# Quantum algorithmic tools for simulating open quantum systems

András Gilyén Alfréd Rényi Institute of Mathematics

joint work with

Anthony (Chi-Fang) Chen, Michael Kastoryano, and Fernando Brandão

CQCWS1, IPAM, UCLA, CA, US 4th October 2023



▶ We have a (small) quantum system that is coupled to a (large) environment

- ► We have a (small) quantum system that is coupled to a (large) environment
- Early proposal: simulate both system and environment on a quantum computer

- ▶ We have a (small) quantum system that is coupled to a (large) environment
- Early proposal: simulate both system and environment on a quantum computer
  - See Terhal, DiVincenzo '98

- ▶ We have a (small) quantum system that is coupled to a (large) environment
- Early proposal: simulate both system and environment on a quantum computer
  - ► See Terhal, DiVincenzo '98
  - Very resource intensive: many qubits are devoted to environment simulation

- ▶ We have a (small) quantum system that is coupled to a (large) environment
- Early proposal: simulate both system and environment on a quantum computer
  - See Terhal, DiVincenzo '98
  - Very resource intensive: many qubits are devoted to environment simulation
- We would like to only simulate what happens with system qubits

- ▶ We have a (small) quantum system that is coupled to a (large) environment
- Early proposal: simulate both system and environment on a quantum computer
  - ► See Terhal, DiVincenzo '98
  - Very resource intensive: many qubits are devoted to environment simulation
- We would like to only simulate what happens with system qubits
- Idea describe effective dynamics using Lindbladian Master equation:

- ▶ We have a (small) quantum system that is coupled to a (large) environment
- Early proposal: simulate both system and environment on a quantum computer
  - See Terhal, DiVincenzo '98
  - Very resource intensive: many qubits are devoted to environment simulation
- We would like to only simulate what happens with system qubits
- Idea describe effective dynamics using Lindbladian Master equation:

$$\frac{d\rho}{dt} = \underbrace{-i[H,\rho]}_{\text{(Schrödinger equation)}} + \underbrace{\sum_{j=0}^{m} \underbrace{L_{j}[\rho]L_{j}^{\dagger}}_{\text{dissipative part}} - \underbrace{\frac{decay}{1}_{2} \{L_{j}^{\dagger}L_{j},\rho\}}_{\text{dissipative part}} =: \mathcal{L}[\rho]$$

- ▶ We have a (small) quantum system that is coupled to a (large) environment
- Early proposal: simulate both system and environment on a quantum computer
  - See Terhal, DiVincenzo '98
  - Very resource intensive: many qubits are devoted to environment simulation
- We would like to only simulate what happens with system qubits
- Idea describe effective dynamics using Lindbladian Master equation:

$$\frac{d\rho}{dt} = \underbrace{-i[H,\rho]}_{\text{(Schrödinger equation)}} + \underbrace{\sum_{j=0}^{m} \underbrace{L_j[\rho] L_j^{\dagger}}_{\text{(Issipative part)}} - \underbrace{\frac{decay}{1}}_{\text{(Issipative part)}} =: \mathcal{L}[\rho]$$

L<sub>j</sub>-s describe infinitesimal quantum transitions

- ▶ We have a (small) quantum system that is coupled to a (large) environment
- Early proposal: simulate both system and environment on a quantum computer
  - See Terhal, DiVincenzo '98
  - Very resource intensive: many qubits are devoted to environment simulation
- We would like to only simulate what happens with system qubits
- Idea describe effective dynamics using Lindbladian Master equation:

$$\frac{d\rho}{dt} = \underbrace{-i[H,\rho]}_{\text{(Schrödinger equation)}} + \underbrace{\sum_{j=0}^{m} \underbrace{L_j[\rho]L_j^{\dagger}}_{\text{(Jachrödinger equation)}} - \underbrace{\frac{decay}{1}_{2} \{L_j^{\dagger}L_j,\rho\}}_{\text{dissipative part}} =: \mathcal{L}[\rho]$$

