

# Randomness (and Pseudorandomness)

Avi Wigderson  
IAS, Princeton

# Plan of the talk

# Perfect Randomness

Uniform distribution  
Full entropy

Prob[T]=Prob[H]= $\frac{1}{2}$ , Independent

HHTHTTTTHHHTHTTHTTTTHHHHTHTTTTTTHHHTHH  
HH



The amazing utility of randomness -  
lots of examples. Are they real?

# Pseudorandomness

Deterministic structures which share *some* properties of random ones

# Lots of examples & applications

## Surviving weak or no randomness

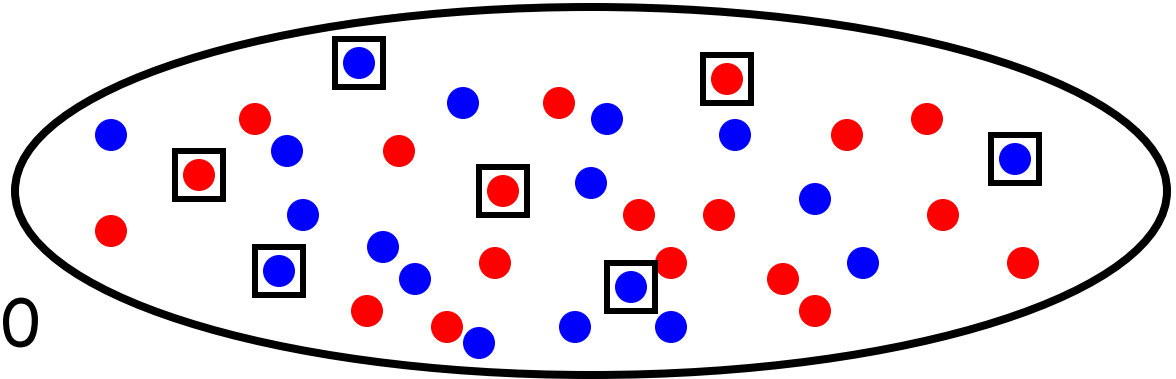
# The remarkable utility of randomness



# Fast Information Acquisition

Population: 280 million, voting blue or red

Random  
Sample: 2,000



Theorem: With probability  $> .99$   
% in sample = % in population  $\pm 2\%$   
independent of population size

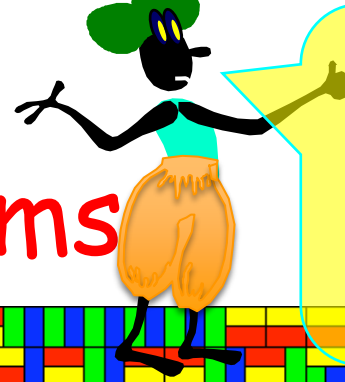
Deterministically, need to  
ask the whole population!



Where are  
the random  
bits from?

