior 	 Difficult to include discrete decisions 	
Fast computing times	PDE-constrained optimization problems	
 Opportunity to introduce non-linearities 		





Optimization Problem

Consider network with entry/exit suppliers



- Given are time-dependent inflow profiles
- Controls are the distribution rates at vertices

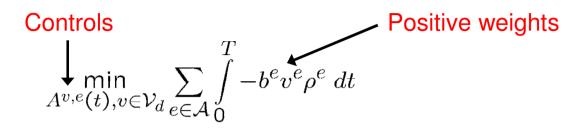
What is the optimal distribution of parts among the network?

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Optimization Problem

Cost functional



Constraints

subject to
$$e \in \mathcal{A}, v \in \mathcal{V}, t \in (0,T), x \in [a^e, b^e]$$

Processor
$$\longrightarrow \partial_t \rho^e + v^e \partial_x \rho^e = 0$$
,

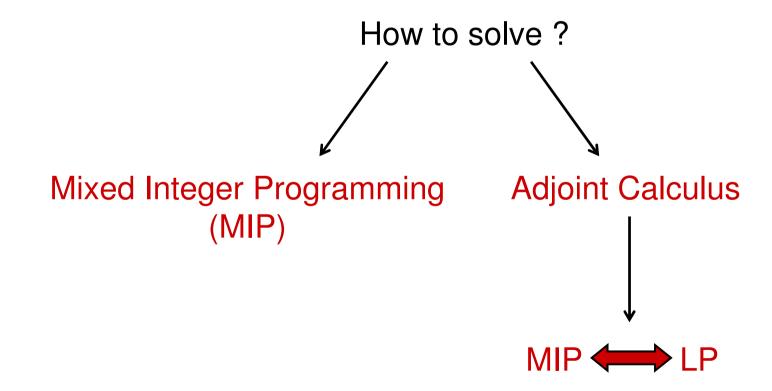
Outline
$$\partial_t q^e = h^e(\rho^e, A^{v,e}) - \psi^e(q^e)$$

$$\psi^e(q^e) = \begin{cases} \min\{h^e(\rho^e, A^{v,e}), \mu^e\}; & q^e(t) = 0\\ \mu^e; & q^e(t) > 0 \end{cases}$$

Initial conditions
$$\longrightarrow \rho^e(x,0) = 0, q^e(0) = 0$$

Solution Techniques

Optimal control problem with PDE/ODE as constraints!





Mixed Integer Program (MIP)

See Fügenschuh, Göttlich, Herty, Klar, Martin (2006): A Discrete Optimization Approach to Large Scale Networks based on PDEs

A mixed-integer program is the minimization/maximization of a **linear** function subject to linear constraints.

Problem: Modeling of the queue-outflow in a discete framework

$$\psi^{e}(q^{e}) = \begin{cases} \min\{h^{e}(\rho^{e}, A^{v, e}), \mu^{e}\}; & q^{e}(t) = 0\\ \mu^{e}; & q^{e}(t) > 0 \end{cases}$$

Solution: Reduce complexity (as less as possible binary variables)

Relaxed queue-outflow
$$\psi^e(q^e) = \min\left\{\frac{q^e(t)}{\epsilon}, \mu^e\right\}$$

¹See **Armbruster et al. (2006)**: Autonomous Control of Production Networks using a Pheromone Approach



Derivation MIP

Problem: Suitable discretization of the min-nonlinearity

$$\psi(q^e(t)) = \min\left\{\frac{q^e(t)}{\epsilon}, \mu^e\right\}$$

Idea: Introduce binary variables (decision variables) ξ_t^e

$$\psi(q_t^e) = \min\left\{\frac{q_t^e}{\epsilon}, \mu^e\right\} \qquad \longleftrightarrow \qquad \xi_t^e \in \{0,1\}, M >> 1$$

$$\mu^e \xi_t^e \leq \psi(q_t^e) \leq \mu^e$$
 Outflow of the queue
$$= \text{Inflow to a processor} \qquad \frac{q_t^e}{\epsilon} - M \xi_t^e \leq \psi(q_t^e) \leq \frac{q_t^e}{\epsilon}$$



Example: $\mu^e \leq \frac{q_t^e}{\epsilon}$ implies $\xi_t^e = 1$ and $\psi(q_t^e) = \mu^e$

The problem has $|\#Arcs| \cdot |\#Timesteps|$ binary variables!