L<sub>j</sub>-s describe infinitesimal quantum transitions

After time t the induced channel is the superoperator

 $\exp(t\mathcal{L}[\cdot])$ 

#### **Continuous-time Markov chains**

▶ We have a continuous-time Markov chain with generator L

- ▶ We have a continuous-time Markov chain with generator L
  - ▶ The off-diagonal entries of *L* are the (non-negative) jump rates

- ▶ We have a continuous-time Markov chain with generator L
  - ▶ The off-diagonal entries of *L* are the (non-negative) jump rates
  - ► The diagonal entry is minus the sum of the off-diagonal elements in the column

- ▶ We have a continuous-time Markov chain with generator L
  - ▶ The off-diagonal entries of *L* are the (non-negative) jump rates
  - ► The diagonal entry is minus the sum of the off-diagonal elements in the column
  - ▶ I.e., *L* is the Laplacian of a weighted directed graph

- ▶ We have a continuous-time Markov chain with generator L
  - ► The off-diagonal entries of *L* are the (non-negative) jump rates
  - ► The diagonal entry is minus the sum of the off-diagonal elements in the column
  - ▶ I.e., *L* is the Laplacian of a weighted directed graph
- After time t the Markov transition matrix is exp(tL)

- ▶ We have a continuous-time Markov chain with generator L
  - ► The off-diagonal entries of *L* are the (non-negative) jump rates
  - ► The diagonal entry is minus the sum of the off-diagonal elements in the column
  - ▶ I.e., *L* is the Laplacian of a weighted directed graph
- After time t the Markov transition matrix is exp(tL)

#### **Continuous-time Markov chains**

We have a continuous-time Markov chain with generator L

- The off-diagonal entries of L are the (non-negative) jump rates
- The diagonal entry is minus the sum of the off-diagonal elements in the column
- I.e., L is the Laplacian of a weighted directed graph
- After time t the Markov transition matrix is exp(tL)

(Do not confuse the Laplacian L with the Lindblad operators  $L_{j}$ .)

$$\frac{d\rho}{dt} = \underbrace{-i[H,\rho]}_{\text{(Schrödinger equation)}} + \underbrace{\sum_{j=0}^{m} \underbrace{L_{j}[\rho]L_{j}^{\dagger}}_{\text{dissipative part}} - \underbrace{\frac{1}{2} \{L_{j}^{\dagger}L_{j},\rho\}}_{\text{dissipative part}} =: \mathcal{L}[\rho]$$

#### **Continuous-time Markov chains**

We have a continuous-time Markov chain with generator L

- The off-diagonal entries of L are the (non-negative) jump rates
- The diagonal entry is minus the sum of the off-diagonal elements in the column
- I.e., L is the Laplacian of a weighted directed graph
- After time t the Markov transition matrix is exp(tL)

(Do not confuse the Laplacian L with the Lindblad operators  $L_{j}$ .)



 $L_i$ -s describe infinitesimal quantum transitions

#### **Continuous-time Markov chains**

We have a continuous-time Markov chain with generator L

- The off-diagonal entries of L are the (non-negative) jump rates
- The diagonal entry is minus the sum of the off-diagonal elements in the column
- I.e., L is the Laplacian of a weighted directed graph
- After time t the Markov transition matrix is exp(tL)

(Do not confuse the Laplacian L with the Lindblad operators  $L_{j}$ .)



