# Yoshinaga's Criterion and the topology of complexified complements

## Galen Dorpalen-Barry joint with Graham Denham and Nick Proudfoot

#### **IPAM Workshop**

## Computational Interactions between Algebra, Combinatorics, and Discrete Geometry



## Outline

Motivation

From Geometry to Algebra

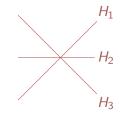
Code and Computation Time

Next Steps

## Motivation

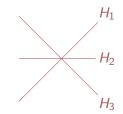
## Arrangements of Hyperplanes in $\mathbb{R}^d$

- A real hyperplane is an affine linear subspace of codimension 1 in V ≅ ℝ<sup>d</sup>.
- A collection of finitely-many (distinct) hyperplanes is an arrangement.



## Arrangements of Hyperplanes in $\mathbb{R}^d$

- A real hyperplane is an affine linear subspace of codimension 1 in V ≅ ℝ<sup>d</sup>.
- A collection of finitely-many (distinct) hyperplanes is an arrangement.



Provide interesting special cases of matroids and oriented matroids, useful tool in for linear optimization, polytopal combinatorics, reflection groups, and geometry.

**Big Question**: what is the topology of the "complexified complement"?

## The Complexified Complement

• The **complexification** of a real vector space V is  $V^{\mathbb{C}} = V + iV$ .

If H is a hyerplane defined as the zero set of f, its complexification H<sup>C</sup> is the zero set of

$$f^{\mathbb{C}}=f+i(f-f(0)).$$

If  $H \subseteq \mathbb{R}^d$ , then  $H^{\mathbb{C}}$  is the zero set of f evaluated on  $\mathbb{C}^d$ .

For a real arrangement A, the complexified complement of A is M(A) = V<sup>C</sup> \ ⋃<sub>H∈A</sub> H<sup>C</sup>.

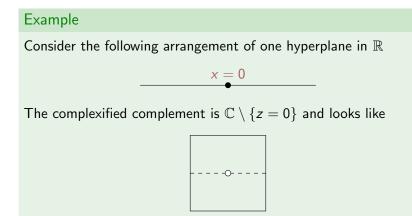
x = 0

#### Example

Consider the following arrangement of one hyperplane in  $\ensuremath{\mathbb{R}}$ 

The complexified complement lives in  $\mathbb C$  ...

## The Complexified Complement



The dashed line represents im(z) = 0 the dot is the puncture where im(z) = 0 and the real part lies on a hyperplane.

## State of the Art

- Cohomology ring of M(A): Orlik-Solomon algebra, depends only on the matroid
- Cohomology ring of a related space M<sub>3</sub>(A): Cordovil algebra, has simple presentation, depends on the oriented matroid (but not the specific orientation)

 $\begin{array}{l} ({\sf Cordovil\ Algebra} = {\sf assoicated}{\sf -graded\ of\ the} \\ {\sf Varchenko}{\sf -Gelfand\ ring} \end{array}$ 

The homotopy type is somewhat understood (there are combinatorial models like Salvetti's complex), but not completely.

**Major Open Question**: Is there a simple combinatorial criterion on  $\mathcal{A}$  that determines when  $\pi_k(\mathcal{M}(\mathcal{A}))$  is trivial for  $k \geq 2$ ?

## Example: Graphical Arrangements

#### Theorem (Stanley, 1972)

Let G be a finite graph and  $A_G$  its graphical arrangement. Then G is supersolvable if and only if every cycle with at least four edges is cut by a chord.

#### Theorem (Terao, 1983)

supersolvable  $\Rightarrow$  the complexified complement is  $K(\pi, 1)$ 

#### Example

A graph G and its graphical arrangement  $\mathcal{A}_G$ . Stanley and Terao tell us that  $M(\mathcal{A}_G)$  is  $K(\pi, 1)$ .

$$2 - 3 | | | 1 - 4 x_1 - x_2 = 0, x_2 - x_3 = 0, x_3 - x_4 x_1 - x_4 = 0, x_1 - x_3 = 0$$

## Example: Graphical Arrangements

#### Theorem (Stanley, 1972)

Let G be a finite graph and  $A_G$  its graphical arrangement. Then G is supersolvable if and only if every cycle with at least four edges is cut by a chord.

#### Theorem (Terao, 1983)

supersolvable  $\Rightarrow$  the complexified complement is  $K(\pi, 1)$ 

#### Example

A graph G and its graphical arrangement  $A_G$ .