Mixed Integer Program

Two-point Upwind discretization (PDE) and explicit Euler discretization (ODE) leads to

maximize outflux

$$\longrightarrow \min_{A_t^{v,e}, v \in \mathcal{V}_d} \sum_{e} \sum_{t} -\Delta t \, b_t \, v^e \rho_t^{e,b}$$
 subject to

processor

$$\longrightarrow \rho_{t+1}^{e,b} = \rho_t^{e,b} + \frac{\Delta t}{\Delta x} \left(\psi(q_t^e) - v^e \rho_t^{e,b} \right)$$

queue

$$q_{t+1}^e = q_t^e + \Delta t \left(h_t^e - \psi(q_t^e) \right)$$

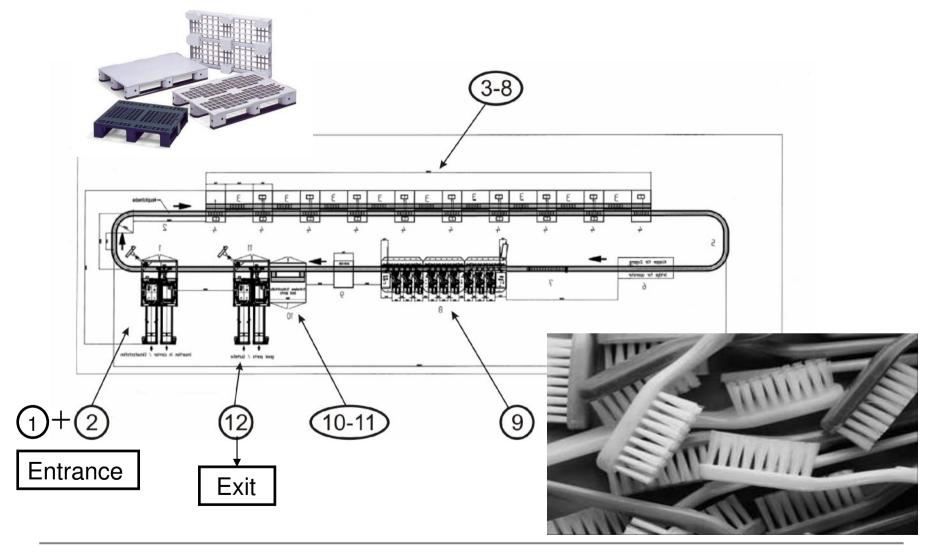
queue outflow

$$\frac{\mu^e \xi_t^e \le \psi(q_t^e) \le \mu^e}{\frac{q_t^e}{\epsilon} - M \xi_t^e \le \psi(q_t^e) \le \frac{q_t^e}{\epsilon}}, \qquad \xi_t^e \{0, 1\}$$

initial conditions

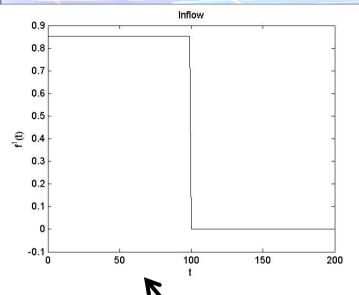
$$q_0^e = 0, \ \rho_0^{e,b} = 0, \ \psi(q_0^e) = 0$$

