## Modular quantum signal processing

 with gadgets
## [2309.16665]

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(a)

$\square$ interlink
$\longrightarrow$ correction


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Quantum algorithms: tough to design and interpret

Despite beautiful work in QPLs [1, 2], we still reason at gate level, most of the time

What can quantum programming techniques \& abstractions learn from their classical cousins?

## Quantum algorithm design



## Overview

Pinpoint strengths of algorithmic class:

Argue function-first algorithmic design with above is (1) desirable but (2) not immediately possible

Turn a desire into a formal problem

Show solution utility, efficiency, functional
programming flavor, and introduce package

## QSP: functional interpretability

Circuit map:
$\Phi \in \mathbb{R}^{n+1}, s \in\{0,1\}^{n} \mapsto \mathrm{q}$. gates.
$-\underbrace{}_{s_{k}=0,1}=A(a)^{s_{k}} B(b)^{1-s_{k}}$


Space- and query-efficient algorithms for spectral mapping; unifies most quantum algorithms [3, 4]

## The temptation

QSP and QSVT achieve* polynomial transforms

$$
\langle 0| U(\Phi, x)|0\rangle=P(x)
$$

Given protocols achieving $f(x), g(y)$, and $h(z, w)$, it feels like we should be able to achieve

$$
h \circ(f, g) \equiv h(f(x), g(y)) .
$$

Moreover, the composite protocol ought to be able to use $f(x)$ and $g(y)$ as contiguous subroutines


## The problem

QSP encodes privileged basis: expects $X$ rotations as input but does not produce them as output

Even worse lifting to QSVT or with many oracles

QSP protocols aren't first class functions

## Why bother

Modularity and reusability in coherent setting
QSP is great: space efficient, infinity norm dependent, strong numerical footing ...

Multivariable analogues of QSP are hard, while multivariable functions (comparators, control-flow, loss) are ubiquitous [5]

## A note on QSVT

Commuting operators and spectral mapping Anything we do in SU(2) is done in preserved QSVT subspaces. Let $\Pi, \tilde{\Pi}$ proj. and $A=\tilde{\Pi} U \Pi$, then

$$
\begin{aligned}
& A=\sum_{i=1}^{d_{\min }} \xi_{i}\left|\tilde{\psi}_{i}\right\rangle\left\langle\psi_{i}\right| \mapsto P^{S V}(A) \equiv \sum_{i=1}^{d_{\min }} P\left(\xi_{i}\right)\left|\tilde{\psi}_{i}\right\rangle\left\langle\psi_{i}\right|, \\
& U=\cdots \oplus \bigoplus_{\xi_{i} \neq 0,1}\left[\begin{array}{cc}
\xi_{i} & \sqrt{1-\xi_{i}^{2}} \\
\sqrt{1-\xi_{i}^{2}} & -\xi_{i}
\end{array}\right]_{\tilde{\mathcal{H}}_{i}}^{\mathcal{H}_{i}} \oplus \cdots
\end{aligned}
$$

## Gadgets 气

An $(a, b)$ gadget is a unitary ${ }^{*}$ superoperator:

$$
\bigotimes_{k \in[a]} e^{i \theta_{k} \sigma_{x}} \mapsto \bigotimes_{j \in[b]} e^{i \phi_{j} \sigma_{z}} e^{i \psi_{j} \sigma_{x}} e^{-i \phi_{j} \sigma_{z}}
$$

over $\theta$ in known range. Each of $b$ output legs achieves $\cos \psi_{j}$, often polynomial in $\cos \theta_{k}, k \in[a]$.

Functionally interpretable boxes: a inputs, b outputs

## Gadget Rosetta Stone

(a) Self-composition of circuits

(b) Self-composition of polynomials

(c) Self-composition of $(2,2)$ gadgets (n.b. superoperators)


## Thm 1: linking gadgets

Let $\varepsilon, \delta>0, \mathfrak{G}$ an $(a, b)$ gadget and $(\Xi, S)$ an atomic ( $c, d)$ gadget, where $\Xi \equiv\left\{\Phi_{0}, \ldots, \Phi_{d-1}\right\}$ and $S \equiv\left\{s_{0}, \ldots, \boldsymbol{s}_{d-1}\right)$. Suppose $\mathfrak{G}$ achieves

$$
F(x) \equiv\left\{f_{0}\left(x_{0}, \cdots, x_{a-1}\right), f_{1}\left(x_{0}, \cdots, x_{a-1}\right), \ldots, f_{b-1}\left(x_{0}, \cdots, x_{a-1}\right)\right\}
$$

over $x \in[-1,1]^{\times a}$, and $(\Xi, S)$ achieves

$$
G(y) \equiv\left\{g_{0}\left(y_{0}, \cdots, y_{c-1}\right), g_{1}\left(y_{0}, \cdots, y_{c-1}\right), \ldots, g_{d-1}\left(y_{0}, \cdots, y_{c-1}\right)\right\}
$$