Stanley tells us that  $\mathcal{A}_{G}$  is *not* supersolvable, Terao is inconclusive

#### From Falk–Randell, 1986:

	fibered	free	factor	fiber type	SS	LCS	rat'l $K(\pi, 1)$	formal	$K(\pi, 1)$	comb.
R-Refl.	??	т	т	F	F	F'	F 1)	т	т	NA
Simp.	?	F	F	F	F	F	F	т	т	T <sup>so</sup>
C-Refl.	??	т	т	F	F	$\mathbf{F}^{\prime \gamma}$	F	?	??	NA
fibered		F	F	F	F	F	F	т	т	?
free	?		т	F	F	F	F	??	??	??
factor	F	F		F	F	F	F	F	F	T
fiber type	Т	т	т		т	т	т	т	т	т
super- solvable	т	т	т	т		т	т	т	т	т
lower central series	?	F *	F *)	F 1)	F <sup>1)</sup>		F 1)	??	??	т
Rat'l K(π, 1)	?	F *:	F *)	F *)	F °	T 3)		т	??	т
Parallel	?	F	F	F	F	T *3	T *2	T *1	??	т
formal	F	F	F	F	F	F	F		F	. ?
simple	F	F	F	F	F	F	F	F	F	т
$K(\pi, 1)$	??	F	F	F	F	F	F	т		??
*) De wh 1) D <sub>4</sub> 2) C. 3) M.	rows are columns T = tr F = fa NA = no ? = no ?? = no pends o iich has r is not ra bloev con Toda, as Falk [8 mula, pr	are co ue lse of appl of kno of kno n the not app tional npletio comm ] and	icable wn wn, of s assertio eared. $K(\pi, 1)$ n of ge nunicate T. Kohr	ignifica n that [8], t neralize d by H	paral out is l ed pui [. Tera	lel im LCS [7] te brai 10.	. Kohn d group	o, Poinca s, prepri	aré serie: int].	s of the

They updated this in 1998, but this area is still very active (including last week!)

#### From Falk-Randell, 1986:

	fibered	free	factor	fiber type	SS	LCS	rat'l K(π, 1)	formal	<i>K</i> (π, 1)	comb.
R-Refl.	??	т	т	F	F	F?	F <sup>1</sup>	т	т	NA
Simp.	?	F	F	F	F	F	F	т	т	T <sup>s</sup>
C-Refl.	??	т	т	F	F	F '	F	?	??	NA
fibered		F	F	F	F	F	F	т	т	?
free	?		т	F	F	F	F	??	??	??
factor	F	F		F	F	F	F	F	F	T
fiber type	Т	т	т		т	т	т	т	т	т
super- solvable	т	т	т	т		т	т	т	т	т
lower central series	?	F *	F *)	F 1)	F <sup>1)</sup>		<b>F</b> <sup>1)</sup>	??	??	т
Rat'l K(π, 1)	?	F *	F *)	$F^{\left\{ 0\right\} }$	F *)	T <sup>3</sup>		т	??	т
Parallel	?	F	F	F	F	T *)	T *2	T *1	??	т
formal	F	F	F	F	F	F	F		F	. 7
simple	F	F	F	F	F	F	F	F	F	т
$K(\pi, 1)$	??	F	F	F	F	F	F	т		??
*) De wh 1) D <sub>4</sub> Ma 2) C. 3) M.	rows are columns T = tr F = fa NA = no ? = no ?? = no pends o ich has r is not ra ilcev con Toda, as Falk [8 mula, pr	are co ue lse of appl of kno of kno of kno n the tot app tional npletio comm ] and	icable wn wn, of s assertio eared. $K(\pi, 1)$ n of gen nunicate T. Kohn	ignifica n that [8], t seralize d by H	paral out is l ed pui l. Tera	lel imp LCS [T re braic ao.	. Kohn d group	o, Poinca s, prepri	aré serie: int].	s of the

They updated this in 1998, but this area is still very active (including last week!)

In This Talk: Give an easy-to-compute criterion to test if  $M(\mathcal{A})$  fails to be  $K(\pi, 1)$ .

## From Geometry to Algebra

## Overview

Typically, combinatorial properties are stronger than the K(π, 1) property

ex. supersolvable  $\Rightarrow K(\pi, 1)$ 

► Last year, Yoshinaga introduced a new one. It goes the other direction: K(π, 1) ⇒ Yoshinaga, but not the other way around

Unfortunately, this criterion is difficult to compute by hand.

 By translating Yoshinaga's criterion into an algebraic statement, we can use Gröbner bases to do calculations and formulate conjectues.