Toothbrush Manufacturing



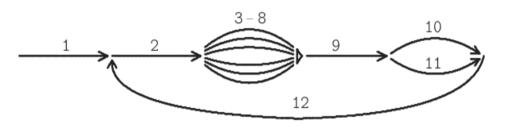


Example



Solved by ILOG CPLEX 10.0

- Solution time = 11.16 sec
- # Variables = 9600



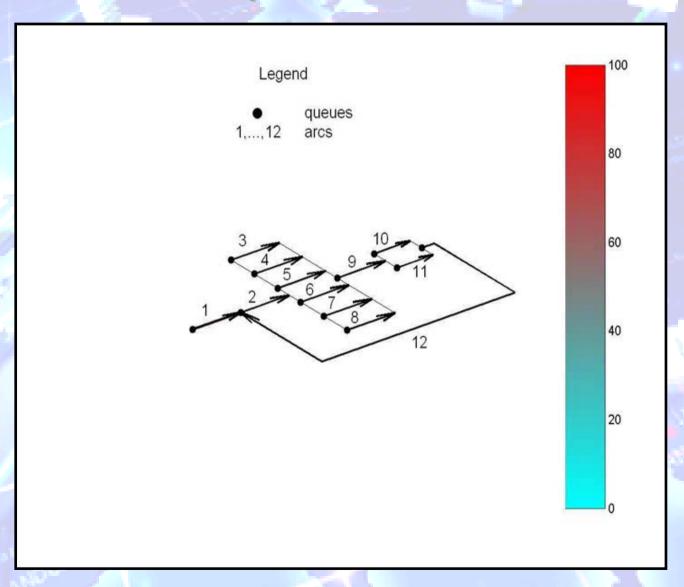
Constant inflow until $t \le 100$

Maximization of outflow, i.e. optimizing the amount of parts passing processor 12

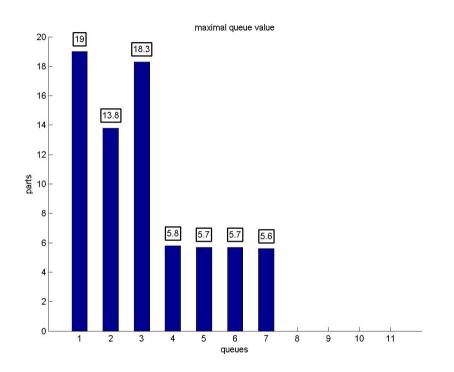
$$\min - \sum_{t} \frac{1}{t+1} g_t^{12}$$

e	μ^e	v^e
1	100	0.01
2	0.71	0.35
3 - 8	0.07	0.01
9	0.71	0.05
10 - 11	0.24	0.12
12	0.71	0.35

Optimization



Results



18
16
14
12
10
10
10
10
10
10
100
150
200

Maximal load of queues

Evolution of queues



Advantage MIP

Constraints can be easily added to the MIP model:

Bounded queues:

$$q_t^e \leq \text{const} \quad \forall e, t$$

Optimal inflow profile:

$$\max \sum_{e=1,t} \mathsf{f}_t^e \qquad \qquad \boxed{ \qquad \qquad } \mathsf{Optimal~\ref{pt}}$$

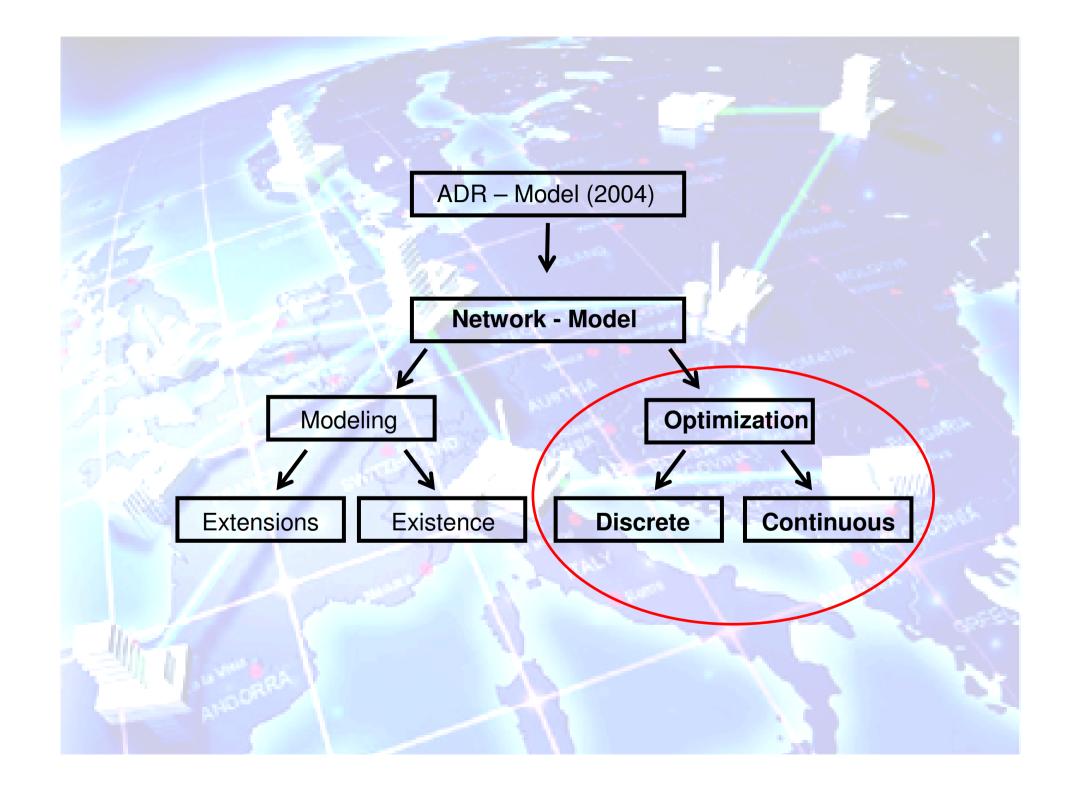
In other words: Find a maximum possible inflow to the network such that the queue-limits are not exceeded.