# Efficient (!!)

# Probabilistic Algorithms



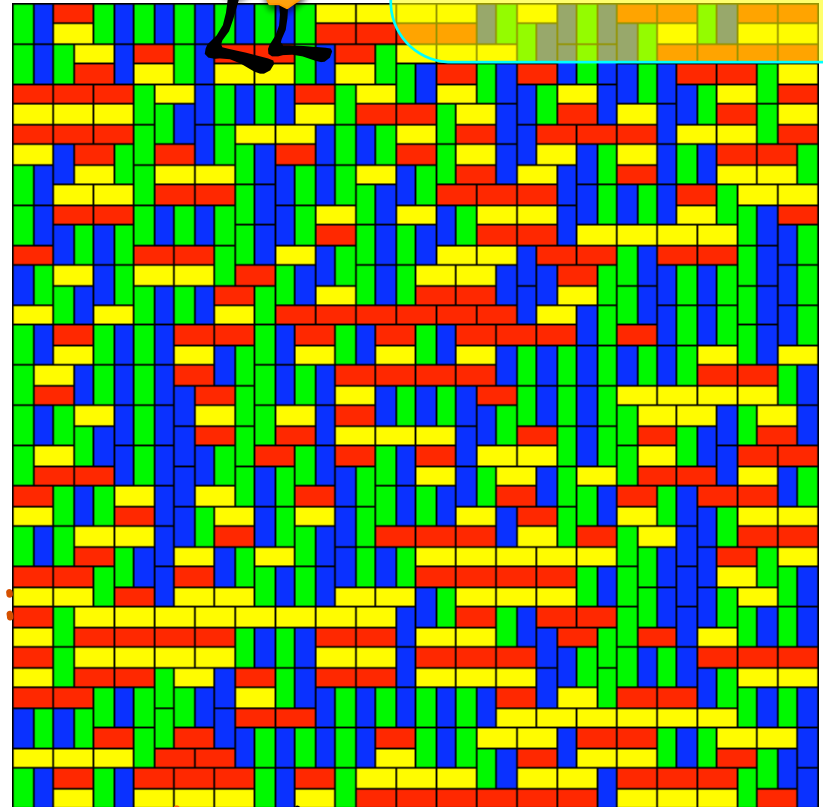
Where are  
the random  
bits from?

Given a region in space, how many domino tilings does it have?

**Monomer-Dimer problem.**

Captures thermodynamic  
properties of matter  
(free energy, phase transitions,...)

**Theorem** [Jerrum-Sinclair-Vigoda]:  
Efficient probabilistic  
approximate counting algorithm  
("Monte-Carlo" method [von Neumann-Ulam])



Best deterministic algorithm **known** requires exponential time!  
**One of numerous examples**

# Probabilistic algorithms in math

## Number theory

- Finding large primes

## Algebra

- Factoring multivariate polynomials over finite fields

## Geometry

- Approximating the volume of convex sets in high dimension

## Analysis

- Computing large Fourier coeffs of multivariate functions

Fast probabilistic algorithms

Best known deterministic algorithms require exponential time



Where are  
the random  
bits from?

# Distributed computation

Randomness makes impossible problems possible

Where are the random bits from?

## The dining philosophers problem

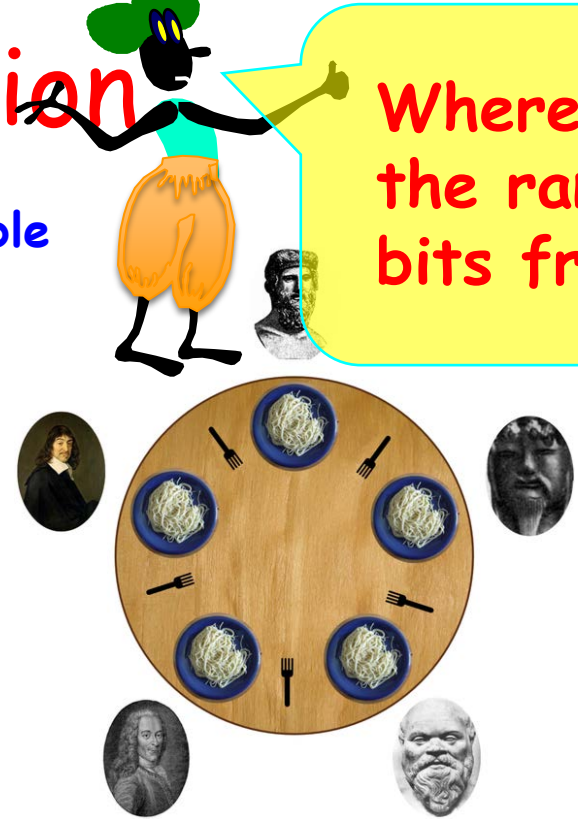
Captures resource allocation and sharing in asynchronous systems

Theorem [Dijkstra]:

No deterministic solution

Theorem [Lehman-Rabin]:

A probabilistic program works



## The Byzantine generals problem

Captures coordination with faults

Theorem [Fisher-Lynch-Paterson]:

No deterministic solution

Theorem [Ben-or, Rabin]:

A probabilistic program works



Byzantine Generals Problem

# Game Theory – Rational behavior

## Chicken game [Aumann]

A: Aggressive

C: Cautious

C	1 1	0 2
A	2 0	-3 -3
	C	A



**Nash Equilibrium:** No player has an incentive to change its strategy given the opponent's strategy.

**Theorem [Nash]:** Every game has an equilibrium in mixed (random) strategies  
( $\Pr[C] = \frac{3}{4}$ ,  $\Pr[A] = \frac{1}{4}$ )

False for pure (deterministic) strategies



Where are the random bits from?



# Cryptography & E-commerce

## Secrets

**Theorem [Shannon]** A secret is as good as the entropy in it. (if you pick a 9 digit password randomly, my chances of guessing it is  $1/10^9$ )

**Public-key encryption** (on-line shopping)

**Digital signature** (identification)

**Zero-Knowledge Proofs** (enforcing correctness)

.....

All require randomness



Where are the random bits from?

# Gambling



Where are  
the random  
bits from?

Search

About 3,730,000 results (0.20 seconds)

Everything

Images

Maps

Videos

News

Shopping

More

Princeton, NJ

Change location

Any time

Past hour

Past 24 hours

Past week

Past month

Past year

Custom range...

All results

Sites with images

Related searches

Timeline

### HotBits: Genuine Random Numbers

[www.fourmilab.ch/hotbits/](http://www.fourmilab.ch/hotbits/)

HotBits is a service which generates **random** data from the decay of radioactive material and sends it over the internet.

### Hardware random number generator - Wikipedia, the free ...

[en.wikipedia.org/wiki/Hardware\\_random\\_number\\_generator](http://en.wikipedia.org/wiki/Hardware_random_number_generator)

Other designs use what are believed to be **true random bits** as the key for a high quality block cipher algorithm, taking the encrypted output as the random bit ...

### RANDOM.ORG - True Random Number Service

[www.random.org/](http://www.random.org/)

ORG offers **true** random numbers to anyone on the Internet. The randomness comes ...  
ORG has generated 1032 billion **random bits** for the Internet community. ...

### Quantis RNG - True Random Number Generator - Overview

[www.idquantique.com/true-random-number.../products-overview.ht...](http://www.idquantique.com/true-random-number.../products-overview.ht...)

Contrary to existing products, Quantis produces random numbers at a very high bite r up to 16Mbps. This is the highest **truly-random bit** rate available to date. ...

### True random number generators

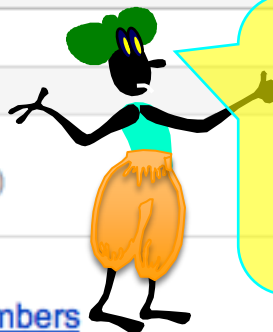
[www.robertnz.net/true\\_rng.html](http://www.robertnz.net/true_rng.html)

**true random** number generator. ... If you sample the output (not too quickly) you (hope to) get a series of **bits** which are statistically independent. These can be ...

### Quantum Random Bit Generator Service

[random.irb.hr/](http://random.irb.hr/)

**true randomness** of data served (high per-bit-entropy of served data); high speed of data generation and serving; high availability of the service (including easy ...



Where are  
the random  
bits from?

Radiative  
decay

Atmospheric  
noise



Is it working?  
[C,VV,MS,...]