 $L_j$ -s describe infinitesimal quantum transitions After time *t* the induced channel is the superoperator  $\exp(t\mathcal{L}[\cdot])$ 

## New generic input model

#### Block-encoding of a Lindbladian

Given a purely dissipative Lindbladian

$$\mathcal{L}[oldsymbol{
ho}] := \sum_{j \in J} \Bigl( oldsymbol{L}_j oldsymbol{
ho} oldsymbol{L}_j^\dagger - rac{1}{2} oldsymbol{L}_j^\dagger oldsymbol{L}_j oldsymbol{
ho} - rac{1}{2} oldsymbol{
ho} oldsymbol{L}_j^\dagger oldsymbol{L}_j \Bigr),$$

we say that a unitary matrix is a block-encoding of the Lindblad operators  $\{L_j\}_{j\in J}$  if

$$(\left\langle 0^{b} \middle| \otimes I \right) \cdot \cdot (\left| 0^{c} \right\rangle \otimes I) = \sum_{j \in J} \ket{j} \otimes L_{j} \quad ext{for} \quad b, c \in \mathbb{N}.$$

## New generic input model

#### Block-encoding of a Lindbladian

Given a purely dissipative Lindbladian

$$\mathcal{L}[
ho] := \sum_{j \in J} \Bigl( oldsymbol{L}_j 
ho oldsymbol{L}_j^\dagger - rac{1}{2} oldsymbol{L}_j^\dagger oldsymbol{L}_j 
ho - rac{1}{2} oldsymbol{
ho} oldsymbol{L}_j^\dagger oldsymbol{L}_j \Bigr),$$

we say that a unitary matrix is a block-encoding of the Lindblad operators  $\{L_j\}_{j\in J}$  if

$$(\left\langle 0^{b} \middle| \otimes I \right) \cdot \cdot (\left| 0^{c} \right\rangle \otimes I) = \sum_{j \in J} \ket{j} \otimes L_{j} \quad ext{for} \quad b, c \in \mathbb{N}.$$

► In principle a block-encoding exists if  $\left\|\sum_{j=1}^{m} \boldsymbol{L}_{j}^{\dagger}\boldsymbol{L}_{j}\right\| \leq 1$ 

## New generic input model

#### **Block-encoding of a Lindbladian**

Given a purely dissipative Lindbladian

$$\mathcal{L}[oldsymbol{
ho}] := \sum_{j \in J} \Bigl( oldsymbol{L}_j oldsymbol{
ho} oldsymbol{L}_j^\dagger - rac{1}{2} oldsymbol{L}_j^\dagger oldsymbol{L}_j oldsymbol{
ho} - rac{1}{2} oldsymbol{
ho} oldsymbol{L}_j^\dagger oldsymbol{L}_j \Bigr),$$

we say that a unitary matrix is a block-encoding of the Lindblad operators  $\{L_j\}_{j\in J}$  if

$$(\left\langle 0^{b} \middle| \otimes I \right) \cdot \cdot (\left| 0^{c} \right\rangle \otimes I) = \sum_{j \in J} \ket{j} \otimes L_{j} \quad \text{for} \quad b, c \in \mathbb{N}.$$

► In principle a block-encoding exists if  $\left\|\sum_{j=1}^{m} \boldsymbol{L}_{j}^{\dagger}\boldsymbol{L}_{j}\right\| \leq 1$ 

• Prior methods depended on  $\sum_{j=0}^{m} \left\| \boldsymbol{L}_{j}^{\dagger} \boldsymbol{L}_{j} \right\| \leq 1$ 

## Weak measurement scheme for Lindbladians

#### **Block-encoding of Lindblad generators**

We say that the unitary U is a block encoding of the purely irreversible generator  $\mathcal{L}$  consisting of Lindblad operators  $\mathbf{L}_j$  if  $(\langle 0^b | \otimes I ) U(|0^a \rangle \otimes I) = \sum_{j=0}^m |j\rangle \otimes \mathbf{L}_j$ .



Here  $Y_{\delta} = e^{-i \operatorname{arcsin} \sqrt{\delta}Y}$  for the Pauli-Y gate.

1. Apply  $\boldsymbol{U}$ .

- 2. Append an ancilla qubit in state  $|0\rangle$  and rotate it with angle  $\arcsin \sqrt{\delta}$  controlled on the  $|0^b\rangle$  state (indicating the successful application of a jump).
- 3. Apply  $\boldsymbol{U}^{\dagger}$  controlled on the ancilla qubit being 0.
- 4. Measure and discard all but the system register.