Advantage MIP

Maintenance shut-down:

$$\begin{split} \phi_t^{\tilde{e}} &\in \{0,1\}, \, \forall t, \, \forall l = 0, \dots, N-1 \\ \mathbf{h}_{t+l}^{\tilde{e}} &\leq \max\{\mu^e : e \in E\} |E| \cdot (1-\phi_t^{\tilde{e}}) \\ &\sum_{t=1}^{N_T} \phi_t^{\tilde{e}} = N \end{split}$$

■ Processor \tilde{e} has to be switched off for N consecutive time intervals during the total run time N_T .



Adjoint Calculus

See Göttlich, Herty, Kirchner, Klar (2006): Optimal Control for Continuous Supply Network Models

Adjoint calculus is used to solve PDE and ODE **constrained** optimization problems. Following steps have to be performed:

1. Define the **Lagrange** – functional:

$$L(\rho^e, A^v, q^e, \Lambda^e, P^e) = \sum_{e \in \mathcal{A}} \int_0^T \int_{a^e}^{b^e} -b^e v^e \rho^e dx dt \quad \text{Cost functional}$$

$$-\sum_{e \in \mathcal{A}} \int_0^T \int_{a^e}^{b^e} \Lambda^e \left[\partial_t \rho^e + v^e \partial_x \rho^e \right] dx dt \quad \text{PDE processor}$$

$$-\sum_{e \in \mathcal{A}} \int_0^T P^e \left[\partial_t q^e - h^e (\rho^e, A^v) + \psi^e (q^e) \right] dt \quad \text{ODE queue}$$

with Lagrange multipliers Λ^e and P^e .



Adjoint Calculus

2. Derive the first order optimality system (**KKT-system**):

Forward (state) equations:

$$\partial_t \rho^e + v^e \partial_x \rho^e = 0, \ \rho^e(x, 0) = 0, \ v^e \rho^e(a, t) = \psi^e(q^e),$$
$$\partial_t q^e = h^e(\rho^e, A^v) - \psi^e(q^e), \ q^e(0) = 0,$$

Backward (adjoint) equations:

$$-\partial_t \Lambda^e - v^e \partial_x \Lambda^e = v^e, \ \Lambda^e(x, T) = 0,$$

$$v^e \Lambda^e(b, t) = \sum_{\bar{e} \in \delta_v^+ \text{ s.t. } e \in \delta_v^-} P^{\bar{e}}(t) \frac{\partial}{\partial \rho^{\bar{e}}} h^{\bar{e}}(\rho^e, A^v),$$

$$-\partial_t P^e = 1 - (P^e - \Lambda^e(a, t)) (\psi^e)'(q^e), \ P^e(T) = 0,$$

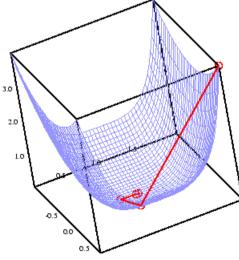
Gradient equation:

$$\sum_{e \in \delta_v^+} P^e \frac{\partial}{\partial A^{v,\overline{e}}} h^e(\rho^e, A^v) = 0.$$

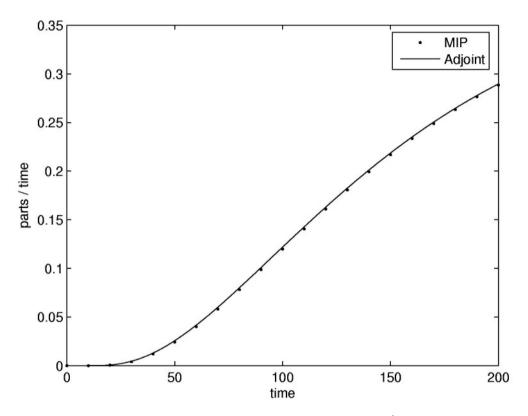
Optimization Algorithm

Projected steepest descent method:

- 1. Choose initial control vector A_0
- 2. Compute for A_0 the solution of state and adjoint equations
- 3. Compute the gradient. If it is zero, then STOP.
- 4. Update the control vector.
- 5. Go to 2.