Photons  
measurement

# Defining randomness



# What is random?



I toss the coin, you guess how it will land

Probability of guessing correctly?  $\frac{1}{2}$

Randomness is in the ~~eye~~ of the beholder  
computational power

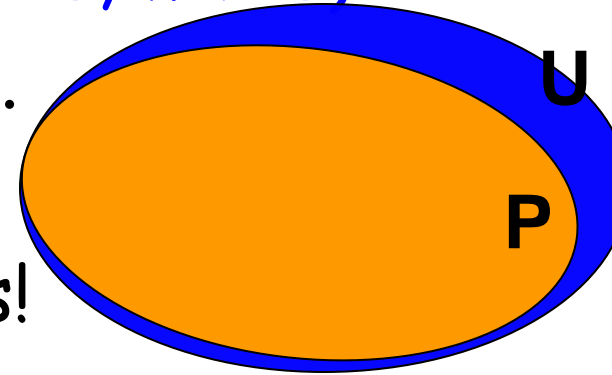
Operative, subjective definition!

# Pseudorandomness

The study of deterministic structures  
(numbers, graphs, sequences, tables, walks)  
with **some** "random-like" **properties**.

Almost all  $x \in U$  is also in  $P$

$P$  - what the limited observer tests!



**Mathematics:** Study of random-like properties  
of *natural* structures: "is  $x_0 \in P$ ?"

**Computer Science:** Efficiently *find* structures  
with random-like properties:

"find any  $x \in P$ "

"find hay in a haystack"



# Normal Numbers

3.1415926535 8979323846 2643383279 5028841971 6939937510 5820974944 5923078164  
0628620899 8628034825 3421170679 8214808651 3282306647 0938446095 5058223172  
5359408128 4811174502 8410270193 8521105559 6446229489 5493038196 4428810975 6659334461  
2847564823 3786783165 2712019091 4564856692 3460348610 4543266482 1339360726  
0249141273 7245870066 0631558817 4881520920 9628292540 9171536436 7892590360  
0113305305 4882046652 1384146951 9415116094 3305727036 5759591953 0921861173 8193261179  
3105118548 0744623799 6274956735 1885752724 8912279381  
8301194912 9833673362 4406566430 8602139494 6395224737 1907021798 6094370277  
0539217176 2931767523 8467481846 7669405132 0005681271 4526356082 7785771342  
7577896091 7363717872 1468440901 2249534301 4654958537 1050792279 6892589235  
4201995611 2129021960 8640344181 5981362977 4771309960 5187072113 4999999837 2978049951  
0597317328 1609631859 .....

- Every digit (e.g. 7) occurs  $1/10$  th of the time,
- Every pair (e.g. 54) occurs  $1/100$  th of the time,
- Every triple (eg 666) occurs  $1/1000$  th of the time,...

in every base!

Pseudorandom  
Property

**Theorem[Borel]:** A random real number is normal

**Open:** Is  $\pi$  normal? Are  $\sqrt{2}$ ,  $e$  normal?

(we know efficient algs for generating normal numbers)

# Major problems of Math & CS are about Pseudorandomness

Clay Millennium Problems - \$1M each

- Birch and Swinnerton-Dyer Conjecture
- Hodge Conjecture
- Navier-Stokes Equations

Pseudorandom  
Property

- **P vs. NP**

Random functions are hard to compute.  
Prove the same for the TSP  
(Traveling Salesman Problem)!



~~Poincaré Conjecture~~

Pseudorandom  
Property

- **Riemann Hypothesis**

Random walks stay close to  
the origin. Prove the same for  
the Möbius walk!

