# Mean Field Equilibria of Dynamic Auctions with Learning

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Joint work with Krishnamurthy Iyer, Stanford University, and Mukund Sundararajan, Google Inc. Many marketplaces employ auctions to conduct trade:

Sponsored search

Google



Online markets





Licensing, etc.

Inspired by auction settings where agents *do not know* their valuation for an item a priori

Through repeated participation, agents **learn** their preferences for goods in the market

*Example*: In sponsored search advertising, an advertiser only learns the value of an ad based on conversion to a sale *after* a user clicks on the ad

### Sponsored search advertising



Major challenges:

- What strategies are "optimal" for bidders?
- Can we characterize market behavior, and in particular the distribution of bids in the market?
- What auction format should the market operator use?
- Should the market operator subsidize learning?

#### In this talk

- We use a mean field model to characterize agents' behavior in presence of learning.
- We establish existence and approximation theorems for MFE.
- 3 We use MFE to study market design: the impact of auction format and reserve prices on the auctioneer's revenue.

## Outline

- 1 Market model
- 2 Dynamic game
- 3 Mean field equilibrium
- 4 Approximation
- 5 Computation
- 6 The auction format
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Finite number of agents

Sequence of second price auctions:

•  $\alpha$  agents (sampled uniformly) per auction

Geometric lifetimes with parameter  $\beta$ :

- After an auction, participating agents leave independently with probability  $1 \beta$ .
- Each departing agent is replaced by a new agent.

In a (static) second price auction:

The highest bidder wins, and pays the second highest bid.

Exercise: It is a *dominant strategy* to bid your true valuation.

Agent *i*'s private valuation  $v_i \in [0, 1]$ : **unknown**, independent.

Valuation determines the reward  $x_{i,t}$ :

$$x_{i,t} = egin{cases} 1 & ext{with probability } v_i; \ 0 & ext{otherwise.} \end{cases}$$

e.g., in sponsored search: reward = sale after a click-through

Observing reward  $x_{i,t}$  informs an agent about her valuation  $v_i$ .

### Initial prior: Beta(m, n)

Density: 
$$f_{(m,n)}(x) \propto x^{m-1}(1-x)^{n-1}$$

Mean: 
$$\mu(m,n) = \frac{m}{m+n}$$

Variance:  $\sigma^2(m, n)$  decreasing in *m* and *n* 

On losing the auction:



On winning the auction, and getting positive rewards:



On winning the auction, and getting zero rewards:



 $s_k$  = Belief parameters after  $k^{th}$  auction

Belief update: after  $k^{th}$  auction,

$$s_k = \begin{cases} s_{k-1} & ext{if the agent does not win;} \\ s_{k-1} + e_1 & ext{if the agent wins and } x_k = 1; \\ s_{k-1} + e_2 & ext{if the agent wins and } x_k = 0, \end{cases}$$

where  $e_1 = (1, 0)$  and  $e_2 = (0, 1)$ .



### Maximize the total expected payoff over the lifetime



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(Per period payoff = reward - payment)

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If private valuation is known:

No learning  $\implies$  auctions decouple

 $\implies$  Bid truthfully each time period

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#### $\implies$ Bid truthfully each time period

What if private valuation is unknown?

If private valuation is unknown:

Value for learning  $\implies$  agents overbid

But how much?

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Hard to determine

Standard tool to analyze dynamic games with *incomplete information* is *perfect Bayesian equilibrium* (PBE).

PBE assumes agents are completely rational:

- Agents track each competitor as long as they stay in the market, and play optimally.
- Agents maintain consistent beliefs about evolution of the entire market, and update them using Bayes' rule.



PBE suffers from two issues:

Intractability: often showing equilibrium exists is difficult.
No structural insight

Implausibility: Not a good model of agent behavior in practice

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#### Inspired by large markets

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Weintraub et al. ('08), Lasry & Lions ('07), Huang et al. ('07), ...
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Agents do not track individual competitors

Each agent plays against a "stationary" market

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Sponsored search

Advertisers use **bid landscape** information to model the rest of the market.

# Mean field equilibrium

### **Optimality:**

Stationary market



Actions are optimal

**Consistency:** 







#### Mean Field Equilibrium = Optimality + Consistency



#### OPTIMALITY

# MFE: Stationary market

#### Suppose the distribution of bids in the market is g



For a fixed agent, in each of her auctions, bids of other  $\alpha - 1$  agents are sampled i.i.d. from *g*.

# MFE: Stationary market

#### Suppose the distribution of bids in the market is g



Probability of winning:  $q(b|g) = g(b)^{\alpha-1}$ 

Expected payment: p(b|g)

Expected payoff in  $k^{th}$  auction:  $q(b|g)\mu(s_k) - p(b|g)$ .

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Geometric lifetime
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Geometric lifetime

## Expected total discounted payoff maximization

From standard dynamic programming, optimal strategy is

- Markovian: bid depends only on current belief
- Stationary: no time dependence

### CONSISTENCY

## MFE: Consistency



# MFE: Consistency



For consistency, we need g to be a **fixed point** of F.