over $y \in[-1,1]^{\times c}$. Let $\mathfrak{I}=(B, C, W)$ an interlink. Then, there exists a gadget $\mathfrak{G}^{\prime}$ which $\varepsilon$-approximately achieves
$H\left(x, y^{\prime}\right) \equiv \bigcup_{k \in[d]} g_{k}\left(\bigcup_{j \in B} f_{W(j)}\left(x_{0}, \ldots, x_{a-1}\right) \cup \bigcup_{k \notin C} y_{k}\right) \cup \bigcup_{k \notin B} f_{k}\left(x_{0}, \ldots, x_{a-1}\right)$
over $\left(F(x), y^{\prime}\right) \in \mathcal{D}$, where $y^{\prime}$ is the subset of $y_{k}$ such that $k \notin C$ and $\mathcal{D}$ is a domain determined by the correction protocol used.
$\mathfrak{G}^{\prime}$ uses description of $(\Xi, S)$ and $\widetilde{\mathcal{O}}\left(d|\Xi|_{\infty} \zeta\right)$ black-box calls to result of running $\mathfrak{G}$.
(a)


- interlink
$\longrightarrow$ correction
(b)

(c)


Generally, $\zeta=\widetilde{O}\left(\operatorname{polylog}\left(\varepsilon^{-1}\right) \operatorname{poly}\left(\delta^{-1}\right)\right)$ for $\delta=\mathcal{O}(1)$. Note $|\Xi|_{\infty}$ is max length of elem in $\Xi$.

## Thm 2: gadget-poly equivalence $1 \pm$

Let $\mathcal{L}\left(x_{1}, \ldots, x_{n}\right)$ polynomials achievable by atomic gadgets over $x_{1}, \ldots, x_{n}$. Suppose $P\left(x_{1}, \ldots, x_{n}\right)$ has degree $D$ and can be split into a tower of $m=\mathcal{O}(\log (D))$ interlinked polynomials,

$$
P=P^{(m-1)} \circ_{\mathfrak{I}_{m-2}}\left(P^{(m-2)} \circ_{\mathfrak{J}_{m-3}} \circ\left(\cdots\left(P^{(1)} \circ_{\mathfrak{I}_{0}} P^{(0)}\right) \cdots\right)\right)
$$

such that $P_{i}^{(j)} \in \mathcal{L}\left(x_{1}, \ldots, x_{n}\right)$ for all $i, j$. If the $P_{k}$ are separated from $\{0, \pm 1\}$ by $\delta \in \mathcal{O}(1)$ for all $k$ over $\mathcal{D}_{k}$. Then, there exists an assemblage of $m$ atomic, snappable gadgets which $\varepsilon$-approximately achieves $P$ on $\mathcal{D}_{k}$ with aggregate query cost $\widetilde{O}(\operatorname{poly}(D)$ polylog $(\varepsilon))$.

| Algebraic structure | Manipulation | Algebraic form |  | Reference |
| :---: | :---: | :---: | :---: | :---: |
| QSP/M-QSP | Polynomial application <br> Multivariable variant ${ }^{\dagger}$ | $\begin{aligned} & x_{0} \mapsto \\ & x_{0}, \ldots \end{aligned}$ | $\rightarrow P\left(x_{0}, \ldots, x_{n}\right),$ | $\begin{aligned} & \text { Thm. D. } 1 \\ & \text { Thm. D. } 2 \end{aligned}$ |
| Polynomial ring ${ }^{\ddagger}$ | Addition <br> Multiplication <br> Scalar multiplication | $\begin{aligned} & x_{0}, x_{1} \\ & x_{0}, x_{1} \\ & x_{0} \mapsto \mathrm{n} \end{aligned}$ | $\begin{aligned} & \left.{ }_{0}+x_{1}\right) / 2, \\ & x_{1}, \\ & \\ & \left.x_{0}, 1\right), \alpha \in \mathbb{R}, \end{aligned}$ | Thm. IV. 3 <br> Thm. IV. 2 <br> Ex. IV.4/5 |
| Composition monoid | Polynomial composition | $f\left(x_{0}\right)$ | $\rightarrow(f \circ g)\left(x_{0}\right)$, | Thm. III. 1 |
| General gadgets | Gadget linking | $\mathfrak{G}, \mathfrak{G}^{\prime}$ |  | Thm. III. 1 |
| Provided examples | Negation/inversion (Ex. IV.1) <br> Angle sum/difference (Ex. IV.2) <br> Affine shift (Ex. IV.7) <br> Step function (Ex. IV.8) <br> General mean (Ex. IV.10) |  | Bandpass function (Ex. IV.11) <br> Majority vote (Ex. IV.12) <br> Functional interpolation (Ex. IV.13) <br> Chebyshev inverse (Thm. IV.1) |  |

## Sum gadget



A: square root gadget, B, C: angle sum/difference gadgets, D: product gadgets.