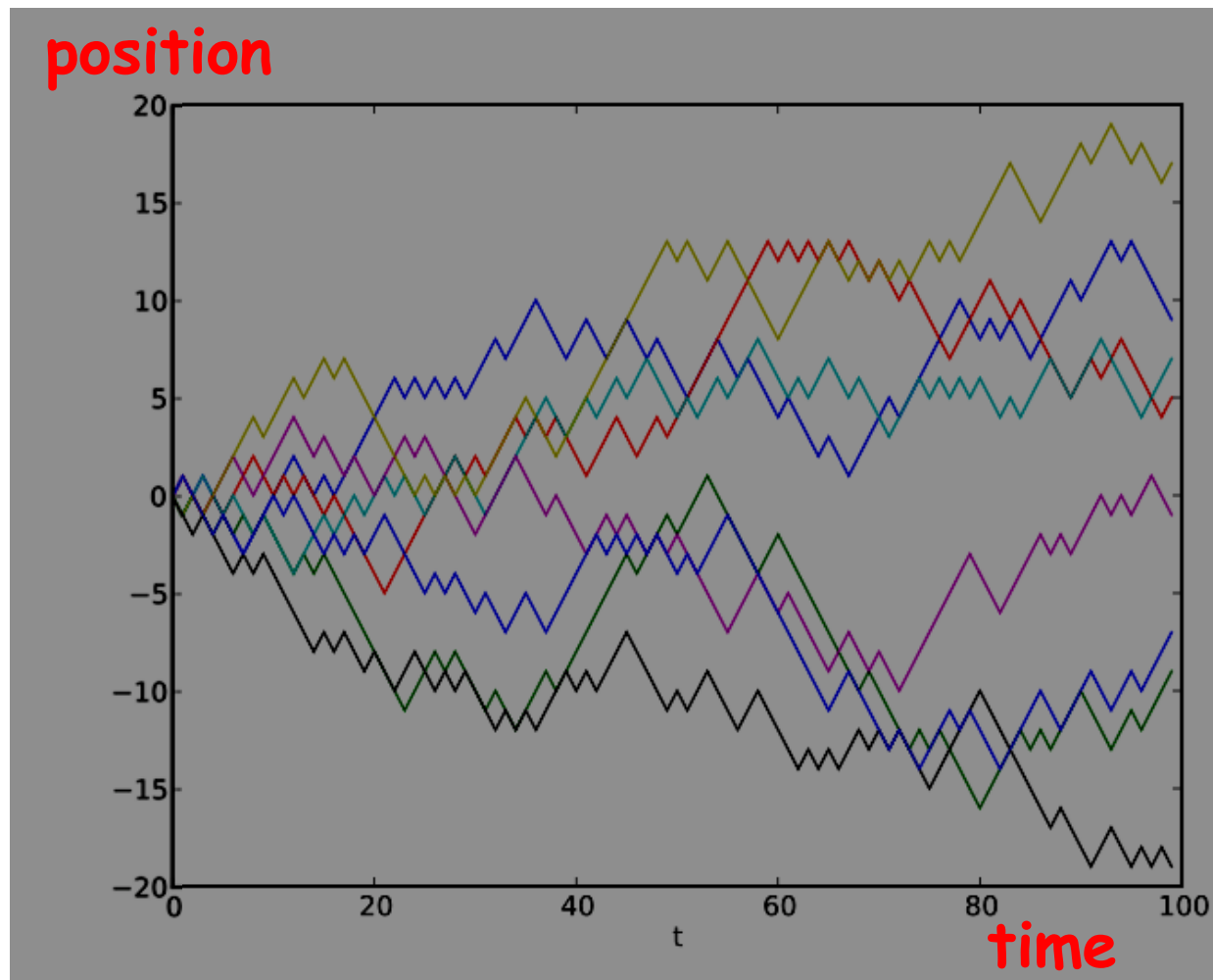
- Yang-Mills Theory



# Riemann Hypothesis & the drunkard's walk

Start: 0  
Each step -  
Up: +1  
Down: -1  
Randomly.

Almost surely,  
after  $N$  steps  
distance from  
0 is  $\sim\sqrt{N}$



# Möbius' walk

$x$  integer,  $p(x)$  number of distinct prime divisors

$$\mu(x) = \begin{cases} 0 & \text{if } x \text{ has a square divisor} \\ 1 & p(x) \text{ is even} \\ -1 & p(x) \text{ is odd} \end{cases}$$

$x =$	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
$\mu(x) =$	1	-1	-1	0	-1	1	-1	0	0	1	-1	0	-1	1	1	0

**Theorem [Mertens 1897]:** These are equivalent:

- For all  $N$   $|\sum_{x \leq N} \mu(x)| \sim \sqrt{N}$
- The Riemann Hypothesis