## Prep for Yoshinaga's Criterion

Let  $\mathcal{A} = \{H_1, \ldots, H_n\}$  a central, essential real arrangement with a fixed **orientation**, i.e. for each  $H_i$  we have two open halfspaces  $H_i^+$  and  $H_i^-$ . Let  $E = \{1, \ldots, n\}$ .

- A signed subset D of E is an ordered pair of disjoint subsets (D<sup>+</sup>, D<sup>-</sup>) of {1,...,n}.
- A signed subset D = (D<sup>+</sup>, D<sup>-</sup>) of E with D<sup>+</sup> ∪ D<sup>-</sup> = E is k-consistent if, for any subset S ⊂ E with #S ≤ k + 1 and any signed subset D = (D<sup>+</sup>, D<sup>-</sup>) with D<sup>+</sup> ∪ D<sup>-</sup> = S, we have

$$\bigcap_{i\in D^+}H_i^+\cap \bigcap_{i\in D^-}H_i^- \qquad \text{is nonempty}\,.$$

## Yoshinaga's Criterion

Let  $\mathcal{A} = \{H_1, \dots, H_n\}$  a central, essential real arrangement with a fixed orientation.

- $\Sigma_k = \{S \mid S \text{ is } k \text{-consistent}\}$
- $\blacktriangleright \sigma_k = \# \Sigma_k$

•  $\mathcal{A}$  is **Yoshinaga** of rank k if  $\sigma_k = \sigma_r$ , where r is the rank of  $\mathcal{A}$ .

#### Theorem (Yoshinaga, 2024)

Let A be an arrangement in a real vector space V. If its complexified complement M(A) is  $K(\pi, 1)$ , then A is Yoshinaga of rank 2.

## Yoshinaga's Criterion

For graphical arrangements, need an orientation of the edges of the graph so that no (k + 1)-subset forms an oriented cycle.

#### Example



Check: either way to have an oriented 4-cycle will force a "bad" orientation of one of the three cycles ⇒ Yoshinaga's criterion holds

#### Example

Prep for Algebraic Yoshinaga's Criterion

Let  $\mathcal{A} = \{H_1, \dots, H_n\}$  a central, essential, oriented real arrangement.

A signed dependence  $D = (D^+, D^-)$  is a signed subset of  $\{1, \ldots n\}$  such that

$$\bigcap_{i\in D^+} H_i^+ \cap \bigcap_{j\in D^-} H_j^- = \emptyset \,.$$

- ► The Varchenko–Gelfand ideal is an ideal of the polynomial ring Q[z<sub>1</sub>,..., z<sub>n</sub>] generated by
  - 1. Heaviside Relations:  $z_i(z_i 1)$  for  $i \in E$  and
  - 2. Dependence Relations:

$$\prod_{i \in D^+} z_i \prod_{j \in D^-} (z_j - 1) - \prod_{i \in D^+} (z_i - 1) \prod_{j \in D^-} z_j$$

for each signed dependence  $(D^+, D^-)$ .

## Algebraic Yoshinaga's Criterion

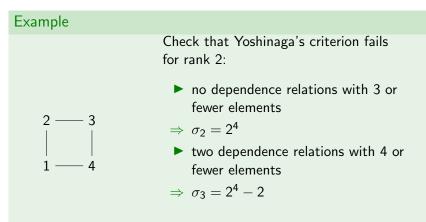
#### Theorem (Denham-DB-Proudfoot, 2025+)

Yoshinaga's criterion holds for rank k if and only if the entire Varchenko–Gelfand ideal is generated by Heaviside relations and dependence relations coming from dependencies of size k + 1 or less.

**Upshot**: Instead exhaustive searches of signed sets, we can check containment of ideals.

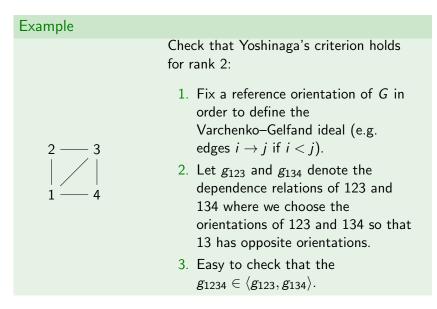
For the experts: this is a statement about an *inhomogeneous* ideal. The corresponding statement fails for the Cordovil relations!