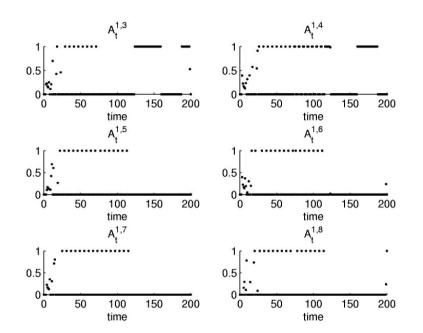


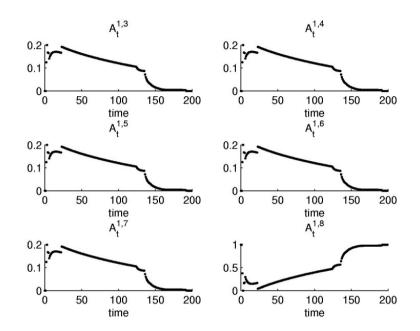
Optimal outflow profile



Optimal functional value: $J^*(\vec{A}^v) = -0.19$

Optimal control





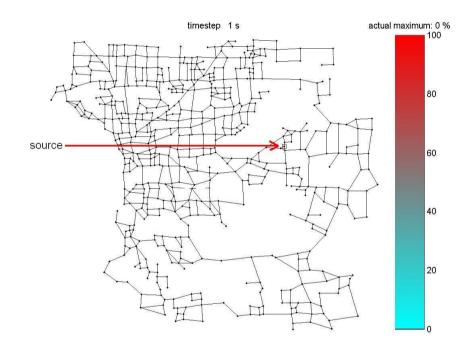
Distribution rates A_t^e , e = 3, ..., 8 computed by the MIP (left) and the adjoint approach (right)

→ No uniqueness!



Motivation

Optimizing large-scale networks with hundreds of arcs and vertices

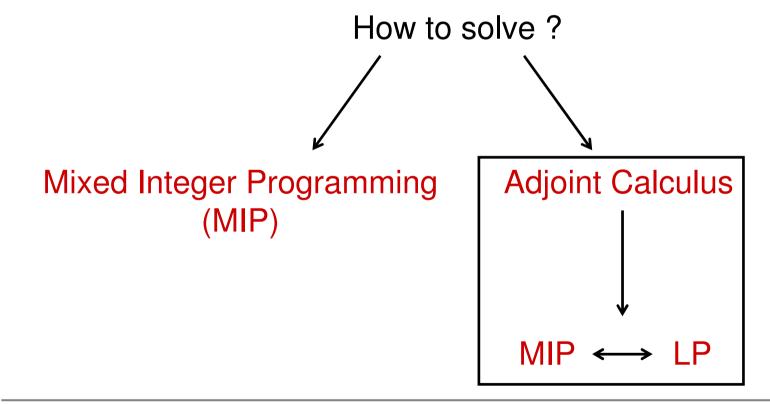


Solving the adjoint system or the MIP is computationally expensive!



Solution Techniques

Optimal control problem with PDE/ODE as constraints!





Linear Program (LP)

Idea: Use adjoint equations to prove the reformulation of the MIP as a LP



$$\psi(q_t^e) = \min\left\{rac{q_t^e}{\epsilon}, \mu^e
ight\}$$



Outflow of the queue



$$u_t^{e\pm} \ge 0$$

$$\psi(q_t^e) = \frac{1}{2}(\mu^e + \frac{q_t^e}{\epsilon}) - \frac{1}{2}(u_t^{e+} + u_t^{e-})$$

$$\mu^e - \frac{q_t^e}{\epsilon} = u_t^{e+} - u_t^{e-}$$

$$u_t^{e+} \cdot u_t^{e-} = 0$$

Remove the complementarity condition!

Remark: The remaining equations remain unchanged!

From MIP to LP

Theorem

Assume the inflow at a fixed vertex is non-zero and $b_{N_T}^e < 0$. Let either the costs b_t^e be monotone increasing in time or $\epsilon > \Delta t$. Then, every optimal solution to the LP automatically satisfies the complementarity condition $u_t^{e-} \cdot u_t^{e+} = 0$.