# Coping in a world without perfect randomness

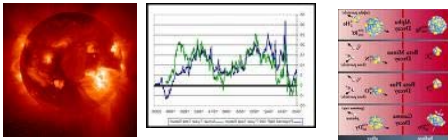
Major developments of  
the last 3 decades

# Possible worlds

- Perfect randomness



- Weak random sources  
Biased, dependent bits



- No randomness  
assuming " $P \neq NP$ "

All require different  
Pseudorandom notions

# Applications

all applications

all algorithms

Extractor theory

[B,SV,NZ,T,...,GUV,DW,...]

purifying randomness

all *efficient* algorithms

Hardness vs. Randomness

[BM,Y,...NW,IW,...]

Every efficient prob alg has  
a deterministic counterpart

# Weak random sources and randomness purification

**Applications:**  
Analyzed on  
perfect  
randomness

Statistics, Cryptography, Algorithms,  
Game theory, Gambling,.....

Unbiased,  
independent



**Extractor Theory**

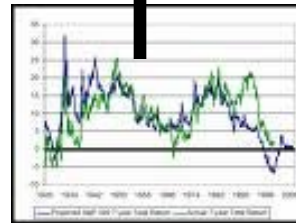
biased,  
dependent



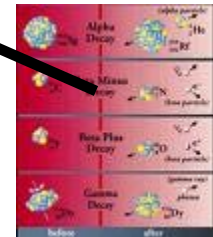
**Reality:**  
Sources of  
imperfect  
randomness



Sun spots



Stock market  
fluctuations



Radioactive  
decay

# Pseudorandom Tables

Random  
nxn table  
Entries in [n]

X	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
1	12	3	10	3	3	7	15	4	9	12	4	1	7	12	11
2	2	3	8	4	3	9	9	11	1	13	1	14	6	6	7
3	1	5	1	8	10	10	3	8	14	2	2	9	8	13	9
4	5	7	12	3	11	4	8	10	4	3	8	8	7	1	11
5	4	15	3	6	7	2	2	1	2	15	8	3	10	2	1
6	2	5	1	8	4	10	3	8	4	2	1	9	8	13	3
7	5	4	14	11	8	7	9	4	9	1	4	3	7	12	8
8	6	2	8	6	13	5	9	11	1	13	14	8	4	6	12
9	8	5	13	2	7	2	2	1	2	15	8	3	10	2	1
10	11	4	14	5	13	5	4	3	6	7	9	2	1	8	8
11	15	4	9	12	4	1	7	12	11	3	6	7	2	12	15
12	1	8	10	10	3	8	14	6	2	8	6	13	5	9	11
13	7	12	3	8	6	5	13	5	9	11	1	2	15	8	3
14	10	3	3	9	8	14	6	2	11	1	2	13	14	8	4
15	12	10	9	1	3	10	3	8	7	15	4	9	2	5	15

Thm: In a random matrix, *every small* window is

"rich": have *many* different entries.

Pseudorandom  
Property

rich: Eg  $k \times k$  windows to have  $k^{1.1}$  distinct  
entries,

for  $k \sim n^{0.1}$

# Addition table

+	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18
4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21
7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22
8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23
9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26
12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27
13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28
14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29
15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30

1983 Newton's "Any one who considers arithmetical methods  
Every man doing with his right hand, of these tables as state of sin."

# Independent-source extractors

**Applications:**  
Analyzed on  
perfect  
randomness

Statistics, Cryptography, Algorithms,  
Game theory, Gambling,.....