Assuming the system register is in the pure state  $|\psi\rangle$ , this circuit C acts as follows:

$$\begin{split} |0\rangle \cdot |0^{c}\rangle |\psi\rangle \stackrel{(1)}{\rightarrow} |0\rangle \cdot \mathbf{U}|0^{c}\rangle |\psi\rangle \\ \stackrel{(2)}{\rightarrow} \left(\sqrt{1-\delta}|0\rangle + \sqrt{\delta}|1\rangle\right) \cdot \left(|0^{b}\rangle\langle 0^{b}| \otimes \mathbf{I}\right) \mathbf{U}|0^{c}\rangle |\psi\rangle + |0\rangle \cdot (\mathbf{I} - |0^{b}\rangle\langle 0^{b}| \otimes \mathbf{I}) \mathbf{U}|0^{c}\rangle |\psi\rangle \\ &= |0\rangle \cdot \mathbf{U}|0^{c}\rangle |\psi\rangle + \sqrt{\delta}|1\rangle \cdot |0^{b}\rangle \underbrace{\left(\langle 0^{b}| \otimes \mathbf{I}\right) \mathbf{U}|0^{c}\rangle |\psi\rangle}_{|\psi_{0}^{c}\rangle :=} - (1 - \sqrt{1-\delta})|0\rangle \cdot \left(|0^{b}\rangle\langle 0^{b}| \otimes \mathbf{I}\right) \mathbf{U}|0^{c}\rangle |\psi\rangle \\ &= |0\rangle \cdot |0^{c}\rangle |\psi\rangle + \sqrt{\delta}|1\rangle \cdot |0^{b}\rangle |\psi_{0}^{\prime}\rangle - (1 - \sqrt{1-\delta})|0\rangle \cdot \mathbf{U}^{\dagger}(|0^{b}\rangle\langle 0^{b}| \otimes \mathbf{I}) \mathbf{U}|0^{c}\rangle |\psi\rangle \\ &= |0\rangle \cdot |0^{c}\rangle |\psi\rangle + \sqrt{\delta}|1\rangle \cdot |0^{b}\rangle |\psi_{0}^{\prime}\rangle - (1 - \sqrt{1-\delta})|0\rangle \cdot |0^{c}\rangle (\langle 0^{c}| \otimes \mathbf{I}) \mathbf{U}^{\dagger}(|0^{b}\rangle \otimes \mathbf{I}) \cdot (\langle 0^{b}| \otimes \mathbf{I}) \mathbf{U}|0^{c}\rangle |\psi\rangle \\ &- (1 - \sqrt{1-\delta})|0\rangle \cdot (\mathbf{I} - |0^{c}\rangle\langle 0^{c}| \otimes \mathbf{I}) \mathbf{U}^{\dagger}(|0^{b}\rangle\langle 0^{b}| \otimes \mathbf{I}) \mathbf{U}|0^{c}\rangle |\psi\rangle \\ &= |0\rangle \cdot |0^{c}\rangle \left(\mathbf{I} - \underbrace{(1 - \sqrt{1-\delta})}_{\frac{\delta}{2} + \mathcal{O}(\delta^{2})} \sum_{j \in J} \mathbf{L}_{j}^{\dagger} \mathbf{L}_{j}\right) |\psi\rangle + \sqrt{\delta}|1\rangle \cdot |0^{b}\rangle \sum_{j \in J} |j\rangle \mathbf{L}_{j}|\psi\rangle - \underbrace{(1 - \sqrt{1-\delta})}_{\frac{\delta}{2} + \mathcal{O}(\delta^{2})} |0\rangle \cdot |0^{c} \perp\rangle, (3.2) \end{split}$$

# Simulation uses quantum Zeno-like effect

## Simulation uses quantum Zeno-like effect

Combine with "compression" cf. "Efficient quantum algorithms for simulating Lindblad evolution" by Richard Cleve, Chunhao Wang (2017)