### Definition

A bid distribution g and a strategy C constitute an MFE if

- **1** Optimality: Given *g*, the strategy *C* is optimal.
- 2 Consistency: g is a fixed point of map F.

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Does an MFE exist?

## Existence of MFE

### Theorem

A mean field equilibrium exists in the auction market.

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Proof uses an infinite dimensional fixed point theorem.

- Show: With the **right** topologies, *F* is continuous.
- Show: Image of *F* is compact

# Existence of MFE

### Theorem

A mean field equilibrium exists in the auction market.

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- Show: With the **right** topologies, *F* is continuous.
- Show: Image of *F* is compact

Existence:

- General belief models
- Multiple units per auction

### **OPTIMAL STRATEGY IN MFE**

$$V(s|g) = \max_{b \ge 0} \left\{ q(b|g)\mu(s) - p(b|g) + \beta q(b|g)\mu(s)V(s + e_1|g) + \beta q(b|g)(1 - \mu(s))V(s + e_2|g) + \beta (1 - q(b|g))V(s|g) \right\}$$

$$V(s|g) = \max_{b \ge 0} \left\{ \underbrace{q(b|g)\mu(s) - p(b|g)}_{(1)} + \beta q(b|g)\mu(s)V(s + e_1|g) + \beta q(b|g)(1 - \mu(s))V(s + e_2|g) + \beta(1 - q(b|g))V(s|g) \right\}$$

(1) Expected payoff in current auction

$$\begin{split} V(s|g) &= \max_{b \ge 0} \left\{ q(b|g)\mu(s) - p(b|g) + \underbrace{\beta q(b|g)\mu(s)V(s+e_1|g)}_{(2)} \right. \\ &+ \beta q(b|g)(1-\mu(s))V(s+e_2|g) + \beta (1-q(b|g))V(s|g) \right\} \end{split}$$

### (2) Future expected payoff on



$$V(s|g) = \max_{b \ge 0} \left\{ q(b|g)\mu(s) - p(b|g) + \beta q(b|g)\mu(s)V(s + e_1|g) + \underbrace{\beta q(b|g)(1 - \mu(s))V(s + e_2|g)}_{(3)} + \beta (1 - q(b|g))V(s|g) \right\}$$

(3) Future expected payoff on



$$V(s|g) = \max_{b \ge 0} \left\{ q(b|g)\mu(s) - p(b|g) + \beta q(b|g)\mu(s)V(s + e_1|g) + \beta q(b|g)(1 - \mu(s))V(s + e_2|g) + \underbrace{\beta(1 - q(b|g))V(s|g)}_{(4)} \right\}$$

### (4) Future expected payoff on



$$V(s|g) = \max_{b \ge 0} \left\{ q(b|g)\mu(s) - p(b|g) + \beta q(b|g)\mu(s)V(s + e_1|g) + \beta q(b|g)(1 - \mu(s))V(s + e_2|g) + \beta (1 - q(b|g))V(s|g) \right\}$$

# MFE: Agent's decision problem

Given g, agent's value function satisfies Bellman's equation:

$$V(s|g) = \max_{b \ge 0} \left\{ q(b|g)\mu(s) - p(b|g) + \beta q(b|g)\mu(s)V(s + e_1|g) + \beta q(b|g)(1 - \mu(s))V(s + e_2|g) + \beta(1 - q(b|g))V(s|g) \right\}$$

## MFE: Agent's decision problem

## Rewriting:

$$V(s|g) = \max_{b \ge 0} \left\{ q(b|g)C(s|g) - p(b|g) \right\} + \beta V(s|g),$$

### where

$$C(s|g) = \mu(s) + \beta \mu(s)V(s + e_1|g) + \beta(1 - \mu(s))V(s + e_2|g) - \beta V(s|g).$$

# MFE: Optimality

Agent's decision problem is

$$\max_{b\geq 0} \Big\{q(b|g)C(s|g) - p(b|g)\Big\}$$

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### Same decision problem as in

- Static second-price auction
- **a**gainst  $\alpha 1$  bidders drawn i.i.d. from g
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- **a**gainst  $\alpha 1$  bidders drawn i.i.d. from g
- with agent's known valuation C(s|g).

We show  $C(s|g) \ge 0$  for all s

 $\implies$  Bidding C(s|g) at posterior s is optimal!

## C(s|g): Conjoint valuation at posterior s

$$C(s|g) = \mu(s) + \beta \mu(s)V(s + e_1|g) + \beta(1 - \mu(s))V(s + e_2|g) - \beta V(s|g)$$

### C(s|g): Conjoint valuation at posterior s

 $C(s|g) = \mu(s) + \beta\mu(s)V(s + e_1|g) + \beta(1 - \mu(s))V(s + e_2|g) - \beta V(s|g)$ 

### Conjoint valuation = Mean + Overbid

(We show Overbid  $\geq 0$ )

### Overbid: $\beta \mu(s)V(s+e_1|g) + \beta(1-\mu(s))V(s+e_2|g) - \beta V(s|g)$

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#### Overbid

Expected marginal future gain from **one additional observation** about private valuation

If private valuation is unknown:

Value for learning  $\implies$  agents overbid

But how much?