In other words:

We can solve the LP model and obtain the same solution as for the complementary problem!



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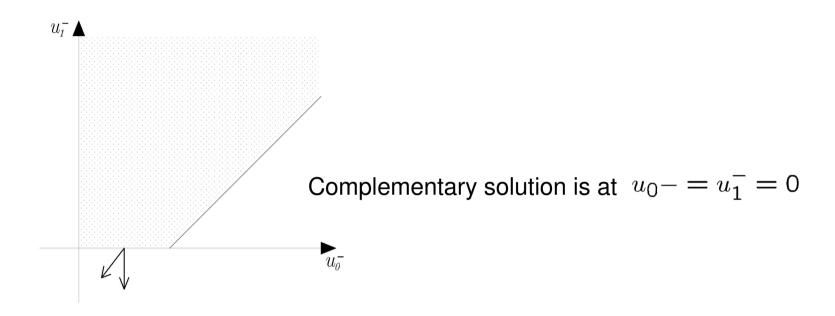
Idea of the Proof: Part I

Constraints can be reformulated as follows:

$$\begin{aligned} & \min_{A^{v,e}} \sum_{e} \sum_{t} -\Delta t b_{t}^{e} v^{e} \rho_{t}^{e,b} = \sum_{e} \sum_{t} -\Delta t b_{t}^{e} x_{t}^{e} + \dots \\ & q_{t+1}^{e} = \dots, \ \rho_{t+1}^{e,b} = \dots \\ & x_{t}^{e} + u_{t}^{e-} = \frac{q_{t}^{e}}{\epsilon} \\ & x_{t}^{e} + u_{t}^{e+} = \mu^{e} \\ & u_{t}^{\pm} \geq 0 \end{aligned}$$

- Complementary condition is satisfied whenever the inflow to the arc is maximized
- $\hbox{\bf Monotone increasing costs b^e_t imply maximizing x^e_t }$
- Maximizing x_t^e is minimizing either u_t^{e-} or u_t^{e+} \longrightarrow complementary condition is satisfied

Idea of the Proof: Part II



- Consider a single supplier with constant costs $b_t^e = -1$ and time horizon T=2
- Complementary solution is still an optimal solution
- But we obtain another solution if we store all incoming parts in the queue for one time-step and extract later --> complementary condition is not satisfied
- For $\epsilon = \Delta t$ the cost vector is orthogonal to the boundary

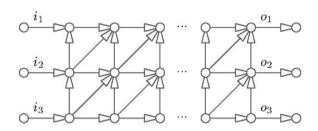


Numerical Results I

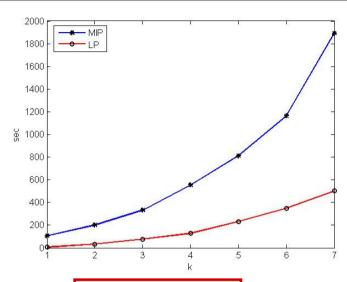
Example

Solved by ILOG CPLEX 10.0

$$\frac{f_{\{i_1,i_2,i_3\}}^{in}(t)}{\Box}$$



 $k \times 3$ interior network points



k	Presolve MIP	Solution MIP	Presolve LP	Solution LP
$\overline{1}$	78.87	105.33	0.11	6.26
2	159.09	201.64	0.32	33.54
3	236.59	332.78	0.54	76.91
4	352.08	555.11	0.75	127.24
5	477.50	810.44	0.99	232.68
6	590.09	1163.25	1.22	348.53
7	983.55	1891.90	1.42	502.72
8	907.01	infeasible infeasible	1.73	768.47

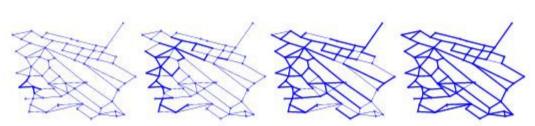
Numerical Results II

A production network consisting of 418 arcs and 233 vertices.

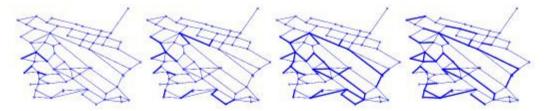


~15 Mio. variables!

Clipping area



Simulation results for t=10,23,36,50



Optimization results for t=10,23,36,50

MIP: no result after one day!

LP: ca. 5 hours



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