Unbiased,  
independent



**Extractor**

biased,  
dependent

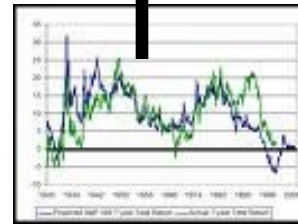


**Reality:**  
*Independent*  
weak sources  
of randomness



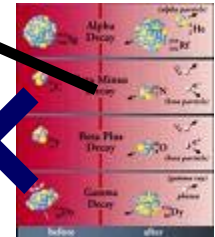
Sun spots

+



Stock market  
fluctuations

×



Radioactive  
decay



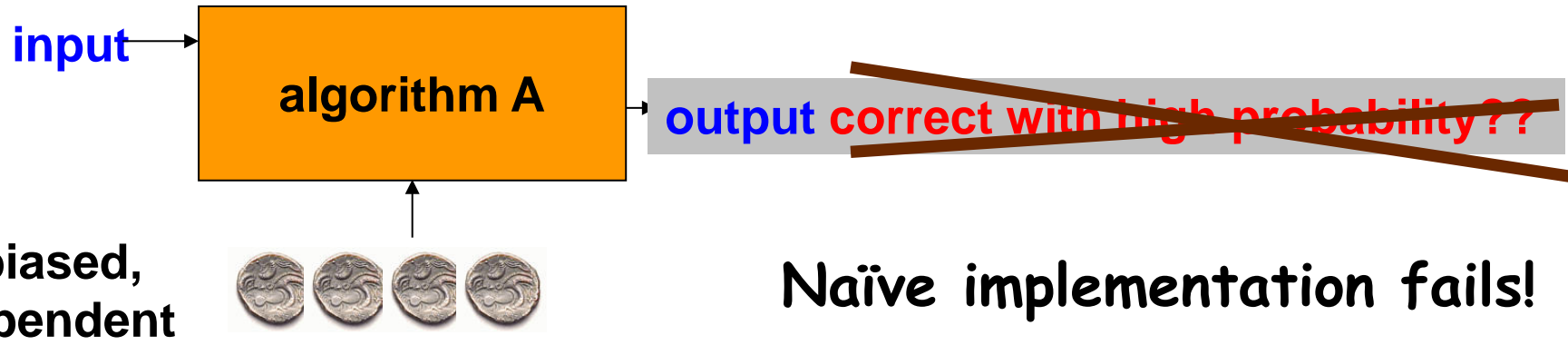
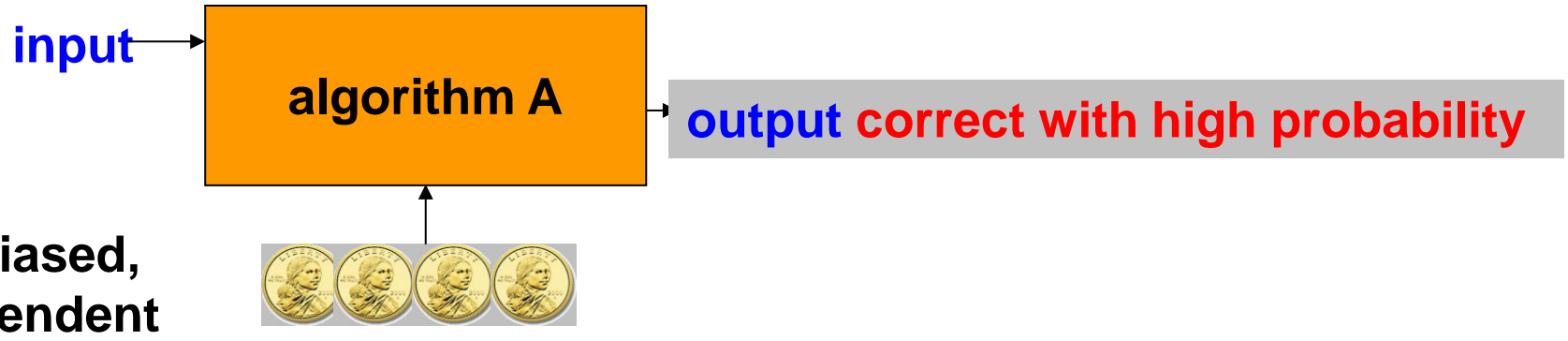
# Summary

- **Randomness** is in the eye of the beholder:  
A pragmatic, subjective definition
- **Pseudorandomness** “tests”  $\leftrightarrow$  “applications”  
Capture many basic problems and areas in Math & CS
- **Applications** of randomness survive in a world without perfect (or any) randomness
- **Pseudorandom** objects often find uses beyond their intended application (expanders, extractors,...)