# **Davies generator – simulating thermalization?**

$$\mathcal{L}_{Davies}(\rho) = \sum_{a \in A} \sum_{\nu \in B} \gamma(\nu) \left( \underbrace{\mathbf{A}_{\nu}^{a} \rho(\mathbf{A}_{\nu}^{a})^{\dagger}}_{\text{transition rates}} - \underbrace{\frac{1}{2} ((\mathbf{A}_{\nu}^{a})^{\dagger} \mathbf{A}_{\nu}^{a} \rho + \rho(\mathbf{A}_{\nu}^{a})^{\dagger} \mathbf{A}_{\nu}^{a})}_{\text{decay rates}} \right)$$

where the quantum mechanical transition rates  $\mathbf{A}_{\nu}^{a}$  are defined as

$$oldsymbol{A}^a_{
u}:=\sum_{E_i-E_j=
u}oldsymbol{P}_{E_i}oldsymbol{A}^aoldsymbol{P}_{E_j}$$
 for Bohr frequency  $u\in spec(oldsymbol{H})-spec(oldsymbol{H})=:B(oldsymbol{H}),$ 

## **Davies generator – simulating thermalization?**

$$\mathcal{L}_{Davies}(\rho) = \sum_{a \in A} \sum_{\nu \in B} \gamma(\nu) \left( \underbrace{\mathbf{A}_{\nu}^{a} \rho(\mathbf{A}_{\nu}^{a})^{\dagger}}_{\text{transition rates}} - \underbrace{\frac{1}{2} ((\mathbf{A}_{\nu}^{a})^{\dagger} \mathbf{A}_{\nu}^{a} \rho + \rho(\mathbf{A}_{\nu}^{a})^{\dagger} \mathbf{A}_{\nu}^{a})}_{\text{decay rates}} \right)$$

where the quantum mechanical transition rates  $\mathbf{A}_{\nu}^{a}$  are defined as

$$oldsymbol{A}^a_
u := \sum_{E_i - E_j = 
u} oldsymbol{P}_{E_i} oldsymbol{A}^a oldsymbol{P}_{E_j}$$
 for Bohr frequency  $u \in spec(oldsymbol{H}) - spec(oldsymbol{H}) =: B(oldsymbol{H})$ 

► Folklore knowledge: "unphysical exponential relaxation", etc.

Simulating Nature on a quantum computer should be "easy"

- Simulating Nature on a quantum computer should be "easy"
- Preparing Gibbs states can be QMA-hard!

- Simulating Nature on a quantum computer should be "easy"
- Preparing Gibbs states can be QMA-hard!
- So Nature might get stuck at local minima ...

- Simulating Nature on a quantum computer should be "easy"
- Preparing Gibbs states can be QMA-hard!
- So Nature might get stuck at local minima ...
- "Local minima in quantum systems" by Chi-Fang Chen, Hsin-Yuan Huang, John Preskill, Leo Zhou Sep 29 2023 arXiv:2309.16596

## **Back to first principles**

"Coarse graining can beat the rotating-wave approximation in quantum markovian master equations" Christian Majenz, Tameem Albash, Heinz-Peter Breuer, and Daniel A. Lidar. Phys. Rev. A, (2013)

## **Back to first principles**

 "Coarse graining can beat the rotating-wave approximation in quantum markovian master equations" Christian Majenz, Tameem Albash, Heinz-Peter Breuer, and Daniel A. Lidar. Phys. Rev. A, (2013)

$$\mathcal{L}_{\beta}[\rho] := \sum_{a \in \mathcal{A}} \int_{-\infty}^{\infty} \gamma(\omega) \Big( \hat{\boldsymbol{A}}^{a}(\omega) \rho \hat{\boldsymbol{A}}^{a}(\omega)^{\dagger} - \frac{1}{2} \Big\{ \hat{\boldsymbol{A}}^{a}(\omega)^{\dagger} \hat{\boldsymbol{A}}^{a}(\omega), \rho \Big\} \Big) d\omega$$

where  $\gamma(\omega) = \min(1, e^{-\beta\omega})$  may be the Metropolis weight or the (smoother) Glauber dynamics weight  $\gamma(\omega) = (e^{\beta\omega} + 1)^{-1}$ , or something similar and

$$\hat{\mathbf{A}}^{a}(\omega) :\propto \int_{-T/2}^{T/2} e^{-i\omega t} e^{i\mathbf{H}t} \mathbf{A}^{a} e^{-i\mathbf{H}t} dt$$
 for each  $a \in \mathbf{A}$  and  $\omega \in \mathbb{R}$ .