## Algebraic Yoshinaga's Criterion



 $\sigma_2 \neq \sigma_3 \Rightarrow$  Yoshinaga fails at rank 2

## Algebraic Yoshinaga's Criterion



## Implementation (Sage vs Macaulay2)

## Why does run-time matter to us?

- There are partial characterizations of the K(π, 1) property for subarrangements of the type B reflection arrangement starts with Edelman–Reiner's counterexamples to Saito's conjecture, continued by Bailey, Proudfoot–Falk, and Suyama, Daisuke, Torielli, Tsujie
- This algebraic condition is the right thing to help us complete the characterization
- In order to figure out the right statement, we ran exhaustive searches (for small n) on the Whistler compute cluster at Texas A&M.
  - With the naïve implementation, this would have taken months or years.
  - With the algebraic version, it took us a few weeks.

### Two Implementations

#### Naïve version (Sage):

	_implementation.py 3, U × 🗅 🗸 🖽 🐇
🔹 naiv	e_implementation.py >
	import time
	<pre>def is_flat_consistent(X,topes,eps):</pre>
	INPUT:
	A – a real arrangement
	X – list of the hyperplanes containing a flat of A
	e - a sign vector
	OUTPUT:
	A boolean, true if e is consistent with the flat X and false otherwi
	EXAMPLES:
	<pre>sage: A = hyperplane_arrangements.Coxeter("B2"); topes = [list(A.sig</pre>
	<pre>sage: X = {0,1,2}; eps = [+1,+1,-1,+1]; is_flat_consistent(X,topes,e</pre>
	False
	<pre>sage: X = {0,1,2}; eps = [+1,+1,+1,+1]; is_flat_consistent(X,topes,e</pre>
	True
	<pre>sage: X = {0,1}; eps = [+1,+1,-1,+1]; is_flat_consistent(X,topes,eps</pre>
	True
	<pre>new.eps = [eps(k] for k in list(X)]</pre>
	new_eps = [eps(k) for k in list(k)] new topes = []
	for T in topes:
	new topes: new topes.append[[T[k] for k in list(X]])
	return new eps in new topes
	retorn new_eps in new_copes

#### Using rings (Macaulay2):

sitig ،	nore <sup>%6</sup> sigma.tex
≅ cod	e.m2
65	checkContainments = (A,R) -> (
69	firstTwo := byDegree0 + ideal for c in circs list
71	return all(circs, c -> ((# c == 3) or
72	(relationFromCircuit(nrmls,c) % firstTwo) == 0))
73	) is a second
74	
75	bigVGIdeal = (A,R) -> (
76	<pre>byDegree0 := ideal for i from 0 to # hyperplanes A = 1 list (e_i = 1) + e_i ;</pre>
77	nrmls := coefficients A;
78	circs := circuits A;
79	return byDegree0 + ideal for c in circs list relationFromCircuit(nrmls,c);
88	1;
81 82	shortVGIdeal = (A,R,k) -> (
83	idempoties := ideal for i from 0 to # (hyperplanes A) - 1 list (e i - 1) * e i :
84	nrmls := coefficients A:
85	circs := select(circuits A, c $\rightarrow$ (#c $\ll$ k + 1));
86	return idempoties + ideal for c in circs list(relationFromCircuit(nrmls.c)):
87	):
88	
89	shortCircuits= (A,R,k) -> (
98	<pre>shortCircs := select(circuits A, c -&gt; (#c &lt;= k + 1));</pre>
91	checkDigit := true;
92	while checkDigit == true do(
93	checkDigit = false;
94	for c in shortCircs do(
95	for d in shortCircs do[
96	if #(set(c) * set(d)) == 1 then(
97	<pre>newCirc := set(c) + set(d) - set(c)+set(d); binnestElement := set/max_talist/full(irch);</pre>

## Two Implementations

#### Macaulay2 version uses the "HyperplaneArrangements" pacakge

← → C O A https://macaulay2.com	/doc/Macaulay2/shar	☆ ♡	۵ ۵	=
Macaulav2 latest release (1.24.11) × a				
Documentation	Se	earch		
Packages * HyperplaneArrangements ::	next   previor	us   forward	backwa	rd
HyperplaneArrangements		up	index	toc
HyperplaneArrangements -	- manipulating	hyperpla	ine	
arrange	ments			
Description				
the hyperplanes. The tools provided allow the and to compute various algebraic invariants of Introductions to the theory of hyperplane arra textbooks:	of arrangements.	-		old,
<ul> <li>Alexandru Dimca, Hyperplane arrangen ISBN: 978-3-319-56221-6</li> </ul>	nents, Universitext,S	Springer, Ch	am, 2017	
<ul> <li>Peter Orlik and Hiroaki Terao, Arrangerr mathematischen Wissenschaften 300,S 978-3-662-02772-1</li> </ul>				
<ul> <li>Richard P. Stanley, An introduction to hy</li> </ul>				
Combinatorics, 389-496, IAS/Park City I Mathematical Society, Providence, RI, 2				
mainemalical acciety, Providence, RI, 2	UUT. IODIN: 978-1-4	+/04-3912-5		
Authors				
Autnors				
Graham Denham				