# Thank you!



# Single-source extractors



Naïve implementation fails!



**Reality:** one weak random source (all randomness correlated)

# Single-source extractors

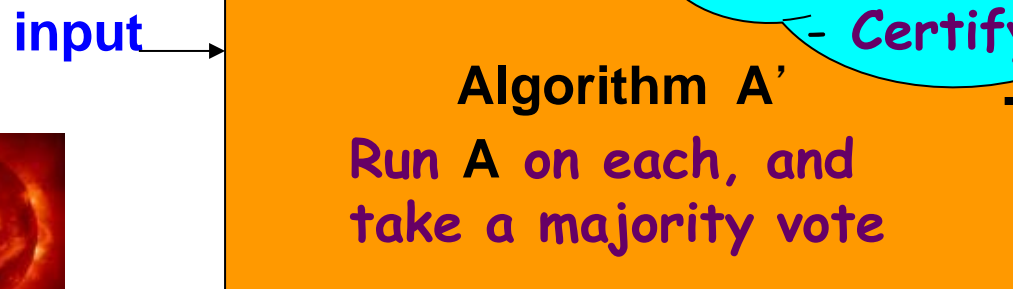
Probabilistic algorithms with 1 weak random source



Many other applications:

- Hardness of approx
- Derandomization,
- Data structures
- Error correction
- **Certifying randomness**

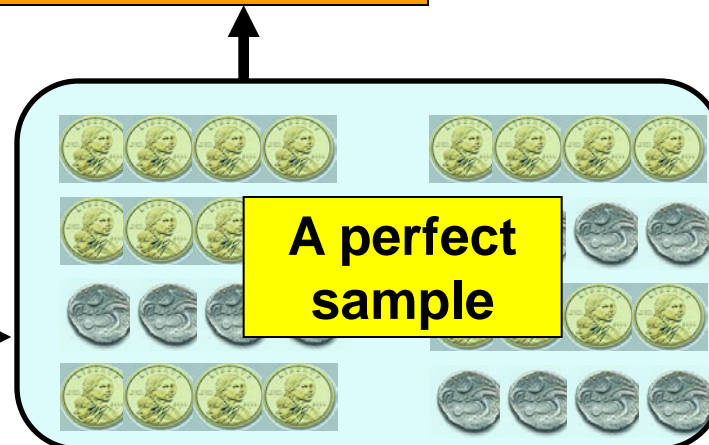
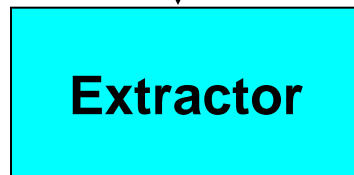
correct whp!!



Run A on each, and  
take a majority vote



$n^3$  biased,  
dependent  
entropy  $n^2$



BI, SV, NZ,  
.....  
Ta, Tr, ISW,  
.....  
LRVW, GUV  
Dv, DW,...

# Deterministic de-randomization

## Hardness vs. Randomness

input

*Efficient*  
algorithm A

output correct with high probability

$n$  unbiased,  
independent



All efficient probabilistic  
algorithms have efficient  
deterministic counterparts

input

Run A on  $n$  unbiased samples  
take a majority vote

ways!

“ $P \neq NP$ ”:  
TSP is  
hard

A perfect  
sample

0010001 1111101  
1100000 0111110  
1100000 1000000  
0010001 0100001  
0101101 1000001

BM, Y, ...  
.....  
NW, IW, ...  
.....  
IKW, KI, ...



**Reality:** Universe is deterministic