# **Discretization + Operator Fourier transform**

► We want to turn *L* to

$$\sum_{\nu} \gamma(
u) |
u 
angle \mathcal{L}^{(
u)}$$
 where  $\mathcal{L}^{(
u)} = \sum_{\psi, \psi' : |\nu| = E_{\psi'} - E_{\psi}} |\psi' \chi \psi'| \mathcal{L} |\psi \chi \psi|$ 

## **Discretization + Operator Fourier transform**

We want to turn L to

$$\sum_{
u} \gamma(
u) |
u 
angle oldsymbol{L}^{(
u)}$$
 where  $oldsymbol{L}^{(
u)} = \sum_{\psi, \psi' \colon 
u = E_{\psi'} - E_{\psi}} |\psi' 
angle \psi' |oldsymbol{L}| \psi 
angle \psi|$ 

► (Flat) Operator Fourier transform:

$$\sum_{t=-T/2}^{T/2} \frac{1}{T} |t\rangle \boldsymbol{L} \to \sum_{t} \frac{1}{T} |t\rangle \underbrace{\exp(iHt)\boldsymbol{L}}_{\sum_{v} \exp(ivt)\boldsymbol{L}^{(v)}} \to \sum_{v} \underbrace{\sum_{w} \hat{f}(w-v) |w\rangle}_{\text{peak at } v} \boldsymbol{L}^{(v)}$$

(Show figure from paper block-encoding the jumps.)

Block-encoding the Lindblad operators via quantum operator Fourier transform:



## A coherent Szegedy quantum walk type variant

(Discretized) discriminant proxy

$$\sum_{\mathsf{a}\in\mathsf{A}}\int\sqrt{\gamma(\omega)\gamma(-\omega)}\hat{\boldsymbol{\mathsf{A}}}^{\mathsf{a}}(\omega)\otimes\hat{\boldsymbol{\mathsf{A}}}^{\mathsf{a}*}(\omega)-\frac{\gamma(\omega)}{2}\big(\boldsymbol{\mathsf{A}}^{\mathsf{a}}(\omega)^{\dagger}\hat{\boldsymbol{\mathsf{A}}}^{\mathsf{a}}(\omega)\otimes\boldsymbol{\mathsf{I}}+\boldsymbol{\mathsf{I}}\otimes\hat{\boldsymbol{\mathsf{A}}}^{\mathsf{a}}(\omega)^{\dagger*}\hat{\boldsymbol{\mathsf{A}}}^{\mathsf{a}}(\omega)^{*}\big)\,\mathsf{d}\omega$$

If the mixing time is sufficiently short then provides access to (approximately) the following specific purification of the Gibbs state

$$ig|\sqrt{oldsymbol{
ho}_eta}ig
angle\propto\sum_i e^{-eta E_i/2}|\psi_i
angle\otimesig|\psi_i^*
angle$$
 where  $oldsymbol{H}=\sum_i E_i|\psi_i
angle\psi_i|$ 

In which (physical) systems can we expect rapid convergence?

- In which (physical) systems can we expect rapid convergence?
- How to bound the gap of the generator?

- In which (physical) systems can we expect rapid convergence?
- How to bound the gap of the generator?
- How noise resilient is this algorithm?

- In which (physical) systems can we expect rapid convergence?
- How to bound the gap of the generator?
- How noise resilient is this algorithm?
- Finally a quadratic improvement for carbon capture?