Avi Steiner <avi.steiner@gmail.com>

Vareio

#### Using rings (Macaulay2):

sitig ،	nore <sup>%6</sup> sigma.tex ≅ code.m2 ×
≡ cod	e.m2
65	checkContainments = (A,R) -> (
69	<pre>firstTwo := byDegree0 + ideal for c in circs list</pre>
71	return all(circs, c -> ((# c == 3) or
72	<pre>(relationFromCircuit(nrmls,c) % firstTwo) == 0))</pre>
73	);
74	
75	bigVGIdeal = (A,R) -> (
76	<pre>byDegree0 := ideal for i from 0 to # hyperplanes A - 1 list (e_i - 1) * e_i ;</pre>
77	nrmls := coefficients A;
78	circs := circuits A;
79 80	return byDegree0 + ideal for c in circs list relationFromCircuit(nrmls,c); ):
88	11
82	shortVGIdeal = (A.R.k) -> (
83	<pre>idempoties := ideal for i from 0 to # (hyperplanes A) - 1 list (e i - 1) * e i ;</pre>
84	nrmls := coefficients A:
85	circs := select(circuits A, $c \rightarrow (\phi c \ll k + 1))$ ;
86	return idempoties + ideal for c in circs list(relationFromCircuit(nrmls,c));
87	- Bi
88	
89	shortCircuits= (A,R,k) -> (
98	<pre>shortCircs := select(circuits A, c -&gt; (#c &lt;= k + 1));</pre>
91	checkDigit := true;
92	while checkDigit == true do(
93	checkDigit = false;
94	for c in shortCircs do(
95 96	<pre>for d in shortCircs dol     if #(set(c) * set(d)) == 1 then(</pre>
96 97	<pre>if #(set(c) * set(d)) == 1 then( newCirc := set(c) + set(d) - set(c)*set(d):</pre>
97	highertElement in cellenr to ist/full/istla

## Comparing Run-Times of the Two Implementations

Naïve version (Sage) - several minutes

```
sage: A = hyperplane_arrangements.braid(6)
sage: %time yoshinagas_criterion(A)
CPU times: user 2min 29s, sys: 510 ms,
total: 2min 30s
Wall time: 2min 30s
True
```

Using rings (in Macaulay2) - a few seconds

## Ring calculations are consistently fast

```
Using rings (in Macaulay2) - a few seconds
```

```
i2 : A = typeA(6, QQ); R = setupRing(A);
time checkContainments(A,R)
-- used 8.18428s (cpu); 4.16699s (thread); 0s (gc)
o4 = true
```

Using rings (in Macaulay2) - a few seconds

i2 : A = typeA(7, QQ); R = setupRing(A); time checkContainments(A,R) -- used 470.341s (cpu); 210.63s (thread); 0s (gc) o4 = true

## Where do we go from here?

## Successes so Far

- Fast check to determine if an arrangement *could* be  $K(\pi, 1)$
- Uses the Varchenko–Gelfand ring to gather information about the *homotopy groups* of M(A)

The Varchenko–Gelfand ring and its associated graded (aka the Cordovil algebra) are interesting in their own right, and we'll see them again tomorrow morning in Ayah and Sarah's talks.

- For a natural notion of chordal oriented matroid, we can show that chordal implies Yoshinaga's criterion
- In the case of subarrangements of the Type B reflection arrangement (including graphical arrangements, braid arrangements, threshold arrangements, etc), get more precise results because we can precisely characterize dependent sets

## Loose Threads

Are there other interesting families of arrangements for which we can say concretely whether Yoshinaga's criterion fails or not?

This is interesting for ranks  $\geq$  2, and has a topological interpretation.

Can we make the calculations more efficient by exploiting the symmetry of the arrangement? Interesting special case: graphical arrangements where the underlying graph is not chordal.

Can this criterion be "upgraded" to get a combinatorial condition equvivalent to K(π, 1)-ness?

## Thank you for your attention!