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Expected marginal value of one additional observation

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Simple description of agent behavior!

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### Does an MFE capture rational agent behavior in finite market?

Issues:

- Repeated interactions ⇒ agents no longer independent.
- Keeping track of history will be beneficial.

Hope for approximation only in the asymptotic regime

Look at the market as an interacting particle system.

Interaction set of an agent: all agents influenced by or that had an influence on the given agent. Look at the market as an interacting particle system.

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Intuition: As market size increases, any two agents' interaction sets become disjoint with high probability.

# Approximation

#### Theorem

As the number of agents in the market **increases**, the maximum additional payoff on a **unilateral** deviation converges to zero.

As the market size increases,

Expected payoff under optimal strategy, given others play  $C(\cdot|g)$  Expected payoff under  $\mathbf{C}(\cdot|\mathbf{g})$ , given others play  $C(\cdot|\mathbf{g})$ 

 $\rightarrow$  0

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#### Theorem

As the number of agents in the market **increases**, the maximum additional payoff on a **unilateral** deviation converges to zero.

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Mean field equilibrium is **good** approximation to agent behavior in finite large market.

Good approximation even in small markets due to **behavioral** reasons

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A natural heuristic inspired by model predictive control.

Implicitly encodes a learning algorithm for the agents

Closely models market evolution when agents optimize given current average estimates.



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# Algorithm

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# Algorithm

- **1** Initiate the market at bid distribution  $g_0$ .
- **2** Given  $g_k$ , compute conjoint valuation  $C(\cdot|g_k)$ .
- 3 Evolve the market *one time period*, assuming each agent bids her conjoint valuation.
- 4 Compute the new bid distribution  $g_{k+1}$ .
- 5 Repeat until  $||g_{k+1} g_k||_{\infty} < \epsilon$ .

# Heuristic converges to MFE within 30-50 iterations in practice, for reasonable error bounds ( $\epsilon\sim 0.0015)$

Computation takes  $\sim$  30 mins on a laptop.

## Overbidding



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## Repeated standard auctions

Standard auction:

- 1 Winner has the highest bid.
- 2 Zero bid implies zero payment.

Example: First price, all pay, third price, etc.

#### Theorem

*MFE* exists in any repeated **standard auction** if the static auction has a symmetric Bayes-Nash equilibrium.

#### Dynamic revenue equivalence

Expected revenue to the auctioneer is **same** irrespective of static auction format

(This is an analog of a similar result in *static* auction theory.)

*Moral:* Auctioneer's expected revenue not affected by choice of auction format.

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Setting a reserve can increase auctioneer's revenue.

Effects of a reserve:

- Relinquishes revenue from agents with low conjoint valuation.
- 2 Extracts more revenue from those with high conjoint valuation.

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Effects of a reserve:

- Relinquishes revenue from agents with low conjoint valuation.
- 2 Extracts more revenue from those with high conjoint valuation.
- 3 Imposes a learning cost:
  - Precludes agents from learning, and reduces incentives to learn.

Due to learning cost, agents change behavior on setting a reserve.

Auctioneer sets a reserve r and agents behave as in an MFE with reserve r.

Defines a **game** between the auctioneer and the agents.

# **Optimal reserve**

Two approaches:

- **Nash equilibrium**: Ignores learning cost. Auctioneer sets reserve assuming bid distribution is fixed, and agents behave as in MFE with reserve *r*.
- Stackelberg equilibrium: Includes learning cost. Auctioneer computes revenue in MFE for each *r*, and sets the maximizer *r*<sub>OPT</sub>.

We compare these two approaches using numerical computation.

By definition,  $\Pi(r_{OPT}) \ge \Pi(r_{NASH})$ .

 $\Pi(r_{OPT}) - \Pi(0)$  is greater than  $\Pi(r_{NASH}) - \Pi(0)$  by  $\sim 1 - 30\%$ .

Ignoring learning incurs a potentially significant cost.

Improvement depends on the distribution of initial beliefs of arriving agents.

## CONCLUSION

The methodology of MFE allows for

- Tractability: many analytical insights possible
- Plausibility: conjoint valuation captures nicely the value of learning

Numerical computation feasible  $\implies$  questions of practical relevance such as optimal reserve can be answered through computation.

Proof of convergence of the heuristic

How does efficiency of such repeated one-shot mechanisms compare with optimal mechanisms?

Empirical validation

Other models:

- Unit demand bidders (eBay, Amazon, etc.)
- Budget constrained bidders

## THANK YOU

Paper at http://ssrn.com/abstract=1799085