## Prior work

Cascaded classical filters for sharpening. [6, 7]
Recursively defined quantum (computational) subroutines. $[8,9,10,11,12,13,14]$

Single-variable self-embedded QSP/QSVT. [15, 16]
Multivariable QSP/QET and LCU-based methods.
Functional programming techniques: quantum and classical. [1, 2, 17]

## Correction protocol

Main idea: repeated use of twisted signal to obliviously, approximately cancel its own twist
$e^{i \phi \sigma_{z}} e^{i \psi \sigma_{x}} e^{-i \phi \sigma_{z}} \mapsto\left\{e^{i \phi \sigma_{z}} e^{i \psi \sigma_{x}} e^{-i \phi \sigma_{z}}, e^{2 i \phi \sigma_{z}}\right\} \mapsto e^{i \psi \sigma_{x}}$.
Off diagonal element's phase sent to $\approx \pi / 2$

$$
\left[\begin{array}{cc}
P & i Q \sqrt{1-x^{2}} \\
* & *
\end{array}\right] \mapsto\left[\begin{array}{cc}
P & i \sqrt{1-P^{2}} \\
* & *
\end{array}\right]
$$



## Computing costs - -

For $P(x)$ close to $\pm 1$, QSP suffers 'gimbal lock'
Take $\delta=\mathcal{O}(1)$, and $P(x) \in[-1+\delta, 1-\delta]^{b}$ over $x \in[-1+\delta, 1-\delta]^{\text {a }}$ : simple bounding box

Cost per correction is $\mathcal{O}\left(\operatorname{poly}\left(\delta^{-1}\right) \operatorname{polylog}\left(\varepsilon^{-1}\right)\right)$

## On query/space complexity $\square$

An (a, b) gadget has a cost for each output leg, up to desired precision, in queries to each input leg:
$a \times b$ cost matrix $C$
Up to padding rows \& cols, C of gadget computed by multiplying $C_{k}$ of sub-gadgets.

Space/query complexity breaks into cases: single-variable, psd $Q$, controlled access.

## Is this efficient?

Poly-many, poly-size gadgets, log depth.
Maintains QSP's infinity norm scaling.
Poly-logarithmic in inverse precision.
No post-selection required.
Hierarchical, distributed, modular.
Linking high-degree gadgets is where method shines.
Potentially huge space saving (psd $Q$ ); coupled to questions in poly decomposition theorems.

| Method | Queries | Space | Norm |
| :--- | :--- | :--- | :--- |
| $[18]$ | poly $(r) d^{\text {poly }(r)}$ | $\operatorname{poly}(r)[$ poly $(d)+\operatorname{poly}(s)]$ | $\\|\cdot\\|_{1}$ |
| Ours | $d$ | $s+c$ | $\\|\cdot\\|_{\infty}$ |

Table: Asymptotic query and space complexity, as well as relevant norm. Here $d$ is generalized degree, $r$ is number of variables, and $s$ is qubits needed for block encodings. Note the case that $c=0$ is well-understood.

## Functional programming

## Semantics and syntax

Monad: monoid in the category of endofunctors

$$
\begin{aligned}
& \text { unit: } T \rightarrow M T \\
& \text { bind }:(M T, T \rightarrow M U) \rightarrow M U
\end{aligned}
$$

Unwrapped type is scalar value (or $X$ rotation), wrapped is QSP output; this work constructs bind

Bonus: attribute grammar \& language of gadgets

## Python package

| Q gadget_init.py UNREGISTERED |  |
| :---: | :---: |
| 1 | xi_ $0=[[0,0,0,0],[0,0,0,0]]$ |
| 2 | xi_1 $=[[0,0,0,0],[0,0,0,0]]$ |
| 3 |  |
| 4 | s_0 $=[[0,1,0],[1,0,1]]$ |
| 5 | s_1 $=[[0,1,0],[1,0,1]]$ |
| 6 |  |
| 7 | \# Create two (2, 2) gadgets named "g0" and "g1". |
| 8 | g0 = AtomicGadget (2, 2, "g0", xi_0, s_0) |
| 9 | g1 = AtomicGadget (2, 2, "g1", xi_0, s_1) |
| 10 |  |
| 11 | \# Link first output of g0 to first input of g1. |
| 12 | linking_guide = [(("g0", 0), ("g1", 0))] |
| 13 | a0 = g0.wrap_gadget() |
| 14 | a1 = g1.wrap_gadget() |
| 15 | a2 = a0.link_assemblage(a1, linking guide) |
| $\square 1$ | 15, Column 43 Tab Size: 4 |

github.com/ichuang/pyqsp/tree/beta


## To take home

Efficient, function-first quantum algorithm design
Coherent use of QSP protocols as QSP oracles with functional interpretability

Readymade repo at ichuang/pyqsp/tree/beta
Formal syntax and semantics, coupled to deep lit in (Q)PLs: compilation, verification [1, 19, 20]

Ancilla-free gadgets: minimal functional description?
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