Noise, complexity, and information dynamics in quantum circuits





TS, Yao PRL 131, 160402 (2023) TS, Yao forthcoming (2023) Thomas Schuster IPAM, November 2023

Many-body quantum science

Coherent large-scale quantum systems may achieve large advantages in computation, physics, chemistry, sensing, cryptography, ...





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However, near-term quantum devices are impacted by noise



Quantum error correction can address this, but requires additional capabilities/overhead

What can we do with noisy quantum devices?

Physics experiments

 $|\psi_{qsL}\rangle = \left| \begin{array}{c} & & \\ &$ +

Harvard/MIT (2022)

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NISQ algorithms

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Harvard/MIT (2022)

NISQ algorithms



Benchmarking noise



Marcus group (2015)

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Cerezo et al (2021)

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Error mitigation









Gate-by-gate: benchmark individual gates, hope they perform the same when together

Sampling: measure bitstrings, compare with classical simulation



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Loschmidt echo: thought expt \sim 1870's, NMR \sim 1980's, single-particle theory \sim 2000's,

NISQ expts ~2010s Pines, Pastawski, Peres, Jalabert, Zurek, Cappellaro, Sanchez,...

















Experiments show the Loschmidt echo decay depends on U itself

How can we understand this in many-body quantum systems?



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• Recent results show that random noisy circuits are classically simulable (in scaling with *n*) Yung, Gao (2017), Gao, Duan (2018), Barak et al (2020), Gao et al (2021), Fontana et al (2023), Aharonov et al (2023)



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Does this extend to more general classes of quantum circuits?





Noise \Leftrightarrow **Information dynamics** \Leftrightarrow **Complexity**

Operator growth in open quantum systems TS, Yao PRL 131, 160402 (2023)



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A polynomial-time classical algorithm for almost any noisy quantum circuit

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Quantum information dynamics

Operator growth: How does local information evolve in time?



e.g. $U = e^{-iHt}$

Linked to recent developments in quantum gravity, quantum chaos, quantum state transfer, tensor network algorithms, quantum thermalization...

Zooming in... Operator size distributions

Roberts et al. (2018), Streicher, Qi (2019)

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Define the **size** (a.k.a. "weight") of a Pauli string = its number of non-identity elements

$$\hat{R} = \mathbbm{1} \otimes \mathbbm{1} \otimes X \otimes \mathbbm{1} \otimes Y \otimes Z \otimes X \otimes X \otimes \mathbbm{1} \otimes Y$$
 Size = 6

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Time-evolved operators have a **size distribution**:

$$\hat{Q}(t) = \sum_{R} c_R \hat{R}$$



Roberts et al. (2018), Streicher, Qi (2019)






Intuition: noise propagates according to the circuit's operator growth dynamics



Intuition: highly non-local operators are more sensitive to noise



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 $\hat{M}_t = X + Y \otimes \mathbb{1} \otimes Z + \ldots + Y \otimes Z \otimes \mathbb{1} \otimes X \otimes Z$

Intuition: highly non-local operators are more sensitive to noise





More precise: Noise decays Pauli strings at a rate proportional to their size $\sim \gamma \mathcal{S}$





(exact for single-qubit decoherence;

otherwise, approximation for high-size components under ergodic dynamics)



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$$\partial_t \mathcal{N} = -\gamma \overline{\mathcal{S}} \mathcal{N}$$



equal to Loschmidt echo!



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$$\partial_t \overline{\mathcal{S}} = (\text{unitary}) - \gamma \delta \mathcal{S}^2$$



equal to average OTOCs!



Interplay between operator growth and noise determines the Loschmidt echo

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Noisy: Size *plateaus* to constant value!

$$\partial_t \overline{\mathcal{S}} \approx \lambda \overline{\mathcal{S}} - \gamma \delta \mathcal{S}^2$$

 $\delta \mathcal{S} \sim \mathcal{S}$
 $\mathcal{S}_{\mathrm{sat}} = rac{\lambda}{\gamma}$





Example: All-to-all dynamics

Unitary: Size grows exponentially in time. Distribution is broad.

Noisy: Size *plateaus* to constant value!



 \Rightarrow Loschmidt echo decay is *independent* of noise rate!

Echoes seminal results in single-particle quantum chaos Jalabert, Pastawski (2001)



Recent NMR experiments

NMR expts w/ \sim 1000 spins

Domínguez et al. (2021)



All-to-all random circuits

TS, Yao, arXiv:2208.12272 (2022)





Adamantane lattice 192 nearest neighbors



All-to-all random circuit

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We provide an efficient classical algorithm for "almost any" noisy quantum circuit

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i.e. pick initial bitstring at random, calculate O for bitstring



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Our result: A classical algorithm to compute expectation values with root-meansquare error $\varepsilon \cdot \|O\|_F$ in time

$$n^{\mathcal{O}\left(\frac{1}{\gamma}\log\left(\sqrt{T}/\varepsilon\right)\right)}$$



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Idea: Simulate in Heisenberg picture, keep only low-size components

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Sensitivity to noise \Leftrightarrow Complexity

Decompose operator in Pauli string basis at each time step

Close correlation between "paths" of Pauli strings that are hard to simulate, and those that are strongly damped by noise [Aharonov, Gao, Landau, Liu, Vazirani (2023)]





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To bound contribution from high weights, previous works used that in random circuits, Pauli paths on average do not coherently interfere

Builds upon recent algorithms for noisy random circuit sampling, but with modifications in the algorithm + proof techniques [Aharonov, Gao, Landau, Liu, Vazirani (2023)]

Algorithm: At each time step, truncate any component of O(t) with size above a threshold ℓ



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Truncate *T* times, each with error $e^{-\gamma \ell}$:

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Simple extensions

Corollary 1: Can replace ensemble of input states with a single *highly mixed* state:

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Corollary 2: Ensemble of input states \rightarrow ensemble of circuits with *spatial disorder*



"Push" randomness from initial state into circuit

Encompasses quantum simulation of disordered spin models

Our bounds are weak for noise rates on leading quantum devices

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But designing NISQ circuits that take advantage of these noise budgets is challenging!

see e.g. White, Refael, Pollman, Rakovsky, von Keyserlingk, Ye, Machado, Nahum, Zhou... also IBM (2023) + Kedezchi et al, Anand et al, Begusic et al, Rudolph et al, ...

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Classically simulable Insensitive to noise Potentially complex Sensitive to noise
Implications for NISQ experiments?

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Corollary 3: Complex quantum circuits **must be** highly sensitive to noise:

Any quantum experiment* with classical complexity $\chi(n, \epsilon)$, must only succeed for

$$\gamma \leq \mathcal{O}\left(\frac{\log(n)\log\left(\sqrt{T^*}/\varepsilon\right)}{\log(\chi(n,\varepsilon))}\right)$$

*For computing expectation values with small rms error

e.g. if
$$\chi \sim \exp(n)$$
, require $\gamma \lesssim 1/n$

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Quantum information dynamics provides a natural language for understanding noise and complexity in NISQ experiments

Future Directions

- Benchmarking Loschmidt echo predictions with NMR experiments; connection to noise-induced transitions Google Quantum AI (2023), Ware et al (2023)
- "Experimental tests" for efficient classical simulations
- Dephasing noise, related e.g. to Clifford + T gate simulations
- What is the sensitivity to noise of physical many-body dynamics? Can we identify classes of dynamics that are / are not sensitive? White, Refael, Pollman, Rakovsky, von Keyserlingk, Ye, Machado, Nahum, Zhou...

TS, Yao PRL 131, 160402 (2023) TS, Yao forthcoming (2023)

Norman Yao

